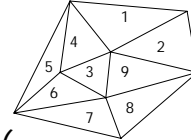


# Point location and Persistent Data structures

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## First problem

- Given a subdivision of the plane into triangles, create a data structure, such that given a query point, we can find in which triangle it lies.

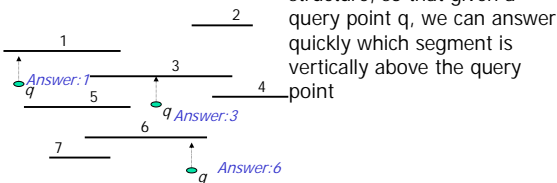


Too hard ;-(

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## A question that we would solve...

- A simpler problem



Given a set of horizontal segments, create a data structure, so that given a query point  $q$ , we can answer quickly which segment is vertically above the query point

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## A really really really easy question...

- Same question, but for a set of horizontal segments, all having the same x-coordinates.

(4,20) \_\_\_\_\_ (18,20)  
 (4,16) \_\_\_\_\_ (18,16)  
 (4,13)  $\uparrow$  Answer: 16 (18,13)  
 (4,7) \_\_\_\_\_ (18,7)

For simplicity, we use the y-coordinate of a segment as the "name" of the segment

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## Solution: Skip list

Solution: Store the y-coordinates of the segments in a SkipList.

Once a query point  $(x,y)$  is given, perform  $succ(y)$

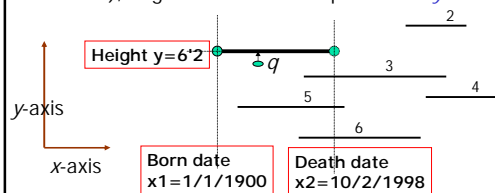
(4,20) \_\_\_\_\_ (18,20)  
 (4,16) \_\_\_\_\_ (18,16)  
 (4,13)  $\uparrow$  Answer: 16 (18,13)  
 (4,7) \_\_\_\_\_ (18,7)



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## Different question

- In a city, people are born, and die.
- Each person is recognized by its height (no two people have the same height). We denote height by the letter  $y$ , and (birth/death) date by  $x$ .
- Need a DS that would find  $Find(x,y)$  - who was the person that was alive at date  $x$ , and her/his height is  $y$  (or if not exists), larger and as close as possible to  $y$ .

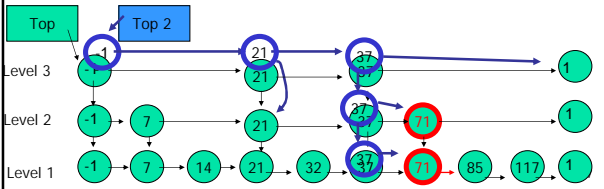


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## Different problem

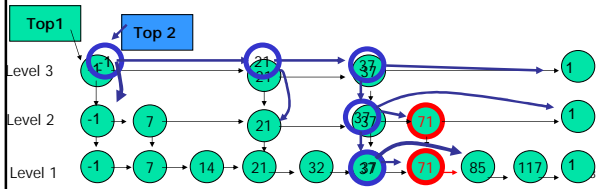
How to delete an element from the skipList, without destroying it?  
Assume we want to delete(71)

- Idea #1: Copy the whole SkipList, and delete - too much memory
- Idea #2: Copy the path that changes during the deletion, then modify this path.



## Virtually copying SL

- To create the new virtual copy:
- Start from **TOP** of the old SL, create a new top, named **Top2**
- Do **find(x)** /\* x is the key to be deleted \*/
- Copy and connect every element along the search path.
- Delete **x** from the SL pointed to by **Top2** /\*it does not affect the SL pointed to by **Top1** - only blue pointers change\*/
- Need to be a bit careful in the deletion (next slide)*



## Need a little new function

- `Follower( SL* sl, int x , int d){`
  - // Returns a pointer to the smallest cell at level
  - // d, with key strickly greater than x
- `P=sl->top; int d1=sl -> l ;`
- `while(1) {`
  - `while(p->key < x ) p = p->nxt ;`
  - `if( d1 == d ) return p ;`
  - `assert( p->down != NULL ) ; //add #include<assert.h>`
  - `p = p -> down ;`
  - `d1-- ;`
- `}`

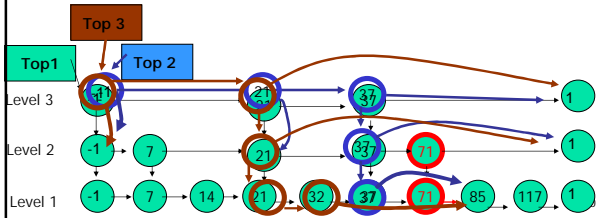
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## Virtually copying SL

To delete 37 -

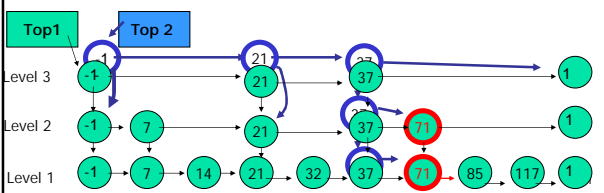
We copy as before the search path (brown path)

In each level d at which appear, we delete it using the command  
`p->nxt = follower( sl, 37, 3 )`



## Something fishy

- But it should not be.
- The new SL (lets call it SL(2) ) is a perfectly legal SL. It has the same keys of SL(1), excluding 71 that was deleted.

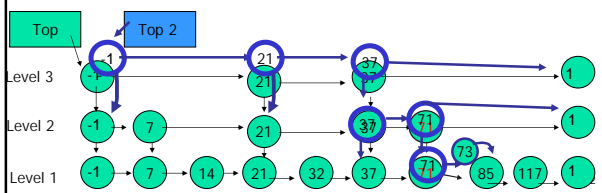


## Inserting a key

How to insert a key into the skipList, without destroying it ?

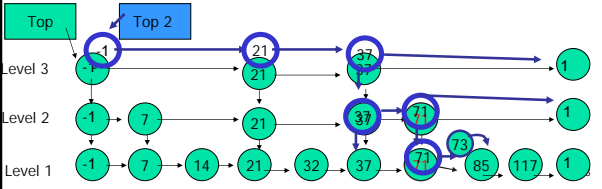
Same idea: Assume we want to **insert(75)**

- Do **search(75)**.
- Copy every element that the search goes through.
- Let **Top2** point to the top of the list.
- Insert(72)** into the SL pointed to by **Top2** - only blue elements change



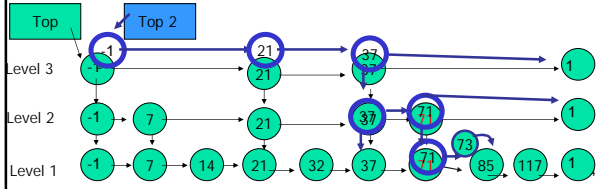
## Again - a brandnew SL

- Note - again we obtained a perfectly legal SL.
- We have two SkipLists - one contains 73, the other one does not contain 73.
- We can now insert/delete elements into/from **SL(2)**
- Remember:** to access a SL, one only need the **root** - the **top**.



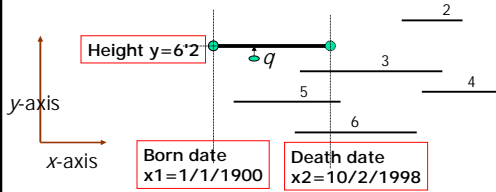
## How much space

- We saw that the average length of a path is  $O(\log n)$ , so each insert/delete takes  $O(\log n)$  time and space



## Back to the birth/death question

- In a city, people are born, and die.
- Each person is recognized by its height (no two people have the same height). We denote height by the letter  $y$ , and (birth/death) date by  $x$ .
- Need a DS that would find  $Find(x,y)$  - who was the person that was alive at date  $x$ , and her/his height is  $y$  (or if not exist), larger and as close as possible to  $y$ .



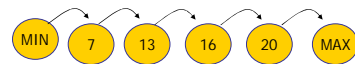
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## And remembering that this one is easy...

- All births/deaths start at the same date
- Call this problem the *same-population problem* (no births no deaths)

(4,20) (18,20)  
 (4,16) (18,16)  
 (4,13) **Answer: 16** (18,13)  
 (4,7) (18,7)

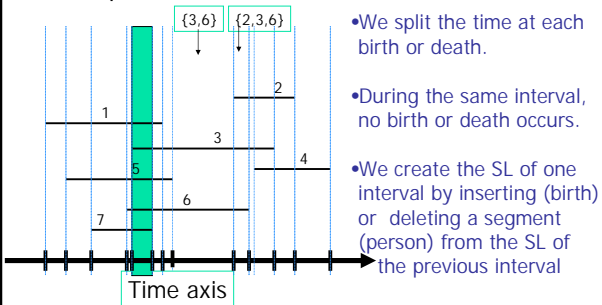
Easily solved via standard skip list



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## We can solve this one by...

- We split the time axis into *time-intervals*.
  - We split the time at each birth or death.
  - During the same interval, no birth or death occurs.
  - We create the SL of one interval by inserting (birth) or deleting a segment (person) from the SL of the previous interval



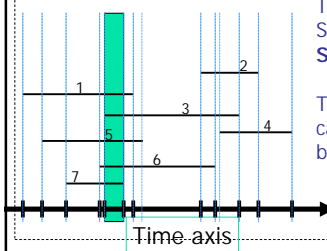
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## How to access the different SL ?

So we obtain  $2n+1$  SkipLists, (one for time intervals).

The roots of all different SL are Stored in a new SL (call it **dates SL**), sorted by date.

The SkipLists of the people are called **People SL**, sorted by height.

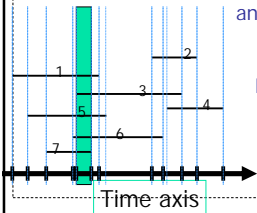


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## How to access the different SL ?

Alg: Create an empty people-SL  
Create the dates-SL, and insert each birth/death event.

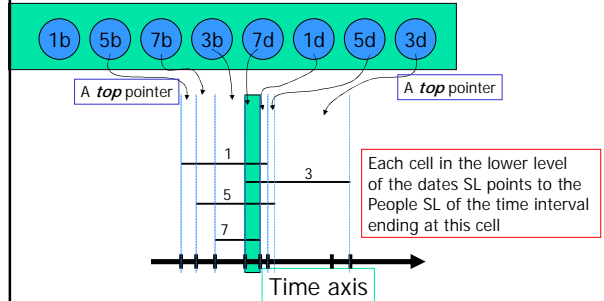
Scan these event (in increasing dates),  
and for each event (date) do  
Insert/delete (virtually) into/from the  
people SL.  
Link the TOP of the new people-SL  
with the current key of the  
dates-SL



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## Overall data structure

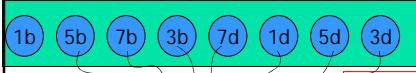
- Dates SL (only lower level shown)



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## Answering a query

A query point  $(x,y)$  - which segment is vertically above the point  $(x,y)$



Use the people SL to find the time interval contain  $x$ , and the corresponding people SL (use *successor(x)* operation in Dates SL).

Perform *successos(y)* operation in this people SL to find the segment above  $(x,y)$ .

Time axis

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