A Demand Response Calculus with Perfect Batteries

Dan-Cristian Tomozei

Joint work with Jean-Yves Le Boudec

CCW, Sedona AZ, 07/11/2012



Demand Response by Quantity

- distribution network operator may interrupt / modulate power
- elastic loads support graceful degradation
- Thermal load (Voltalis), water heaters (Romande Energie «commande centralisée»),

Tableau électrique



Lisison directe ou CPL Voltalis Bluepod switches off thermal load for 60 mn

Compteur



e-cars

Network Calculus? Service curve?

Voltalis:

At most 30 mn of interruption total per day

"Service curve" contract *Guaranteed energy delivered in* $(s,t) \ge \beta(t-s), \quad \forall 0 \le s \le t$ $\beta(t) =$ superadditive function.

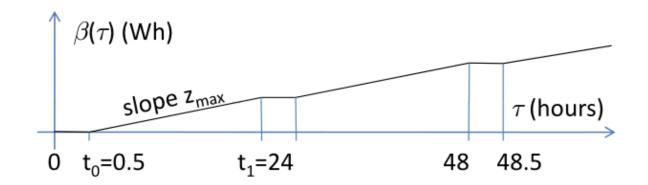


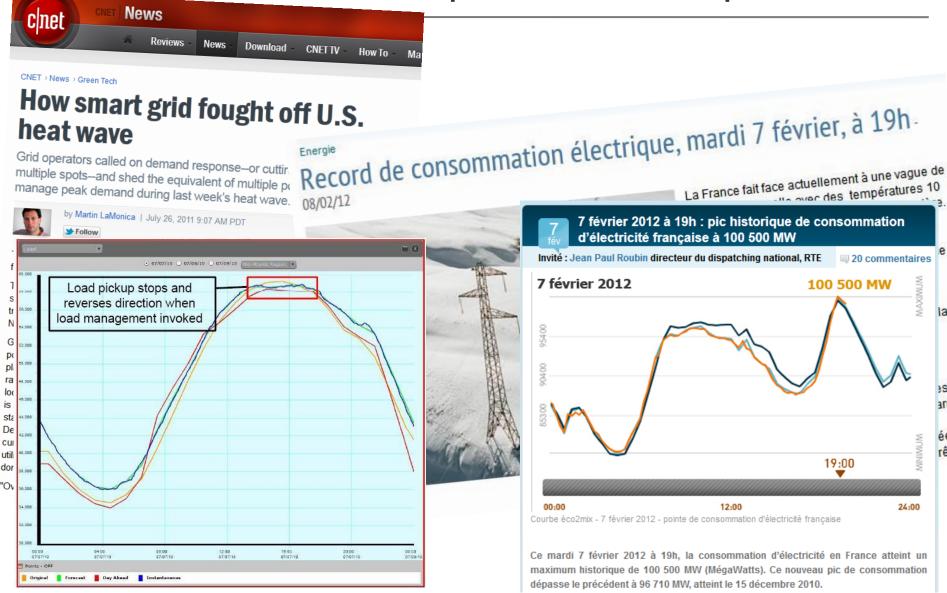


Tableau électriqu

BluePod



Real Situation: Unexpected Consumption Peaks





#4

Is Demand Response a good solution?

Today

Aggregate demand is predictable



- Operators foresee "reserve" (primary, secondary, tertiary of the second seco
 - E.g., gas turbines
- Reserve is expensive (capacity) / rare event \rightarrow demand response
 - Delay ("buffer") demand until the peak has passed ~ virtual energy storage

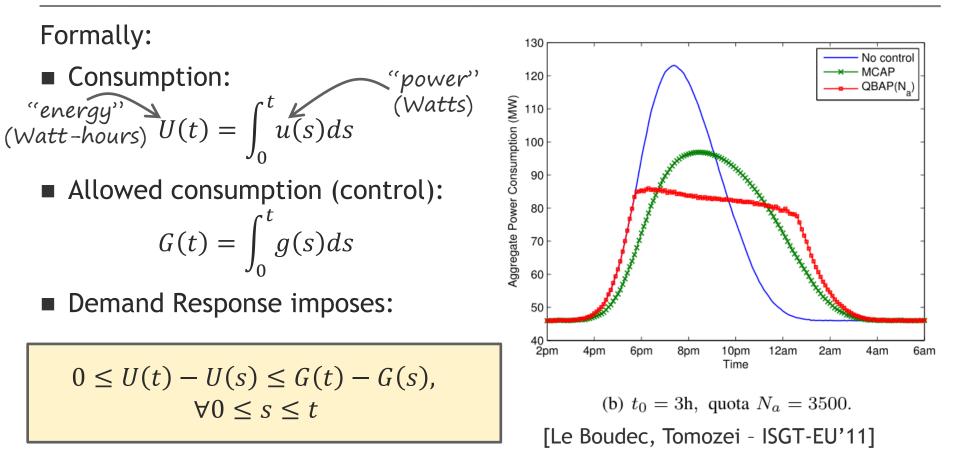
Tomorrow?

- High penetration of renewables → Large (unaffordable) reserve requirements
- E.g., fleet of e-cars \rightarrow DR exploits load flexibility





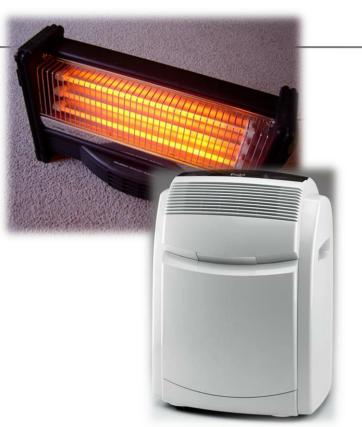
Demand Response by Quantity



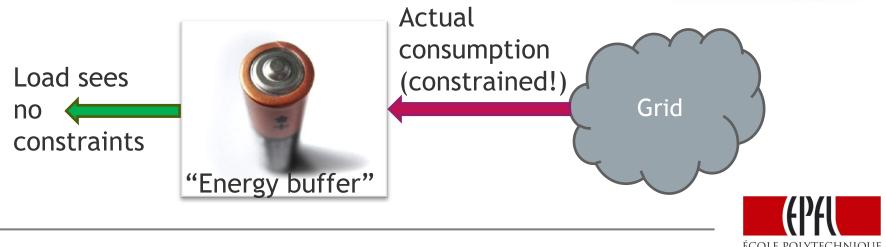


Inelastic load = lights out?

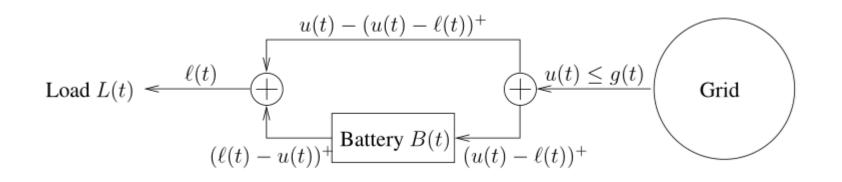
- Inelastic (non-dispatchable) loads
 - Lamps, TVs, Microwaves, ...
- Elastic (dispatchable) loads
 - Heating, A/C (TCLs)
- Make it dispatchable!
 - Inelastic load $L(t) = \int_0^t \ell(s) ds$
 - Use a large enough battery!



FÉDÉRALE DE LAUSANNE



The Perfect Battery

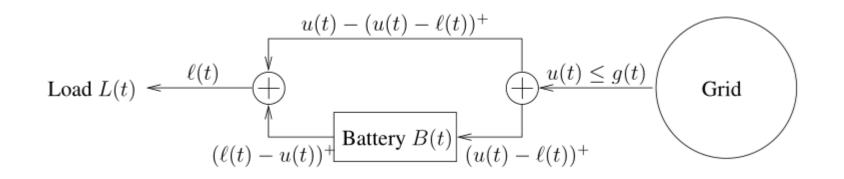


Battery may be charged $(u(t) > \ell(t))$ or discharged $(u(t) < \ell(t))$ Load $\ell(t)$ is given

Problem is to determine a power schedule u(t), subject to $0 \le u(t) \le g(t)$ and within battery constraints



System Equations for the Perfect Battery

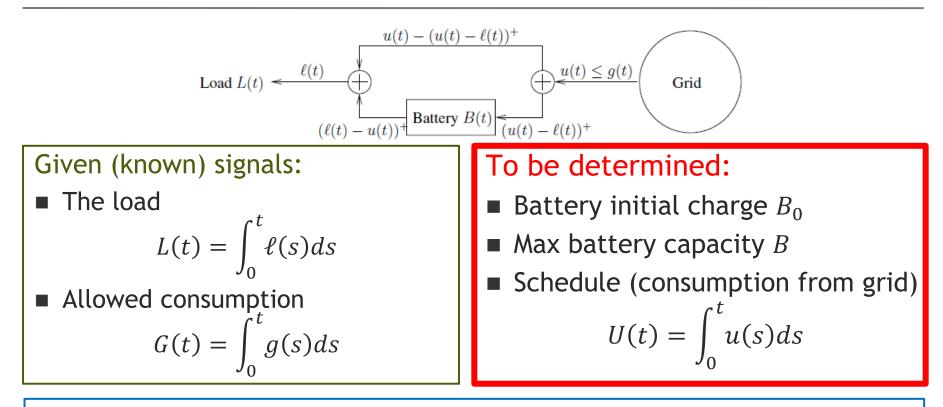


1. $L(t) \le B_0 + U(t)$ no underflow 2. $U(t) - L(t) + B_0 \le B$ no overflow 3. $U(t) - U(s) \le G(t) - G(s), \forall s \le t$ power constraint

where U(t), L(t), G(t) are cumulative functions such as $U(t) = \int_0^t u(s) ds$



Omniscient Problem



Constraints:

- Demand Response: $0 \le U(t) U(s) \le G(t) G(s), \forall s \le t$
- Perfect battery constraints: $L(t) \le B_0 + U(t)$ $U(t) - L(t) + B_0 \le B$



#10

Main Result

Theorem

• There exists a feasible schedule if and only if

$$\begin{cases} B_0 \ge \sup_t \left(L(t) - G(t) \right) \\ B \ge \sup_{0 \le s \le t} \left(L(t) - L(s) - G(t) + G(s) \right) \end{cases}$$

Moreover, if this is the case, then there exist a "minimal" and a "maximal" schedule:

$$U_*(t) = 0 \lor \sup_{\tau \ge t} (G(t) - G(\tau) + L(\tau) - B_0)$$

$$U^{*}(t) = G(t) \wedge \inf_{s \le t} (G(t) - G(s) + L(s) + B - B_{0})$$

$$U_*(t) \le U(t) \le U^*(t), \quad \forall t \ge 0$$

 The maximal schedule is causal & corresponds to the greedy policy (maximizes battery charge)



Service Curve Approach to Demand Response

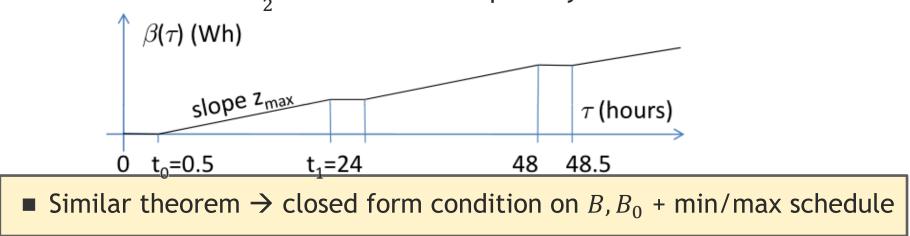
Assume we do not know the control signal G(t)

Instead: service curve contract [Le Boudec, Tomozei, ISGT-EU'11] $G(t) - G(s) \ge \beta(t - s), \quad \forall 0 \le s \le t$

 $\beta(t)$ = superadditive function.

Example:

- At most 30 mn of interruption total per day
- Or reduction to $\frac{z_{max}}{2}$ for 60mn total per day





Service Curve + Arrival Curve

Assume we don't know the load L(t) either!

Instead, L(t) is constrained by a subadditive arrival curve: $L(t) - L(s) \le \alpha(t - s), \quad \forall 0 \le s \le t$

Smallest arrival curve - obtained via min-plus deconvolution: $\alpha(t) \coloneqq \sup_{s \ge 0} \{L(s+t) - L(s)\}$

G(t) is well behaved (according to superadditive service curve): $G(t) - G(s) \ge \beta(t - s), \quad \forall 0 \le s \le t$

Theorem

For all $(B \ge)B_0 \ge B^* \coloneqq \sup_s \{\alpha(s) - \beta(s)\}$, there exists a feasible online (causal) schedule, valid for all loads and control signal compatible with $\alpha(\cdot)$ and $\beta(\cdot)$ respectively.

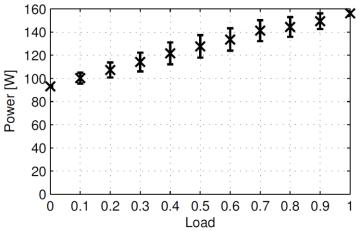


Application: Transparent DR for data centers

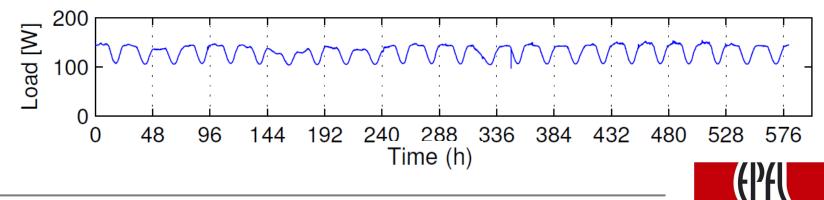
Akamai data set [Qureshi et al, SIGCOMM 2009]

#1

- Traffic at Akamai (millions of hits over 24 days)
- Measured power consumption of a desktop (SPEC)

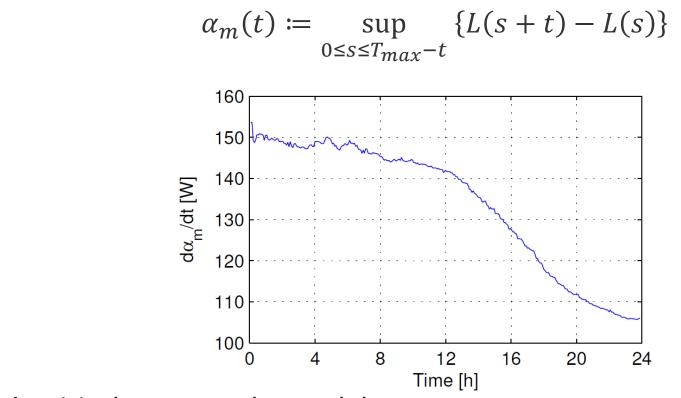


- Uniform repartition of tasks => consumption of one server



FÉDÉRALE DE LAUSANNE

Empirical arrival curve

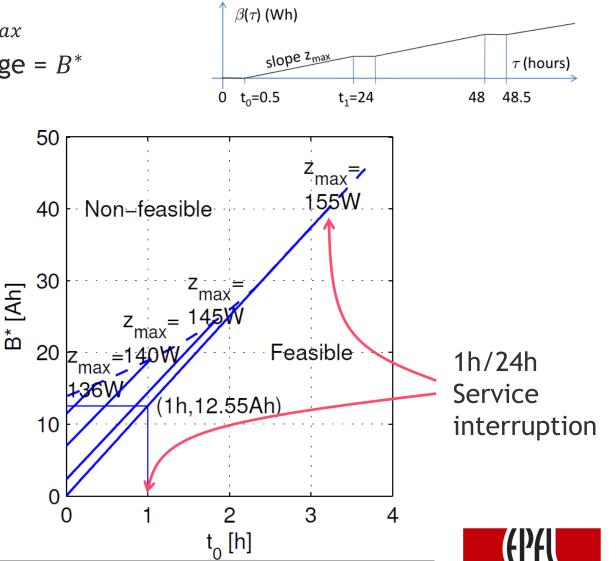


Intuitively = worst observed day



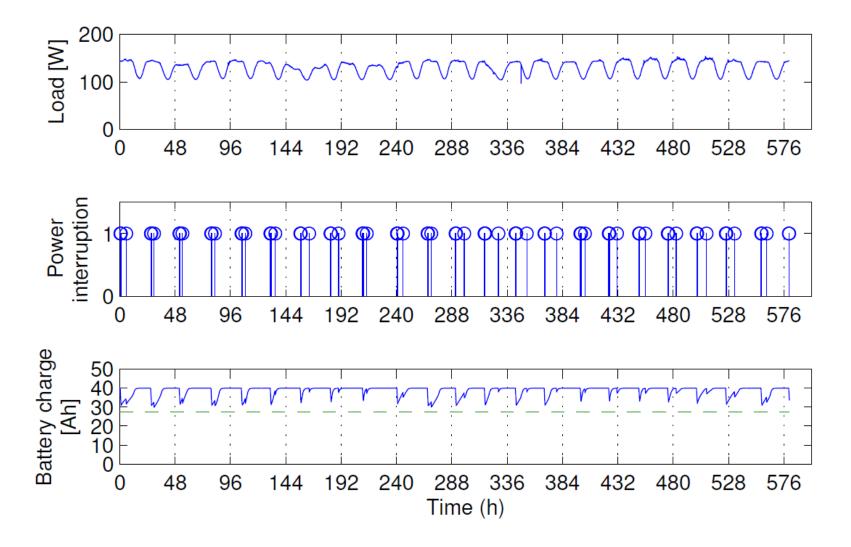
Choosing a SC contract and a battery

- Interruption time = t_0
- Maximum power = z_{max}
- Required battery charge = B*



ÉCOLE POLYTECHNIQUE Fédérale de Lausanne

A run of the system using the greedy policy





Conclusion

Another application of Network Calculus: Smart Grids

Theoretical results for perfect battery

Practical battery sizing problem

Easy to compute

Ongoing work

Realistic battery model (losses, aging, ...)

