Subdivision Tree Representation of Arbitrary Triangle Meshes

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January 15, 1998
Abstract

We investigate a new way to represent arbitrary triangle meshes. We prove that a large class of triangle meshes, called normal triangle meshes, can be represented by a subdivision tree, where each subdivision is one of four elementary subdivision types. We also show how to partition an arbitrary triangle mesh into a small set of normal meshes. The subdivision tree representation can be used to encode mesh connectivity information. Our theoretical analysis shows that such a coding scheme is very promising.
0.1 Introduction

Triangle meshes are popular in graphics and other application areas. One problem in dealing with meshes is to store topological information or connectivity of meshes. Recently, some methods have been presented to efficiently code meshes [5, 3, 4, 13, 6, 9]. A recent survey can be found in [12].

A triangle mesh has two different parts: topological information and geometric information. The topological information (also called connectivity) is about the relationship between vertices, edges and triangles in the mesh, while the geometrical information is about the position of vertices. Usually, the topological information is independent from geometric information, except some special cases such as regular grid, or Delaunay Triangulations, where geometric information is able to determine connectivity.

Most current methods try to generate long triangle strips [5, 3, 13, 1] to code mesh connectivity so that they can be used in some graphical library such as OpenGL. Some methods use vertex split/merge to code topological information implicitly [6, 9]. The hierarchical triangulation presented in [4] discussed the representations of meshes in tree structures. A recursive refinement of meshes naturally forms tree representations of triangle meshes. However, paper [4] uses links to code vertices, edges and triangles explicitly. It may result in large storage in storing these links (indices for vertices, edges and triangles).

It is true that a recursive refinement has a tree representation. But for a given mesh, whether we can generate its tree representation is not obvious. We have not seen any work on this reverse problem. The main purpose of this paper is to answer solve this problem. We are going to find the necessary and sufficient conditions for a triangle mesh to be represented by a tree.

This paper is organized as follows: In Section 2 gives basic terminologies. Our main results are presented in Section 3.

0.2 Notation and Definition

We first list some topological concepts used in our study. Some terms are defined in $\mathbb{R}^n$ in their general forms, although we are only interested in orientable bounded
surfaces in $\mathbb{R}^3$. For details about these terms and related information, readers may refer to topology textbooks, such as [8].

**Definition 1** The Euclidean distance between $x = (x_1, \ldots, x_n)$ and $y = (y_1, \ldots, y_n)$ in $\mathbb{R}^n$ is given by
\[
\|x - y\| = \sqrt{(x_1 - y_1)^2 + \cdots + (x_n - y_n)^2}
\]

**Definition 2** A topological space is a set $X$ with a collection $\mathcal{B}$ of subsets $N \subseteq X$, called neighborhoods, such that

- every point is in some neighborhood, i.e., $\forall x \in X, \exists N \in \mathcal{B}$ such that $x \in N$
- $N_1, N_2 \in \mathcal{B}$ with $x \in N_1 \cap N_2, \exists N_3 \in \mathcal{B}$ such that $x \in N_3 \subseteq N_1 \cap N_2$.

The set, $\mathcal{B}$, of all neighborhoods is called a basis for the topology on $X$.

**Definition 3** Two topological spaces $A$ and $B$ are homeomorphic if there is a continuous invertible function $f: A \to B$ with continuous inverse $f^{-1}: B \to A$. Such a function $f$ is called a homeomorphism.

**Definition 4** An $n$-cell is a set whose interior is homeomorphic to the $n$-dimensional disc $D^n = \{x \in \mathbb{R}^n : \|x\| < 1\}$ with the additional property that its boundary must be divided into a finite number of lower-dimensional cells, called the faces of the $n$-cell.

- A $0$-dimensional cell is a point.
- A $1$-dimensional cell is a line segment.
- A $2$-dimensional cell is a triangle.
- A $3$-dimensional cell is a tetrahedron.

**Definition 5** A complex $K$ is a finite set of cells, $K = \cup\{\sigma : \sigma \text{ is a cell}\}$ such that:

- if $\sigma$ is a cell in $K$, then all faces of $\sigma$ are elements of $K$.
- if $\sigma$ and $\tau$ are cells in $K$, then $\text{Int}(\sigma) \cap \text{Int}(\tau) = \emptyset$. 

where $\text{Int}(A)$ denotes the interior of a set $A$. The dimension of $K$ is the dimension of its highest-dimensional cell.

**Definition 6** An $n$-dimensional manifold is a topological space such that every point has a neighborhood homeomorphic to an $n$-dimensional open disc

$$D^n = \{ x \in \mathbb{R}^n : \| x \| < r \}$$

We further require that any two distinct points have disjoint neighborhoods. A 2-manifold is often called a surface.

**Definition 7** A $n$-dimensional manifold with boundary is a topological space such that every point has a neighborhood homeomorphic to either a $2$-dimensional open disc or the half-disc $D^n_+ = \{ x = (x_1, \ldots, x_n) \in \mathbb{R}^n : \| x \| < r, x_n \geq 0 \}$. Points with half-disc neighborhoods are called boundary points.

**Definition 8** Let $K$ be a complex. The set of all points in the cells of $K$ is

$$|K| = \{ x : x \in \sigma \in K, \sigma \text{ is a cell in } K \}$$

is the space underlying the complex $K$, or the realization of $K$.

**Definition 9** Let $S_1$ and $S_2$ be two surfaces. Remove a small disc from each of $S_1, S_2$ and glue the boundary circles of these discs together to form a new surface called the connected sum of $S_1$ and $S_2$, written as $S_1 \# S_2$.

**Definition 10** A locally $2$-dimensional topological space $X$ is triangulizable if a $2$-complex structure $K$ can be found with $X = |K|$ and $K$ has only triangle cells satisfying the additional condition that any two triangles are identified along a single edge or at a single vertex or are disjoint. A triangulated complex $K$ is called a simplicial complex or a triangulation on $X$. A cell of a simplicial complex is called a simplex.

**Definition 11** Let $K$ be a complex. The Euler characteristic of $K$ is

$$\chi(K) = \#(0\text{-cells}) - \#(1\text{-cells}) + \#(2\text{-cells}) - \#(3\text{-cells})\ldots$$

where $\#(S)$ denotes the number of elements in a finite set $S$. 
Definition 12 If $S$ is a surface with boundary, the associated surface (without boundary) is $S^*$, where $S^*$ is $S$ with a disc sewn onto each of the boundary circles.

In other words, $S^*$ is the surface $S$ with all its holes patched. If $S$ has $k$ holes or boundary components, then $S^*$ can be considered as $S$ with $k$ discs or 2-cells added so $\chi(S^*) = \chi(S) + k$.

Definition 13 The genus $g$ of a closed surface of Euler characteristic $\chi$ is given by $(2 - \chi)/2$ if $\chi$ is even, and $(1 - \chi)/2$ if it is odd.

The genus of a surface counts how many "handles" the surface has. For orientable surfaces, this is the number of "holes" in the surface.

Definition 14 The genus of a compact surface with boundary $S$ is $g(S) = g(S^*)$ where $S^*$ is the associated surface without boundary.

Theorem 1 Every compact connected orientable surface is homeomorphic to a sphere or a connected sum of $n$ tori.

Intuitively, the orientable surfaces are the sphere and tori with any number of handles. Since a torus is a sphere with a handle, every orientable surface can be described as a sphere with some number of handles. Because an orientable surface with genus $n$ is $n$ tori glued together, and a torus is two closed discs glued together, an orientable surface can be generally decomposed into $2n$ closed discs. A genus 0 surface (sphere) is two closed discs glued together. Based on these facts and for the sake of simplicity, we are going to focus on 2D surfaces homeomorphic to a closed disc, i.e., surfaces in 2D with simple boundaries and no holes.

From our previous definition of triangulation, we have the following definition for 2D triangle meshes:

Definition 15 A 2D triangle mesh is a collection of disjoint triangles in a 2D plane such that they cover a connected region in the plane, and a triangle is not allowed to have a vertex of another triangle in the interior of one of its edge.
In the following, we assume $M$ is a 2D triangle mesh with a simple boundary and no holes. We denote $\partial M$ as the boundary (polygon) of $M$. For a pair of adjacent vertices $v_1, v_2 \in M$, the edge connecting them is denoted by $\overline{v_1v_2}$. A path in $M$ consists of edges $\overline{v_1v_2, \ldots, v_{n-1}v_n}$ is denoted as $<v_1, \ldots, v_n>$. For simplicity, a path connecting $v_1, v_2$ is also denoted as $v_1v_2$ if there is no confusion as to which path we are talking about. The length of a path is defined as the number of vertices forming the path.

**Definition 16**: A graph $G = (V, E)$ consists of two sets: a finite set $V$ of elements called vertices and a finite set $E$ of vertex pairs called edges.

Generally speaking, a 2D triangulation cannot be represented by a graph non-ambiguously. For instance, when a triangle mesh has a hole which is a single triangle, the mesh has the same graph as the mesh without the hole.

**Definition 17**: A dual graph of mesh $M$ Dual$(M)$ is a 2D graph obtained as follows. For each triangle of $M$ we have a vertex in Dual$(M)$. If two triangles share a common edge $e$ in $M$ then the corresponding vertices in Dual$(M)$ will be adjacent, i.e., connected by an edge $\hat{e}$. We call $\hat{e} \in$ Dual$(M)$ the dual of $e \in M$ (Fig. 1-1).

![Fig. 1-1. M (in solid lines) and its dual Dual(M) (in dotted lines)](image)

**Definition 18**: Two graphs $G_1$ and $G_2$ are said to be isomorphic if there exists a one-to-one correspondence between their vertex sets and a one-to-one correspondence between their edge sets such that the corresponding edges of $G_1$ and $G_2$ are incident on the corresponding vertices of $G_1$ and $G_2$. 
Definition 19: Two triangle meshes $M_1$ and $M_2$ are called topologically equivalent if their dual graphs are isomorphic.

Definition 20 All vertices/edges on $\partial M$ are called boundary vertices/edges, the rest vertices/edges in $M$ are called interior vertices/edges.

It is easy to see that a boundary edge is contained by only one triangle in $M$, while an interior edge is shared by exactly two triangles in $M$. An edge is interior in $M$ iff it has a dual edge in Dual$(M)$.

Definition 21 For a vertex $v$ in a graph $G$, the number of edges incident to $v$ is called the valence (or degree) of $v$, denoted by $\rho(v)$.

Obviously, $\rho(v) \geq 3$ for an interior vertex $v$, and $\rho(v) \geq 2$ for a boundary vertex $v$.

Definition 22 For a vertex $A \in \partial M$, $A$ is called subdivisible if $\rho(A) > 2$.

Definition 23 For a point $A \in \partial M$, we define its cone $\text{cone}(A, M)$ as

$$\text{cone}(A, M) = \{ \triangle AV_1V_2 | \triangle AV_1V_2 \in M \}$$

i.e., $\text{cone}(A, M)$ is the collection of all triangles (in $M$) containing vertex $A$. The collection of all edges in $\text{cone}(A, M)$ excepting those containing vertex $A$ is called the frontier of $\text{cone}(A, M)$.

The frontier of $\text{cone}(A, M)$ is a connected path. Fig. 1-2 shows a cone of vertex $A$ and its frontier $< U1, U2, U3, U4, U5, U6 >$.

![Fig. 1-2. A cone and its frontier](image)
**Definition 24** A pair of vertices $v_1$ and $v_2$ in $M$ are called interior-connected if there exists $k (k \geq 0)$ vertices $w_1, \ldots, w_k$ such that $w_i$'s $(i = 1, \ldots, k)$ are distinct interior vertices in $M$, and $\overline{v_1w_1} \subseteq \overline{w_jw_{j+1}}$ $(i = 1, \ldots, k-1)$, and $\overline{w_kw_2}$ are interior edges in $M$. The path $< v_1, w_1, \ldots, w_k, v_2 >$ is called an interior-path between $v_1$ and $v_2$ in $M$.

Based on this definition, we immediately have the following property.

**Property 1** A path $< v_1, w_1, \ldots, w_k, v_2 >$ is not an interior path iff there exists a vertex $w_j$ $(1 \leq j \leq k) \in \partial M$ when $k > 0$ or edge $\overline{v_1v_2}$ is on $\partial M$ (when $k=0$).

Because $M$ is a single triangulated mesh without holes, all vertices are connected (not necessarily interior-connected). In addition, for any pair of vertices in $M$, there are at least two paths (possibly partially overlapping) in $M$ connecting them due to the fact that any vertex has valence larger than one.

**Definition 25** For a pair of vertices $v_1, v_2 \in M$, an edge $e$ in $M$ is called a blocking edge for $v_1$ and $v_2$ (w.r.t. $M$) if the edge $e$ cuts $M$ into two disjoint sub-triangle meshes $M_1, M_2 \subset M$ such that $v_1 \in M_1, v_1 \notin M_2, v_2 \in M_2, v_2 \notin M_1, M_1 \cap M_2 = e$. A vertex $v \in M$ is said to have no blocking edge in $M$ if there is no blocking edge for $v$ and any other vertex $w \in M$; otherwise, we say $v$ has a blocking edge in $M$.

In the dual space $\text{Dual}(M)$, removing the dual edge $\bar{e}$ of the blocking edge $e$ will result in two separated subgraphs of $\text{Dual}(M)$.

**Lemma 1** An edge $e$ is a blocking edge for some vertices $v_1, v_2 \in M$ iff $e$ is an interior edge and both of its end points are on $\partial M$.

**Proof:**

$\Rightarrow$: Assume $e$ is a blocking edge for some vertices $v_1, v_2 \in M$.

If $e$ is a boundary edge, cutting $M$ along $e$ will generate two meshes $M_1 = M$, and $M_2 = e$, thus, both $v_1, v_2 \in M_1$. This contradicts the definition of blocking edge of $e$. Thus, $e$ must be an interior edge.

If at least one of the two end points of $e$ is interior to $M$, cutting along $e$ is not able to separate $M$ into two disjoint parts. Thus, both of the two end points of $e$ must be on $\partial M$. 


If $e$ is an interior edge and both of its endpoints $u_1, u_2$ are on $\partial M$, then $e$ cuts $M$ into two disjoint sub-meshes $M_1, M_2$; each of them has at least one triangle. Let $v_1 \notin e$ be a vertex in $M_1$, and $v_2 \notin e$ be a vertex in $M_2$, then $e$ is a blocking edge for $v_1, v_2$.

\[ \square \]

**Lemma 2** For any pair of non-adjacent vertices $v_1, v_2 \in M$, either there exists an interior-path between $v_1$ and $v_2$ or there is a blocking edge $e$ for $v_1$ and $v_2$.

Proof: If there is a blocking edge $e$ for the vertices $v_1$ and $v_2$, then $v_1$ and $v_2$ are in disjoint submeshes $M_1$ and $M_2$. All paths connecting $v_1$ and $v_2$ must pass through one of the two end points of $e$, which are boundary points. Thus, there is no interior path for $v_1, v_2$.

If there is no blocking edge for non-adjacent vertices $v_1$ and $v_2$, there must be an interior path between $v_1$ and $v_2$. If not, any path connecting $v_1$ and $v_2$ has at least one vertex which is on $\partial M$. Assume $m > 0$ is the minimum number of boundary vertices on all such paths, and the path $\mathcal{P} = v_1, w_1, \ldots, w_k, v_2$ has $m$ vertices $w_{j_1}, \ldots, w_{j_m}$ where $(j_1 < j_2 < \ldots < j_m)$ are on $\partial M$. Without loss of generality, we assume the path $\mathcal{P}$ has the shortest length among such paths.

For convenience, we denote $w_0 = v_1, w_{k+1} = v_2$. Consider the angle (inside $M$) bounded by two edges $\overline{w_{j_1-1}w_{j_1}}, \overline{w_{ji}w_{j_i+1}}$. Assume $\overline{w_{j_1}w_{j_1-1}}, \overline{w_{ji}u_1}, \ldots, \overline{w_{ji}u_r}, \overline{w_{ji}w_{j_i+1}}$ are all edges inside the angle, where $r \geq 0$ (See Fig. 1-3 (a)).

![Diagram](a) ![Diagram](b)

Fig. 1-3

The vertices $u_1, u_2, \ldots, u_r$ must be interior, otherwise, say, $u_s$ ($1 \leq s \leq r$) is on the boundary. Then the edge $\overline{w_{ji}u_s}$ cuts $M$ into two separated parts $M_1$ and $M_2$.  

\[ \naming{\text{Lemma}} \]

\[ \square \]
If \( v_1 \) and \( v_2 \) are not in the same part, \( \overline{w_j u} \) is a blocking edge for \( v_1, v_2 \). This is a contradiction.

If \( v_1, v_2 \) are in the same part (See Fig. 1-3 (b)), then the path \( P \) must pass through \( u_s \) and \( w_t = u_s \) for some \( t > j_1 + 1 \). Therefore, we can shorten the path \( P \) by replacing subpath \( \langle w_{j_1}, w_{j_1+1}, \ldots, w_t \rangle \) by a single edge \( \overline{w_j w} \). This contradicts the shortest length assumption of \( P \).

Thus, all vertices \( u_1, u_2, \ldots, u_r \) must be interior in \( M \). We construct a new path

\[
P' = \langle w_0, w_1, \ldots, w_{j_1-1}, u_1, u_2, \ldots, u_r, w_{j_1+1}, \ldots, w_{k+1} \rangle,
\]

which has only \( m - 1 \) vertices \( v_{j_2}, \ldots, v_{j_m} \) on \( \partial M \). This contradicts the assumption that \( m \) is minimized. Thus, \( v_1, v_2 \) are interior-connected. \( \square \)

Remark: If \( v_1 \) and \( v_2 \) are adjacent, it might be that neither a blocking edge nor an interior-path exists for them. An example is shown in Fig. 1-4.

![Fig. 1-4](image)

**Definition 26** If \( A, B, C \) are three vertices on \( \partial M \), we call \( M \) a triangular patch, denoted by \( P_{\Delta ABC, M} \). The boundary part of \( \partial M \) between \( A \) and \( B \) is called the patch edge \( AB \), denoted by \( PE(A, B) \). Similarly, we have \( PE(A, C), PE(B, C) \). The vertices on \( PE(A, B), PE(A, C), PE(B, C) \) are called the boundary vertices of the patch. The vertices \( A, B, C \) are called the corners of the patch.

Note: For simplicity, we still use \( M \) and \( P_{\Delta ABC, M} \) interchangeably if no confusion is caused.
Definition 27 A edge $e = \overline{u_1u_2}$ in a triangular patch $P_{\Delta ABC,M}$ is called illegal if it is interior and both of its end points are on same patch edge $PE(A, B)$, $PE(B, C)$, or $PE(A, C)$. Otherwise, is called legal.

Definition 28 A triangular patch $P_{\Delta ABC,M}$ is normal if it contains no illegal edges. Otherwise, the patch is called abnormal. If there exists a normal triangular patch $P_{\Delta ABC,M}$ for a mesh $M$, we call $M$ normal, otherwise, abnormal.

Fig. 1-5 show two examples of abnormal triangular patches.

![Fig. 1-5. Examples of abnormal triangular patches](image)

It is easy to see that a sufficient (but not necessary) condition to be a abnormal triangular patch is that valence $\rho(v) = 2$ for some boundary vertex $v$ which is not a corner.

As we mentioned earlier, a surface can be decomposed (cut along some paths on the surface) into a finite number of closed discs. For a 3D triangle mesh, we may choose the cutting paths to be paths of the mesh, i.e., we can cut a 3D triangle mesh into a number of sub-meshes, each of them topologically equivalent to a 2D disc. We should point out that we can select our cutting paths so that no illegal edges will be created by our cutting.
Fig. 1-6. Removing illegal edges by adjust cutting pathes

Fig. 1-6 shows the strategy for avoiding generating illegal edges. In Fig. 1-6, assume we cut a sphere-like triangle mesh into two submeshes $M_1$ and $M_2$. If there is an illegal edge $p = V_1V_4$ (shown as a wide-solid line) on the cutting path (shown as solid lines on the front and dotted lines on the back), we can change our cutting path by replacing sub-path $< V_1, V_2, V_3, V_4 >$ by the edge $p$. The edge $p$ is no longer illegal. Based on this strategy, when we cut a 3D triangle mesh without boundary into several sub-meshes, we can assume that all sub-meshes are topologically equivalent to normal triangular patches.

When a surface has boundaries, we cannot remove illegal edges simply by adjusting cutting edges if illegal edges exist on the existing boundaries. In this case, we may use a construction based on adding a vertex to get rid of all illegal edges immediately as shown in the following example.

Let $M$ be an abnormal 2D triangular patch with the boundary $\partial M = \{V_6, V_1, ..., V_{11}\}$ and three corners $A = V_6, B = V_4, C = V_8$ (Fig. 1-7(a)). The patch $P_{\text{DABC}M}$ has two kinds of illegal edges: those caused by degree-2-vertices – edges $V_6V_2, V_2V_4, V_4V_6, V_6V_8, V_8V_{10}, V_{10}V_6$, and those not caused by degree-2-vertices – edge $V_8V_{10}$. We add a vertex $W$ outside the plane of support of $M$ and construct triangles $\triangle VWV_{i+1}$ for all successive boundary points $V_i, V_{i+1}$ (shown as dotted lines in Fig. 1-7(a)). This results in a 3D triangle mesh denoted by $\tilde{M}$ which is topologically equivalent to a sphere. We consider $\partial M$ as the path cutting $\tilde{M}$ into two sub-meshes, one is the mesh $M$, the other is formed by $W$ with all vertices on $\partial M$. Now, we adjust the cutting path as mentioned above to remove illegal edges. In the example here, our new cutting path is shown in bold lines (Fig. 1-7(a)). Under the new cutting, the mesh $\tilde{M}$ is decomposed into two submeshes $M_1$ and $M_2$, as shown in Fig 1-7(b) and
(c) respectively. Both of them are normal patches.

Fig 1-7 (a)  

Fig 1-7 (b)  

Fig 1-7 (c)  

Removal of illegal edges

Note: The advantage of this technique is that we only need to code two sub-patches if there are many illegal edges. The major drawback is that it may create many new triangles in $M_2$, although the topological information is very cheap to store (we will address this problem later). An alternative way to normalize triangular patches is to cut a patch into many sub-patches so that each sub-patch is normal, but it may reduce the efficiency of coding.

0.3 General subdivision theorem

**Theorem 2** Let $P_{\Delta ABC,M}$ be a normal triangular patch obtained from triangle mesh $M$. We can reconstruct another triangle mesh $\tilde{M}$ by adaptively subdividing a single triangle $\Delta \tilde{A}\tilde{B}\tilde{C}$ such that each subdivision is one of the following four elementary types (called binary, ternary, quaternary, and mitsubishi, respectively) as shown in Fig 2-1(a)(b)(c)(d), $\tilde{M}$ is topologically equivalent to $M$, and the vertices $\tilde{A}, \tilde{B}, \tilde{C}$ correspond to the vertices $A, B, C$. 
Fig. 2-1. Four different ways to subdivide a triangle

Proof: We have two steps in our proof.

In step 1, we prove that $P_{\Delta ABC,M}$ can be subdivided into several sub-triangular patches using one of the four elementary subdivisions. In the second step, we prove the sub-triangular patches generated in the first step can be normalized. Therefore, we can repeat the subdivision on the sub-patches until all sub-patches have only three vertices, i.e., a single triangle.

0.3.1 Step 1

There are two cases.

Case 1: At least one vertex of $A, B, C$ has no blocking edge in $M$

Assume the vertex $A$ has no blocking edges, in other words, for any point $v \in M$, there exists an interior path from $A$ to $v$.

- Case 1.1: If there is a point $v$ on the boundary $PE(B, C)$, then we can subdivide $M$ along the interior path from $A$ to $v$ using ’binary subdivision’ (Fig. 2-2 (a)) and finish our proof.

- Case 1.2: Suppose the patch edge $PE(B, C)$ is a single edge $BC$. Let $\Delta DBC$ be the triangle in $M$ containing the edge $BC$. The vertex $D$ must be an interior point of $M$. Otherwise, either $BD$ or $CD$ will be a blocking edge for the vertex $A$, and thus it contradicts our assumption.

Because $D$ and $A$ are inner-connected by a path $P$, we can subdivide $M$ into three sub-triangular patches by cutting edge $BD, CD$ and the path $P$ (Fig. 2-2 (b)) using the ternary subdivision.
Case 2: All three vertices A, B, C have blocking edges in $M$.

Let edge $A_1A_2$ be a blocking edge for vertex $A$. We can assume that vertices $A_1$ and $A_2$ are on $PE(A, B)$ and $PE(A, C)$, respectively. If this condition is not satisfied, then one of the two vertices $A_1, A_2$, say, $A_1$, must be on the $PE(B, C)$. That leads to $A_1 = B$ and then $B, A_2$ are interior-connected.

We run over all such blocking edges, and assume the blocking edge $A_1A_2$ is selected such that the area $AA_1A_2$ is maximized.

Similarly, we have blocking edges $B_1B_2$ and $C_1C_2$ for vertices $B$ and $C$, which maximize area $BB_1B_2$ and $CC_1C_2$, respectively (Fig 2-3).

We call this assumption the maximum area assumption in the following.

We denote by $M'$ the region bounded by $A_1B_1B_2C_2C_1A_2$.

Based on the maximum area assumption, we immediately get the following result: there are no blocking edges inside the region $M'$, and any interior point $v \in M'$ is interior-connected to all vertices of $A_1, A_2, C_2, C_1, B_2, B_1$. 
If one of the following happens: $A_1 = B_1$ or $A_2 = C_1$ or $B_2 = C_2$, we say there is a degeneracy. We can deal with the degeneracy as follows. If, say, $A_1 = B_1$ (Fig. 2-4), then Lemma 2 implies that $A_2$ and $B_2$ must be interior-connected by a path $P$ in the region bounded by $A_1A_2CB_2$. Otherwise, there will be a blocking edge $A_1Q$ such that $Q \in A_2C$ or $Q \in B_2C$ (shown by a dotted line in Fig. 2-4), which contradicts the maximum area assumption. Thus, we can subdivide $M$ along path $P$ and two blocking edges $A_1A_2$ and $A_1B_2$ using quaternary subdivision.

![Fig. 2-4](image)

From now on, we assume there is no degeneracy, i.e., $A_1 \neq B_1$, $A_2 \neq C_1$ and $B_2 \neq C_2$.

Consider the triangles $\triangle A_1A_2A_3$, $\triangle B_1B_2B_3$ and $\triangle C_1C_2C_3$ in the region $M'$. These triangles exist because of the maximum area assumption and the triangulation of $M'$. In addition, the maximum area assumption guarantees that all three points $A_3$, $B_3$, $C_3$ are interior points in $M'$. Based on Lemma 2, it follows that there exist interior-paths $P_1$, $P_2$, $P_3$ inside $M'$ connecting $A_3$ and $B_3$, $A_3$ and $C_3$, $B_3$ and $C_3$, respectively (Fig. 2-5).

![Fig. 2-5](image)
We let $M_{A,B}$ denote the region bounded by $\overline{A_1A_3}$, $\overline{P_2B_1}$ and $A_1B_1$; $M_{A,C}$ bounded by $\overline{A_2A_3}$, $\overline{P_2C_1}$ and $A_2C_1$; $M_{B,C}$ bounded by $\overline{C_2C_3}$, $\overline{P_3B_2}$ and $B_2C_2$ (Fig 2-5).

We assume paths $P_1$, $P_2$, $P_3$ do not cross, although in general they may overlap. (If the paths cross, we can always re-organize the paths to remove the crossing.) We denote the three possible overlapping sub-paths by $A_3A_4$, $B_3B_4$ and $C_3C_4$ (Fig 2-6). We assume the lengths of these overlapping sub-paths are minimized.

Claim: There exist six edges (Fig. 2-6) $\overline{A_4S_1}$, $\overline{C_4S_2}$, $\overline{B_4T_1}$, $\overline{C_4T_2}$, $\overline{A_4T_1}$, $\overline{B_4T_2}$ such that

- $S_1, S_2 \in A_2C_1$, $T_1, T_2 \in A_1B_1$, $R_1, R_2 \in B_2C_2$,
- $\overline{A_4S_1}, \overline{C_4S_2} \in M_{A,C}$,
- $\overline{B_4T_1}, \overline{C_4T_2} \in M_{B,C}$,
- $\overline{A_4T_1}, \overline{B_4T_2} \in M_{A,B}$.

We now prove the existence of the edge $\overline{A_4S_1}$. If $A_3 = A_4$, we immediately get $S_1 = A_2$ and finish the proof. If $A_3 \neq A_4$, the angle $A_4$ in the region $M_{A,C}$ must be subdivisible. Otherwise, we can reduce the length of path $A_3A_4$, which contradicts our shortest length assumption on $A_3A_4$ above. Now build the cone of $A_4$ inside the region $M_{A,C}$. The frontier of the cone must have a vertex $S_1$ on the boundary $A_2C_1$, otherwise, the frontier can be used to reduce the length of $A_3A_4$.

Similarly, we can prove the existence of the remaining five edges.
It is easy to see that the regions $M_{A,B}, M_{A,C}, M_{B,C}$ are disjoint, but share parts of their boundaries.

Based on Fig. 2-6, we will find subdivision paths in $M$. The problem can be solved when we classify $M$ into the following three cases.

- **Case 2.1:** $A_4 \neq B_4, A_4 \neq C_4, C_4 \neq B_4$. (Fig. 2-7 (a)).

  We claim that there exist six paths $U_1U_2, U_1U_3, V_1V_2, V_1V_3, W_1W_2, W_1W_3$ such that
  - vertex $U_1 \in S_1S_2$, vertices $U_2, U_3 \in A_4C_4$, $U_2 \neq U_3$,
  - vertex $V_1 \in T_1T_2$, vertices $V_2, V_3 \in A_4B_4$, $V_2 \neq V_3$,
  - vertex $W_1 \in R_1R_2$, vertices $W_2, W_3 \in B_4C_4$, $W_2 \neq W_3$,
  - paths $U_1U_2, U_1U_3$ are interior paths in region $A_4C_4S_2S_1$,
  - paths $V_1V_2, V_1V_3$ are interior paths in region $A_4T_1T_2B_4$,
  - paths $W_1W_2, W_1W_3$ are interior paths in region $B_4R_1R_2C_4$,
  - all six paths are interior to $M$.

It suffices to prove the existence of the paths $U_1U_2$ and $U_1U_3$. We can use the following recursive algorithm to find them.

**Algorithm for searching paths $U_1U_2, U_1U_3$ in the region $A_4C_4S_2S_1$:**
init:  
\[ U_1 \leftarrow S_2, \ U_2 \leftarrow C_4, \ U_3 \leftarrow A_4, \]
\[ \text{path } U_1 U_2 = S_2 C_4, \]
\[ \text{done} = \text{FALSE} \]

while done = FALSE do {
  If there is an interior path \( \mathcal{P} \) for \( U_1, U_3 \) in \( A_4 C_4 S_2 S_1 \)
  \[ \text{path } U_1 U_3 \leftarrow \mathcal{P} \]
  \[ \text{done} = \text{TRUE} \]
  else
    if there is a blocking edge \( U_1' U_2' \) in the region \( A_4 C_4 S_2 S_1 \)
    such that \( U_1' \in S_1 S_2 \) and \( U_2' \in C_4 A_4 \) (Fig 2-7 (b))
    \[ U_1 \leftarrow U_1', \ U_2 \leftarrow U_2'. \]
    else /* \( U_1 = S_1 \) */
      \[ U_2 \leftarrow A_3, \ \text{path } U_1 U_3 = S_1 A_4, \ \text{done} = \text{TRUE} \]
  }

The algorithm will terminate because there are only a limited number of vertices on \( S_1 S_2 \).

Now, we are able to subdivide \( M \) along three pathes: \( U_1 U_2, U_2 C_4, C_4 W_3, W_3 W_1, W_1 W_2, W_2 B_4, B_4 V_3, V_3 V_1, V_1 V_2, V_2 A_4, A_4 U_3, U_3 U_1 \) using the quaternary subdivision (Fig. 2-7 (a)).

- **Case 2.2: \( A_4 = B_4 = C_4 \).**

  We find an edge \( \overline{U_{11} U_{12}} \) and a path \( A_4 U_{11} \) such that \( U_{11} \in A_1 T_1, U_{12} \in A_3 A_4, \) and \( \overline{U_{11} U_{12}}, A_4 U_{11} \) are interior to \( M \) as follows.

  - If \( A_1 = T_1 \), we simply select \( U_{11} = A_1, U_{12} = A_3, \) path \( A_4 U_{11} = \) path \( A_4 T_1 \),
    and \( \overline{U_{11} U_{12}} = A_1 A_3 \). Obviously, these two paths are interior to \( M \).
- If $A_1 \neq T_1$ and there is an interior path $\mathcal{P}$ inside the region $A_1T_1A_4 A_3$ (Fig. 2-8(a)), we select $U_{11} = A_1, U_{12} = A_3, \overline{U_{11}U_{12}} = \overline{A_1A_3}$, and path $A_4U_{11} = \mathcal{P}$.

- Suppose $A_1 \neq T_1$ and $A_1, A_4$ are not interior-connected in the region $A_1T_1A_4 A_3$. Based on Lemma 2, there is a blocking edge $u_1u_2$ for $A_1, A_4$. Because $\overline{A_1A_3}, \overline{A_4T_1}$ are edges, the vertices $u_1$ and $u_2$ must be on $A_1T_1$ and $A_3A_4$, respectively (Fig. 2-8 (b)). There might be more than one such blocking edge. We assume $u_1u_2$ is the one closest to $A_4$, i.e., there is a path $\mathcal{P}$ connecting $A_4$ and $u_1$, and $\mathcal{P}$ is an interior path in $M$. Now we assign $U_{11} = u_1, U_{12} = u_2, \overline{U_{11}U_{12}} = \overline{u_1u_2}$, and path $A_4U_{11} = \mathcal{P}$.

Similarly, we can get an edge $\overline{U_{21}U_{22}}$ and a path $A_4U_{21}$ such that $U_{21} \in A_2S_1, U_{22} \in A_3A_4$, and $\overline{U_{21}U_{22}}, A_4U_{21}$ are interior to $M$ (Fig. 2-8 (c)). Together, we have constructed a sub-triangular patch $A_4U_{11}U_{2}$ (we rename $U_1 = U_{11}, U_2 = U_{21}$, and the triangular patch is shown by the shadowed area in Fig. 2-8 (c)) inside the region $AT_1A_4S_1$. The patch $A_4U_{11}U_{2}$ intersects $\partial M$ only at $U_1, U_2$, and its three boundary edges are all interior to $M$.

In the same way, we can obtain two other sub-triangular patches $A_4V_1V_2$ inside $BR_1A_4T_2$, and $A_4W_1W_2$ inside $CS_2A_4R_2$ (Fig. 2-9 (a)). All three sub-triangular patches $A_4U_{11}U_{2}, A_4V_1V_2, A_4W_1W_2$ are disjoint, because they are separated by six edges $\overline{A_4R_1}, \overline{A_4R_2}, \overline{A_4S_1}, \overline{A_4S_2}, \overline{A_4T_1}, \overline{A_4T_2}$.
If \( U_1 \neq V_1, V_2 \neq W_2, W_1 \neq U_2 \) (Fig. 2-9 (a)), we can subdivide \( M \) along the boundaries of these three sub-triangular patches using mitsubishi subdivision.

If one of the following happens: \( U_1 = V_1, V_2 = W_2, W_1 = U_2 \), and, say, \( U_1 = V_1 \) (Fig. 2-9 (b)), we can subdivide \( M \) along paths \( U_1U_2, V_1V_2, U_2A_4 \) and \( V_2A_4 \) using quaternary subdivision.

\[ \text{Fig. 2-9} \]

- **Case 2.3**: Exactly one of the following is true: \( A_4 = B_4, A_4 = C_4, C_4 = B_4 \).

Assume \( A_4 = C_4, A_4 \neq B_4 \) (Fig. 2-10 (a)). Using the exact same methods as in Case 2.1 and Case 2.2, we can get sub-triangular patches \( A_4U_1U_2, A_4X_1X_2 \), and paths \( V_1V_2, V_1V_3, W_1W_2, W_1W_3 \) (Fig. 2-10 (b)). We obtain a third sub-triangular patch bounded by the following eight paths: \( A_4V_3, V_3V_1, V_1V_2, V_2B_4, B_4W_3, W_3W_1, W_1W_2, W_2A_4 \). Therefore, we get a result similar to Case 2.2. The rest of the argument is exactly same as in Case 2.2.
0.3.2 step 2

In step 1, we proved that any normal triangular patch $M$ can be subdivided into several sub-triangular patches. But these sub-triangular patches may be abnormal, so we cannot apply our general subdivision theorem on the sub-patches. Fortunately, we can normalize abnormal sub-patches in the way we described earlier.

Consider the result of using binary subdivision (Fig. 2-11). Assume we subdivide $M$ along the path $CD$ (in solid lines). There are three new boundaries $AD$, $DB$, $CD$ in the two sub-triangular patches $ACD$ and $BCD$. Obviously, paths $AD$ and $BD$ do not have any blocking edges because $M$ is normal. If path $CD$ has a blocking edge, say $UV$ (shown in a dotted line), we can subdivide $M$ using a new path formed by path $CU$, edge $UV$ and path $VD$. Thus, the blocking edge is removed in the new subdivision.

It is easy to see that this technique applies to all abnormal sub-triangular patches generated by the other three subdivisions.
Remark: From our arguments before, we know that the 'normality' of a triangle mesh is a weak condition: all closed 3D triangle meshes can be subdivided into normal triangular patches, and an abnormal triangular patch with an arbitrary number of illegal edges can always be decomposed into two normal triangular patches if we add one point.

In CAGD and computer graphics, it is common that a single triangular surface defined on a triangle domain, say, a triangular Bezier surface or B-spline surface, is subdivided into many triangular sub-patches by a 'divide and conquer' technique. No illegal edges are created in the subdivision procedure. Therefore

Corollary 3 If triangle mesh $M$ is generated by arbitrarily subdividing a singular triangular surface defined on a triangle domain, then $M$ can be represented as a subdivision tree where each subdivision is one of the four elementary types defined above.

Corollary 4 A necessary and sufficient condition for a triangle mesh to be represented by a single subdivision tree is that the mesh be normal.

0.4 Efficiency analysis

0.4.1 About entropy coding

Before we analyze the efficiency of our subdivision tree representation for topological coding, we need to briefly introduce some basic concepts from information theory.

Information theory uses entropy as a measure of how much information is contained in a message [10]. The higher the entropy of a message, the more information it contains. For a message consisting of symbols $S = \{s_1, ..., s_m\}$, each $s_i$ with probability $p_i$, the entropy of the information is defined as

$$H = \sum_{i=1}^{m} p_i \log_2 \frac{1}{p_i}$$

(0.1)

If we use $l_i$ bits to encode each symbol $s_i$ (i=1,..,m), the average codeword length is lower bounded by the entropy, i.e.,

$$l_{average} = \sum_{i=1}^{m} p_i l_i \geq \sum_{i=1}^{m} p_i \log_2 \frac{1}{p_i}$$
When $\log_{p_i} 1$ are integers, we can code each $s_i$ using $l_i = \log_{p_i} 1$ bits with a Huffman coding scheme. Otherwise, a more complicated scheme called arithmetic coding can be used so that the average coding length $l_{average}$ is as close to the entropy bound as possible [2]. Huffman and arithmetic coding are both often called entropy coding techniques.

When we analyze efficiency of a compression method theoretically, we can use the entropy of its output message without actually encoding the message.

It can be verified that the entropy defined by Eq. (0.1) has the maximum value of $\log_2 m$ when $p_i = \frac{1}{m}$ for all $i = 1, ..., m$. Therefore, we can always use an average $\log_2 m$ bits to encode each symbol in a message which consists of $m$ symbols $s_1, ..., s_m$.

### 0.4.2 Efficiency analysis

Our general subdivision theory tells us that the topological structure of a normal triangular patch can be represented (thus encoded) as a tree. Each node represents a triangle and each node has only four different ways to generate its children. Because the first elementary subdivision (binary subdivision) can involve adding a point to any of three edges, it has three cases. In addition, we have to include a 'leaf' case which has no subdivision. Thus we have seven subdivision cases to represent. Based on information theory, we can roughly estimate that each case can be coded by at most $\log_2 7 \approx 2.81$ bits.

When a normal triangular patch is represented by a subdivision tree, the representation is usually not unique. Fig 2-12 (b)(c) are two different subdivision tree representations for the same triangular patch shown in Fig 2-12(a). Excluding leaf nodes, the tree in Fig. 2-12(b) has 6 subdivisions while Fig 2-12(c) has 11 subdivisions for the total 14 triangles. That brings up the question: how can we get the most efficient subdivision?
We define BT rate as 'bits/triangle (BT)' to measure the coding efficiency. The BT rate for each of our four elementary subdivisions is $\frac{\log_7 2}{2} (\approx 1.40)$, $\frac{\log_7 3}{3} (\approx 0.94)$, $\frac{\log_7 4}{4} (\approx 0.70)$ and $\frac{\log_7 9}{9} (\approx 0.31)$, respectively. Obviously, the first elementary (binary) subdivision has the least efficiency, the last (mitsubishi) has the greatest.

The subdivision tree for a normal triangular patch with $n$ triangles has $n$ leaf nodes and at most $n-1$ non-leaf nodes. Therefore, if we use $\log_2 7$ bits for each subdivision case, the BT rate is upper bounded by $\frac{(2n-1) \log_7 2}{n} \leq 2 \log_2 7 \approx 5.62$ (BT). That happens when all non-leaf nodes use binary subdivisions. However, taking the probabilities into account, the BT rate of entropy coding is much smaller than the 5.62. Actually, we have the following theorem.

**Theorem 5** For a normal triangular patch with $n$ triangles ($n \gg 1$), the entropy of BT rate for any subdivision tree representation of the patch is lower bounded by $\frac{2}{8} \log_2 9 - 3 \approx 0.57$ and upper bounded by 2.52.

Proof: Let’s assume there are $n_i$ nodes using the $i$-th elementary subdivision ($i=1,2,3,4$) and the total number of triangles is $n$. Notice that the four elementary subdivisions increase the number of triangles by 1, 2, 3, and 8, respectively. Therefore, we have the following relations:

$$\sum_{i=1}^{4} k_i n_4 = n - 1 \approx n$$

(0.2)

where $k_1 = 1, k_2 = 2, k_3 = 3, k_4 = 8$.

With entropy coding, the BT rate is expressed as

$$E = \frac{n \log_2 \frac{n}{N} + \sum_{i=1}^{4} (n_i \log_2 \frac{n_i}{N})}{n}$$

(0.3)
where $N = n + \sum_{i=1}^{4} n_i$. Maximizing $E$ under the constraint Eq. (0.2), we immediately get the following equations using Lagrange multipliers:

$$\frac{\log_2 \frac{n_i}{k_i}}{k_i} = \alpha = \text{constant}, \text{ for } i = 1, 2, 3, 4$$

Denoting $x = 2^\alpha$, we have

$$\sum_{i=1}^{4} n_i = \sum_{i=1}^{4} x^{k_i}(n + \sum_{i=1}^{4} n_i)$$

i.e.

$$\sum_{i=1}^{4} n_i = n \frac{\sum_{i=1}^{4} x^{k_i}}{1 - \sum_{i=1}^{4} x^{k_i}}$$

So,

$$n_i = x^{k_i}(n + \sum_{i=1}^{4} n_i) = n \frac{x^{k_i}}{1 - \sum_{i=1}^{4} x^{k_i}}$$

From Eq. (0.2), we get

$$\frac{\sum_{i=1}^{4} k_i x^{k_i}}{1 - \sum_{i=1}^{4} x^{k_i}} = \frac{n - 1}{n} \approx 1$$

i.e.,

$$\sum_{i=1}^{4} (k_i + 1)x^{k_i} = 1 \quad (0.4)$$

Eq. (0.4) has a unique positive solution $x \approx 0.30435$, and thus $\alpha = \log_2 x \approx -1.7162$.

Based on the values of $x$ and $\alpha$, we also have the following $\sum_{i=1}^{4} n_i \approx 0.7399n$, $n_1 = x(n + \sum_{i=1}^{4} n_i) \approx 0.5295n$, $n_2 = x^2(n + \sum_{i=1}^{4} n_i) \approx 0.1612n$, $n_3 = x^3(n + \sum_{i=1}^{4} n_i) \approx 0.0490n$, $n_4 = x^4(n + \sum_{i=1}^{4} n_i) \approx 0.000128n$.

The maximum value of the BT rate can be calculated from Eq. (0.3): $E_{\max} = 2.52$.

Because Eq. (0.4) has only one positive root, which leads to the maximum value of $E$, the minimum value of $E$ must be reached at the boundary of the region: $\{n_1 \geq 0, n_2 \geq 0, n_3 \geq 0, n_4 \geq 0, n_1 + 2*n_2 + 3*n_3 + 8*n_4 = n - 0\}$. It is easy to verify that $E$ gets its minimum value when $n_1 = n_2 = n_3 = 0, n_4 = n - 1 \approx n$, i.e., the minimum value is $E_{\min} \approx \frac{9}{8} \log_2 9 - 3 \approx 0.57$. 

$\square$
The proof above suggests us that a smaller BT rate can be obtained when fewer binary subdivisions are used and more quaternary and mitsubishi subdivisions are used.

Although topological information is only a small part of mesh representation, the efficiency of topological coding is still an important problem, especially for hardware supported rendering.

Assume there are \( n \) vertices in a big mesh. According to the Euler’s theory, there are roughly \( 2n \) triangles. The simplest way to code the topological information is to store each triangle separately by indexing its three vertices. That costs \( 3\log_2 n \) (BT). OpenGL allows triangles to be coded in strips. A triangle of \( k \) triangles can be coded by indexing only \( k+2 \) vertices. That reduces costs to \( \frac{k+2}{k} \log_2 n \) (BT). A generalized triangle strip method includes swap commands in a triangle strip [5]. Its BT rate is roughly \( \frac{1}{2} \log_2 n + c \) for \( c \approx 2 \) (BT). Deering developed the notion of generalized triangle mesh to reuse a limited number previously appearing vertices stored in a buffer [3]. The BT rate is about \( \frac{1}{16} \log_2 n + 4 \) (BT). PM in [6] uses about \( \frac{1}{2} \log_2 n + 2.5 \) (BT) for topological coding. And modified PM in [11] uses \( \frac{1}{2} \log_2 n + 3.5 \) (BT). Li’s method [9] has an average BT rate of \( \frac{1}{2} \log_2 n + 4.5 \) (BT).

We notice that all coding schemes above have a term \( \log_2 n \) in their BT rates. The reason is that they access triangles randomly. Random accessing is useful in rendering, especially in view-dependent rendering.

One big improvement made by Taubin and Rossignac is to arrange all vertices and triangles in a fixed way so that no indexing term \( \log_2 n \) appears [13]. An average cost of 2 (BT) is reported in their testing. More importantly, when all vertices are arranged in a fixed order, their geometric information can be coded efficiently using a bit-plane coding scheme. However, the method cannot generate hierarchical topology, and is not good for rendering.

Our method uses subdivision trees to put all vertices in a fixed positions. Once a tree is constructed, the topological information coded by the tree is complete. We need not spend \( \log_2 n \) bits for indexing a vertex. The BT rate in our scheme is upper bounded by 2.52 (BT) for a normal triangular patch. The BT rate is lower than all existing methods, except [13]. Taubin et al. reported 2 (BT) in their tests, but didn’t give theoretical analysis of their coding efficiency.

The efficiency of our method will be reduced when we cut a big mesh into many
normal triangular patches. We mentioned earlier that we can always cut a mesh with genus $g$ and $k$ boundaries into $2n + k$ normal triangular patches with a penalty of adding $n_k$ triangles and $k$ vertices, where $n_k$ is the total number of vertices in $k$ boundaries. Thus, our efficiency reduces at most by $\frac{2n_k \log_2 \tau}{2n} = \frac{n_k \log_2 \tau}{n}(BT)$. We should point out that for most cases, $n_k \ll n$. In addition, the actual cost may much less than the upper bound we estimated here. Therefore, our coding efficiency is believed to be very high.

The efficiency of topological coding is not the main goal of our compression. Most storage is spent on geometry coding, not topology coding. We use subdivision tree representation because it helps us to construct wavelets on an arbitrary mesh, and wavelets are expected to have excellent performance in geometry coding.
Bibliography


