

# CSc 372

## Comparative Programming Languages

### 27: Prolog — Grammars

Christian Collberg

[collberg+372@gmail.com](mailto:collberg+372@gmail.com)

Department of Computer Science  
University of Arizona

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# Prolog Grammar Rules

- A DCG (**definite clause grammar**) is a phrase structure grammar annotated by Prolog variables.
- DCGs are translated by the Prolog interpreter into normal Prolog clauses.
- Prolog DCG:s can be used for generation as well as parsing. I.e. we can run the program backwards to generate sentences from the grammar.

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## Prolog Grammar Rules...

```
s      --> np, vp.
vp     --> v, np.
vp     --> v.
np     --> n.
n      --> [john].    n      --> [lisa].
n      --> [house].
v      --> [died].    v      --> [kissed].
```

```
?- s([john, kissed, lisa], []).
```

yes

```
?- s([lisa, died], []).
```

yes

```
?- s([kissed, john, lisa], []).
```

no

## Prolog Grammar Rules...

```
?- s(A, []).
A = [john,died,john] ;
A = [john,died,lisa] ;
A = [john,died,house] ;
A = [john,kissed,john] ;
A = [john,kissed,lisa] ;
A = [john,kissed,house] ;
A = [john,died] ;
A = [john,kissed] ;
A = [lisa,died,john] ;
A = [lisa,died,lisa] ;
A = [lisa,died,house] ;
A = [lisa,kissed,house] ;
A = [lisa,died] ;
```

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## Implementing Prolog Grammar Rules

- Prolog turns each grammar rule into a clause with one argument.

- The rule  $S \rightarrow NP VP$  becomes

```
s(Z) :- np(X), vp(Y), append(X,Y,
```

- This states that Z is a sentence if X is a noun phrase, Y is a verb phrase, and Z is X followed by Y.

## Implementing Prolog Grammar Rules...

```
s(Z) :- np(X), vp(Y), append(X,Y,Z).  
np(Z) :- n(Z).  
vp(Z) :- v(X), np(Y), append(X,Y,Z).  
vp(Z) :- v(Z).  
n([john]). n([lisa]). n([house]).  
v([died]). v([kissed]).
```

```
?- s([john,kissed,lisa]).
```

```
yes
```

```
?- s(S).
```

```
S = [john,died,john] ;
```

```
S = [john,died,lisa] ; ...
```

## Implementing Prolog Grammar Rules...

- The append's are expensive — Prolog uses **difference lists** instead.

- The rule

```
s(A,B) :- np(A,C), vp(C,B).
```

says that there is a sentence at the beginning of A (with B left over) if there is a noun phrase at the beginning of A (with C left over), and there is a verb phrase at the beginning of C (with B left over).

## Implementing Prolog Grammar Rules...

```
s(A,B) :- np(A,C), vp(C,B).  
np(A,B) :- n(A,B).  
vp(A,B) :- v(A,C), np(C,B).  
vp(A,B) :- v(A,B).  
n([john|R],R). n([lisa|R],R).  
v([died|R],R). v([kissed|R],R).
```

```
?- s([john,kissed,lisa], []).
```

```
yes
```

```
?- s([john,kissed|R], []).
```

```
R = [john] ;
```

```
R = [lisa] ; ...
```

## Generating Parse Trees

- DCGs can build parse trees which can be used to construct a semantic interpretation of the sentence.
- The tree is built bottom-up, when Prolog returns from recursive calls. We give each phrase structure rule an extra argument which represents the node to be constructed.

## Generating Parse Trees...

```
s(s(NP,VP))      --> np(NP), vp(VP).
vp(vp(V, NP))   --> v(V), np(NP).
vp(vp(V))       --> v(V).
np(np(N))       --> n(N).
n(n(john))      --> [john].
n(n(lisa))      --> [lisa].
n(n(house))     --> [house].
v(n(died))      --> [died].
v(n(kissed))    --> [kissed].
```

## Generating Parse Trees...

- The rule

```
s(s(NP,VP)) --> np(NP), vp(VP).
```

says that the top-level node of the parse tree is an s with the sub-trees generated by the np and vp rules.

```
?- s(S, [john, kissed, lisa], []).
S=s(np(n(john)),vp(n(kissed),np(n(lisa))))
?- s(S, [lisa, died], []).
S=s(np(n(lisa)),vp(n(died)))
?- s(S, [john, died, lisa], []).
S=s(np(n(john)),vp(n(died),np(n(lisa))))
```

## Generating Parse Trees...

- We can of course run the rules backwards, turning parse trees into sentences:

```
?- s(s(np(n(john)),vp(n(kissed),
      np(n(lisa)))), S, []).
S=[john, kissed, lisa]
```

# Ambiguity

- An ambiguous sentence is one which can have more than one meaning.

## Lexical ambiguity:

### homographic

- spelled the same
- *bat* (wooden stick/animal)
- *import* (noun/verb)

### polysemous

- different but related meanings
- *neck* (part of body/part of bottle/narrow strip of land)

### homophonic

- sound the same

● to/too/two

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# Ambiguity...

## Syntactic ambiguity:

- More than one parse (tree).
- Many missiles have many war-heads.
  
- “Duck” can be either a verb or a noun.
- “her” can either be a determiner (as in “her book”), or a noun: “I liked her dancing”.

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# Ambiguity...

```
s(s(NP,VP)) --> np(NP), vp(VP).
vp(vp(V, NP)) --> v(V), np(NP).
vp(vp(V, S)) --> v(V), s(S).
vp(vp(V)) --> v(V).
np(np(Det,N)) --> det(Det), n(N).
np(np(N)) --> n(N).
n(n(i)) --> [i].
n(n(duck)) --> [duck].
v(v(duck)) --> [duck].
v(v(saw)) --> [saw].  n(n(saw)) --> [saw].
n(n(her)) --> [her].
det(det(her)) --> [her].
```

```
?- s(S, [i, saw, her, duck], []).
```

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# Pascal Declarations

```
?- decl([const, a, =, 5, ;,
        var, x, :, 'INTEGER', ;], []).
    yes
?- decl([const, a, =, a, ;, var, x,
        :, 'INTEGER', ;], []).
    no

decl --> const_decl, type_decl,
        var_decl, proc_decl.
```

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## Pascal Declarations

```
% Constant declarations
const_decl --> [ ].
const_decl -->
    [const], const_def, [;], const_defs.

const_defs --> [ ].
const_defs --> const_def, [;], const_defs.
const_def --> identifier, [=], constant.

identifier --> [X], {atom(X)}.
constant --> [X], {(integer(X); float(X))}.
```

## Pascal Declarations...

```
% Type declarations
type_decl --> [ ].
type_decl --> [type], type_def, [;], type_defs.

type_defs --> [ ].
type_defs --> type_def, [;], type_defs.
type_def --> identifier, [=], type.

type --> ['INTEGER']. type --> ['REAL'].
type --> ['BOOLEAN']. type --> ['CHAR'].
```

## Pascal Declarations...

```
% Variable declarations
var_decl --> [ ].
var_decl --> [var], var_def, [;], var_defs.

var_defs --> [ ].
var_defs --> var_def, [;], var_defs.
var_def --> id_list, [:], type.

id_list --> identifier.
id_list --> identifier, [','], id_list.
```

## Pascal Declarations...

```
% Procedure declarations
proc_decl --> [ ].
proc_decl --> proc_heading, [;], block.
proc_heading --> [procedure], identifier,
    formal_param_part.
formal_param_part --> [ ].
formal_param_part --> ['(',
    formal_param_section, [')']].
formal_param_section --> formal_params.
formal_param_section --> formal_params, [;],
    formal_param_section.
formal_params --> value_params.
formal_params --> variable_params.
value_params --> var_def.
variable_params --> [var], var_def.
```

# Pascal Declarations – Building Trees

```
decl(decl(C, T, V, P)) -->  
  const_decl(C), type_decl(T),  
  var_decl(V), proc_declaration(P).
```

```
const_decl(const(null)) --> [ ].  
const_decl(const(D, Ds)) -->  
  [const], const_def(D), [;], const_defs(Ds).
```

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# Pascal Declarations – Building Trees...

```
const_defs(null) --> [ ].  
const_defs(const(D, Ds)) -->  
  const_def(D), [;], const_defs(Ds).
```

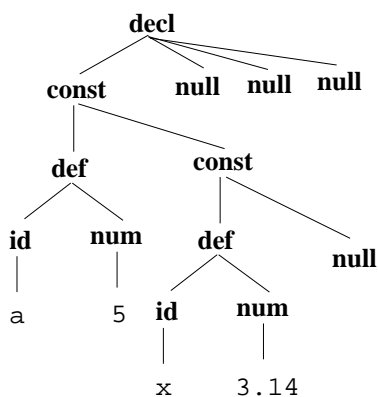
```
const_def(def(I, C)) --> ident(I), [=], const(C).
```

```
ident(id(X)) --> [X], {atom(X)}.  
const(num(X)) --> [X], {(integer(X); float(X))}.
```

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# Pascal Declarations – Example Parse



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# Pascal Declarations – Example Parse...

```
?- decl(S, [const, a, =, 5, ;, x, =, 3.14, ;], [])
```

```
S = decl(  
  const(def(id(a), num(5)),  
    const(def(id(x), num(3.14)),  
      null)),  
  null, null, null)
```

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## Number Conversion

```
- number(V, [sixty, three], []).
V = 63
- number(V, [one, hundred, and, fourteen], []).
V = 114
- number(V, [nine, hundred, and, ninety, nine], []).
V = 999
- number(V, [fifty, ten], []).
no
```

## Number Conversion..

```
number(0) --> [zero].
number(N) --> xxx(N).

xxx(N) --> digit(D), [hundred], rest_xxx(N1),
           {N is D * 100+N1}.
xxx(N) --> xx(N).

rest_xxx(0) --> [ ]. rest_xxx(N) --> [and], xx(N).

xx(N) --> digit(N).
xx(N) --> teen(N).
xx(N) --> tens(T), rest_xx(N1), {N is T+N1}.

rest_xx(0) --> [ ]. rest_xx(N) --> digit(N).
```

## Number Conversion..

```
digit(1) --> [one].      teen(10) --> [ten].
digit(2) --> [two].      teen(11) --> [eleven].
digit(3) --> [three].    teen(12) --> [twelve].
digit(4) --> [four].     teen(13) --> [thirteen].
digit(5) --> [five].     teen(14) --> [fourteen].
digit(6) --> [six].      teen(15) --> [fifteen].
digit(7) --> [seven].    teen(16) --> [sixteen].
digit(8) --> [eight].    teen(17) --> [seventeen].
digit(9) --> [nine].     teen(18) --> [eighteen].
                       teen(19) --> [nineteen].
tens(20) --> [twenty].   tens(30) --> [thirty].
tens(40) --> [forty].    tens(50) --> [fifty].
tens(60) --> [sixty].    tens(70) --> [seventy].
tens(80) --> [eighty].  tens(90) --> [ninety].
```

## Expression Evaluation

- Evaluate infix arithmetic expressions, given as character strings.

```
?- expr(X, "234+345*456", []).
X = 157554
```

```
expr(Z) --> term(X), "+", expr(Y), {Z is X + Y}.
expr(Z) --> term(X), "-", expr(Y), {Z is X - Y}.
expr(Z) --> term(Z).

term(Z) --> num(X), "*", term(Y), {Z is X * Y}.
term(Z) --> num(X), "/", term(Y), {Z is X / Y }.
term(Z) --> num(Z).
```

## Expression Evaluation...

Prolog grammar rules are equivalent to recursive descent parsing. Beware of left recursion!

Anything within curly brackets is “normal” Prolog code.

```
n(C) --> "+", num(C).
n(C) --> "-", num(X), {C is -X}.
n(X) --> int(0, X).

t(L, V) --> digit(C), {V is L * 10 + C}.
t(L, X) --> digit(C), {V is L* 10 + C},
    int(V, X).

digit(X) --> [C], {"0" =< C, C =< "9", X is C-"0"}.
```

## Summary

- Read Clocksin & Mellish, Chapter 9.
- Grammar rule syntax:
  - A grammar rule is written `LHS --> RHS`. The left-hand side (LHS) must be a non-terminal symbol, the right-hand side (RHS) can be a combination of terminals, non-terminals, and Prolog goals.
  - Terminal symbols (words) are in square brackets: `n --> [house]`.
  - More than one terminal can be matched by one rule: `np --> [the,house]`.

## Summary...

- Grammar rule syntax (cont):
  - Non-terminals (syntactic categories) can be given extra arguments: `s(s(N,V)) --> np(N),vp(V) ..`
  - Normal Prolog goals can be embedded within grammar rules: `int(C) --> [C],{integer(C)}`.
  - Terminals, non-terminals, and Prolog goals can be mixed in the right-hand side: `x --> [y], z, {w}, [r], p.`
- Beware of left recursion! `expr --> expr ['+']`  
`expr` will recurse infinitely. Rules like this will have to be rewritten to use right recursion.

## Homework

- Write a program which uses Prolog Grammar Rules to convert between English time expressions and a 24-hour clock (“Military Time”).
- You may assume that the following definitions are available:

```
digit(1) --> [one]. ....
digit(9) --> [nine].
teen(10) --> [ten]. ....
teen(19) --> [nineteen].
tens(20) --> [twenty]. ....
tens(90) --> [ninety].

?- time(T, [eight, am], []).
    T = 8:0 % Or, better, 8:00
```



## Homework...

```
?- time(T, [eight, thirty, am], []).  
   T = 8:30  
?- time(T, [eight, fifteen, am], []).  
   T = 8:15  
?- time(T, [eight, five, am], []).  
   no  
?- time(T, [eight, oh, five, am], []).  
   T = 8:5 % Or, better, 8:05  
?- time(T, [eight, oh, eleven, am], []).  
   no  
?- time(T, [eleven, thirty, am], []).  
   T = 11:30  
?- time(T, [twelve, thirty, am], []).  
   T = 0:30 % !!!
```

## Homework...

```
?- time(T, [eleven, thirty, pm], []).  
   T = 23:30  
?- time(T, [twelve, thirty, pm], []).  
   T = 12:30 % !!!  
?- time(T, [ten, minutes, to, four, am], []).  
   T = 3:50  
?- time(T, [ten, minutes, past, four, am], []).  
   T = 4:10  
?- time(T, [quarter, to, four, pm], []).  
   T = 15:45  
?- time(T, [quarter, past, four, pm], []).  
   T = 16:15  
?- time(T, [half, past, four, pm], []).  
   T = 16:30
```