

CSc 372

Comparative Programming Languages

31 : *Icon — Data Structures*

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Records

- Icon has built-in support for records, lists, tables, and sets.
- These data structures can be freely combined, so that it is easy to construct a list of tables of sets,

Records

- Records and procedures are the only declarations in Icon. They must be declared at the outermost (global) level:

```
record name(field1,field2,...)
```
- You don't give the types of the fields, just their names.
- `type(x)`, where x is a record variable, will return the name (a string) of the record type.
- If R is a record variable, `R.field1` references the field whose name is `field1`.

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Complex Arithmetic Module

```
record complex(re, im)

procedure add(a,b)
    return complex(a.re+b.re, a.im+b.im)
end

procedure main ()
    local x, r, i
    x := complex(5, 4)
    y := complex(1,2)
    z := add(x,y)

    r := z.re      # or r := z[1]
    i := z.im      # or r := z[2]
    t := type(z)  # t="complex"
end
```

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Lists

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Lists

- Lists are a built-in Icon datatype. Lists can be accessed from the beginning (the way you would in LISP, Prolog, etc), the end, or indexed (the way you would access an array in Pascal).
- Lists can be **heterogeneous**, they can contain elements of different type.

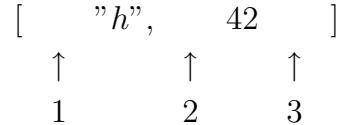
x := ["hello", 1, 3.14, "x", "y"] A list of a string, an integer, a float, and two strings.

y := list(5, "hej") A list of five strings:
["hej", ..., "hej"].

x[2:4] The list consisting of the second, third, and fourth element of x.

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Lists vs. Strings

- Lists are indexed in the same way as strings:

- Strings are **immutable**. This means that when you assign to an element of a string you actually get a **new** string as result.
- Lists are **mutable**. That is, when you assign to an element of a list, the list actually changes.

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List Operations

`s := list()` Create an empty list.
`s := list(n)` Create a list of `n` nulls.
`s := list(n,v)` Create a list of `n` vs.
`s := *x` Number of elements of `x`.
`x ||| y` Concatenate `x` and `y`.
!x Generate all elements of the list, in order, as in `every`
`X := !L do write(X).`

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List Operations...

`x ||| y` Concatenate `x` and `y`.
`put(x, 67)` Add 67 to the end of the list `x`.
`get(x)` Remove and return the last element of `x`.
`push(x, 1024)` Add a new element to the beginning of `x`.
`pop(x)` Remove and return the first element of `x`.
!x Generate all elements of the list, in order, as in `every`
`X := !L do write(X).`
?x Return a random element from list.
`x==y` Succeed if `x` and `y` are the same string.
`x~==y` Succeed if `x` and `y` are different strings.

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Examples

```
][[ L := list(5,10);
    r1 := L1:[10,10,10,10,10] (list)
][[ L[2] := 42;
][[ L;
    r3 := L1:[10,42,10,10,10]
][[ L := [1,2,3,4,5];
][[ L[1:3];
    r5 := L1:[1,2]
][[ L[0:-3];
    r6 := L1:[3,4,5]
][[ every i := !L do write(i);
1
2
3
4
5
```

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Examples

```
][[ L := [[1,2],[3,4],[5,6]];
    r1 := L1:[L2:[1,2],L3:[42,4],L4:[5,6]]
][[ x := pop(L);
][[ x;
    r5 := L1:[1,2] (list)
][[ L;
    r6 := L1:[L2:[42,4],L3:[5,6]]
][[ L := [1,2,3,4,5];
][[ every !L :=: ?L;
][[ L;
    r9 := L1:[2,1,5,3,4]
```

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Fibonacci

```
procedure main()
n := 20
f := [1,1]
repeat {
    i := get(f)
    if i>n then break
    write(i)
    put(f,i+f[1])
}
end
```

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Tables

```
procedure main()
n := 100
p := list(n,1)
every i := 2 to sqrt(n) do
    if p[i]=1 then
        every j := i+i to n by i do
            p[j] := 0
every i := 2 to n do
    if p[i]=1 then
        write(i)
end
```

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Tables

- Tables are **associative arrays**, they map keys to values.
Both values and keys can be of arbitrary type.

Table Operations

Tables are **associative arrays**, they map keys to values.
Both values and keys can be of arbitrary type.

`x:=table(0)` Create a new table `x` whose **default value** is 0. This means that if you look up a key which has no corresponding value, 0 is returned.

`*x` Number of elements in the table.

`?x` An arbitrary element from the table.

`keys(x)` Generate all keys in `x`, one at a time.

`!x` Generate all values, one at a time.

```
every x := keys(T) do  
    write(x, " ==> ", T[x])
```

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Sets

```
x[ "monkey" ] := "banana"  
x[ 3.14 ] := "pi"  
x[ "pi" ] := 3.14  
x[ "pi" ] += 1 Increment pi by 1  
r := x[ "coconut" ] r will be 0  
member(x, 3.14) returns "pi"  
member(x, "banana") fails  
insert(x, "banana", 5) x["banana"] := 5  
delete(x, "monkey") remove "monkey"  
every m := key(x) do write(m) write keys  
every m := !x do write(m) write values
```

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Sets

- Sets are unordered collections of elements.
- `set()` creates an empty set.
- `set(L)` creates a set from a list of elements.
- All the standard set-operations (intersection, etc.) are built-in.

Set Operations

```
x := set([5, 3, "monkey"]) Create a 3-element set  
from a list.  
  
member(x, 5) returns 5  
  
member(x, "banana") fails  
  
insert(x, "banana") add "banana" to x  
  
delete(x, 5) returns the set {3, "banana",  
    "monkey"}  
  
*x number of elements (3)  
?x random element from x  
!x generate the elements
```

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Prime Sieve

```
procedure main()  
n := 100  
p := set()  
every i:=2 to n do insert(p,i)  
every i := 2 to sqrt(n) &  
    member(p,i) &  
    j := i+i to n by i do  
        delete(p,j)  
every i := 2 to n & member(p,i) do  
    write(i)  
end
```

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Set Operations

S := S1 op S2 set union ($op=++$), intersection ($op=**$),
difference ($op=--$).

```
while insert(S, read(f)) Read elements from file f  
into set S
```

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Binary Trees

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```
link ximage
record node (item, left, right)
procedure Preorder (T)
    if \T then {
        write(T.item);
        Preorder(T.left); Preorder(T.right)}
    end

procedure main()
    t := node(1, node(2, &null, &null),
              node(3, &null,
                    node(4, &null, &null)))
    Preorder(t); xdump(t)
end
```

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Readings and References

- Read Christopher, pp 29--34, 105--126.

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```
> icont b
> b
1
2
3
4
R_node_4 := node()
R_node_4.item := 1
R_node_4.left := R_node_1 := node()
R_node_1.item := 2
R_node_4.right := R_node_3 := node()
R_node_3.item := 3
R_node_3.right := R_node_2 := node()
R_node_2.item := 4
```

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Acknowledgments

- Some material on these slides has been modified from Thomas W Christopher's Icon Programming Language Handbook,

<http://www.tools-of-computing.com/tc/CS/iconprog.pdf>.

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