#### Sometimes it is more natural to use an infix notation for **CSc 372** a function application, rather than the normal prefix one: ● 5 + 6 (infix) **Comparative Programming** ● (+) 5 6 (prefix) Languages Haskell predeclares some infix operators in the standard prelude, such as those for arithmetic. 9: Haskell — Curried Functions For each operator we need to specify its precedence and associativity. The higher precedence of an Christian Collberg operator, the stronger it binds (attracts) its arguments: collberg+372@gmail.com hence: Department of Computer Science $3 + 5*4 \equiv 3 + (5*4)$ University of Arizona $3 + 5*4 \neq (3 + 5) * 4$ Copyright © 2005 Christian Collberg [1] [2] -Fall 2005 - 9 372 - Fall 2005 - 9

### **Declaring Infix Functions...**

 The associativity of an operator describes how it binds when combined with operators of equal precedence. So, is

```
5-3+9 \equiv (5-3)+9 = 11
OR
5-3+9 \equiv 5-(3+9) = -7
```

The answer is that + and – associate to the left, i.e. parentheses are inserted from the left.

- Some operators are right associative:  $5^3^2 \equiv 5^3(3^2)$
- Some operators have free (or no) associativity. Combining operators with free associativity is an error:

[0]

```
5 == 4 < 3
```

 $\Rightarrow$  ERROR

# **Declaring Infix Functions...**

**Declaring Infix Functions** 

The syntax for declaring operators:

infixr prec oper	right assoc.
infixl prec oper	left assoc.
infix prec oper	free assoc.

### From the standard prelude:

infix1 7 \*
infix 7 /, `div`, `rem`, `mod`
infix 4 ==, /=, <, <=, >=, >

An infix function can be used in a prefix function application, by including it in parenthesis. Example:

```
? (+) 5 ((*) 6 4)
29
```

	Multi-Argument Functions	
Multi-Argument Functions	<ul> <li>Haskell only supports one-argument functions.</li> <li>An <i>n</i>-argument function f(a<sub>1</sub>,, a<sub>n</sub>) is constructed in either of two ways: <ol> <li>By making the one input argument to f a tuple holding the n arguments.</li> </ol> </li> <li>By letting f "consume" one argument at a time. This is called currying.</li> </ul>	
	TupleCurryingadd :: (Int,Int)->Intadd :: Int->Int->Intadd (a, b) = a + badd a b = a + b	
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Currying	Currying	
<ul> <li>Currying is the preferred way of constructing multi-argument functions.</li> <li>The main advantage of currying is that it allows us to define specialized versions of an existing function.</li> <li>A function is specialized by supplying values for one or more (but not all) of its arguments.</li> <li>Let's look at Haskell's plus operator (+). It has the type (+) :: Int -&gt; (Int -&gt; Int).</li> <li>If we give two arguments to (+) it will return an Int: (+) 5 3 ⇒ 8</li> </ul>	<ul> <li>If we just give one argument (5) to (+) it will instead return a function which "adds 5 to things". The type of this specialized version of (+) is Int -&gt; Int.</li> <li>Internally, Haskell constructs an intermediate - specialized - function: add5 :: Int -&gt; Int add5 a = 5 + a</li> <li>Hence, (+) 5 3 is evaluated in two steps. First (+) 5 is evaluated. It returns a function which adds 5 to its argument. We apply the second argument 3 to this new function, and the result 8 is returned.</li> </ul>	

<ul> <li>To summarize, Haskell only supports one-argument functions. Multi-argument functions are constructed by successive application of arguments, one at a time.</li> <li>Currying is named after logician Haskell B. Curry (1900-1982) who popularized it. It was invented by Schönfinkel in 1924. Schönfinkeling doesn't sound too good</li> <li>Note: Function application (f x) has higher precedence (10) than any other operator. Example: f 5 + 1 ⇔ (f 5) + 1 f 5 6 ⇔ (f 5) 6</li> </ul>	<ul> <li>Let's see what happens when we evaluate f 3 4 5, where f is a 3-argument function that returns the sum of its arguments.</li> <li>f :: Int -&gt; (Int -&gt; (Int -&gt; Int))</li> <li>f x y z = x + y + z</li> <li>f 3 4 5 = ((f 3) 4) 5</li> </ul>
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Currying Example	Currying Example
<ul> <li>(f 3) returns a function f' y z (f' is a specialization of f) that adds 3 to its next two arguments.</li> <li>f 3 4 5 ≡ ((f 3) 4) 5 ⇒ (f' 4) 5</li> <li>f' :: Int -&gt; (Int -&gt; Int)</li> <li>f' y z = 3 + y + z</li> </ul>	<ul> <li>(f' 4) (= (f 3) 4) returns a function f''z (f'' is a specialization of f') that adds (3+4) to its argument.</li> <li>f 3 4 5 ≡ ((f 3) 4) 5 ⇒ (f' 4) 5 ⇒ f'' 5</li> <li>f'' :: Int -&gt; Int f'' z = 3 + 4 + z</li> <li>Finally, we can apply f'' to the last argument (5) and get the result:</li> <li>f 3 4 5 ≡ ((f 3) 4) 5 ⇒ (f' 4) 5 ⇒ f'' 5 ⇒ 3+4+5 ⇒ 12</li> </ul>

### **Currying Example**

The Combinatorial Function: comb :: Int -> Int -> Int • The combinatorial function  $\binom{n}{r}$  "n choose r", computes comb n r = fact n/(fact r\*fact(n-r))the number of ways to pick r objects from n. comb 5 3  $\Rightarrow$  (comb 5) 3  $\Rightarrow$  $comb^5$  3  $\Rightarrow$ 120 / (fact 3 \* (fact 5-3))  $\Rightarrow$  $\left(\begin{array}{c}n\\r\end{array}\right) = \frac{n!}{r!*(n-r)!}$ 120 / (6 \* (fact 5-3))  $\Rightarrow$ 120 / (6 \* fact 2)  $\Rightarrow$ 120 / (6 \* 2)  $\Rightarrow$ In Haskell: 120 / 12  $\Rightarrow$ comb :: Int -> Int -> Int 10 comb n r = fact n/(fact r\*fact(n-r))  $comb^5 r = 120 / (fact r * fact(5-r))$ ? comb 5 3 • comb<sup>5</sup> is the result of partially applying comb to its first 10 argument. [14] [13] -Fall 2005 - 9 372 - Fall 2005 - 9 Associativity What's the Type, Mr. Wolf? If the type of a function f is Function application is left-associative:  $f a b = (f a) b f a b \neq f (a b)$  $t_1 \rightarrow t_2 \rightarrow \cdots \rightarrow t_n \rightarrow t$ and f is applied to arguments The function space symbol `->' is right-associative:  $e_1::t_1, e_2::t_2, \cdots, e_k::t_k,$  $a \rightarrow b \rightarrow c = a \rightarrow (b \rightarrow c)$  $a \rightarrow b \rightarrow c \neq (a \rightarrow b) \rightarrow c$ ● and k < n</p> f takes an Int as argument and returns a function of then the result type is given by cancelling the types type Int -> Int. g takes a function of type Int ->  $t_1 \cdots t_k$ : Int as argument and returns an Int:  $t_1 \rightarrow t_2 \rightarrow \cdots \rightarrow t_k \rightarrow t_{k+1} \rightarrow \cdots \rightarrow t_n \rightarrow t$ f' :: Int -> (Int -> Int) • Hence, f  $e_1 e_2 \cdots e_k$  returns an object of type  $t_{k+1} \rightarrow \cdots \rightarrow t_n \rightarrow t$ . f :: Int -> Int -> Int This is called the Rule of Cancellation. q :: (Int -> Int) -> Int

**Currying Example...** 

## Homework

٩	Define an operator \$\$ so that x \$\$ xs returns True if		
	${f x}$ is an element in ${f xs},$ and ${f False}$ otherwise.		

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?	4 \$\$ [1,2,5,6,4 True	Example:
?	4 \$\$ [1,2,3,5] False	
?	4 \$\$ [] False	

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