

CSc 372 — Comparative Programming Languages

33 : Icon — Generators

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November 30, 2005

1 Expressions as Generators

- Icon expressions are **generators**, they can return a sequence of values.
- Every expression has three possibilities: It can generate
 1. no values (\equiv failure),
 2. one value, or
 3. several values.

2 Expressions as Generators...

- Icon has many built-in generators, e.g. **i to j by k**. The following two statements are equivalent:

```
every i := j to k do p(i)
every p(j to k)
```

- **every e** asks **e** to generate as many values as it possibly can, by backtracking into it until it fails.
- **every e1 do e2** evaluates **e2** for every value generated by **e1**.

3 Expressions as Generators...

- The number of values a generator will produce depends on the environment in which it is invoked:

```
][ write(1 to 5);
1
  r1 := 1 (integer)
][ every write(1 to 5);
1
2
3
4
5
Failure
```

4 find

- `find(e1, e2)` returns the positions within the string `e2` where the string `e1` occurs.
- `find("wh", "who, what, when")` has three possible solutions and hence generates three values.

```
# 123456789012345
][ every i:=find("wh","who, what, when")
    do write(i);
1
6
12
Failure
```

5 Goal-Directed Evaluation

- Expression evaluation in Icon is **goal-directed**; you always try to make every expression succeed and return a value, if at all possible.
- In the example below, `find` first returns 1. This makes `((i :=...) > 10)` fail. Next `find` generates 14 which makes `((i :=...) > 10)` succeed, and `write` is executed.

```
S := "where and at what time?"
if (i := find("wh", S) > 10) then
    write(i)
```

6 Goal-Directed Evaluation...

```
][ 10 < (1 to 12);
    r34 := 11
][ every write( 10 < (1 to 12));
11
12
Failure
```

7 Counting Vowels

```
procedure main()
    v := 0
    while line := read() do {
        every c := !line do
            if c == !"aeiouAEIOU" then
                v += 1
    }
    write("vowels=",v)
end

> vowels
hi there
vowels=3
```

8 find...

- The expression

```
S := "where and at what time?"
][ every i := 10<find("wh",S) do write(i);
14
```

can also be written

```
][ every write(10 < find("wh",S));
14
```

9 File Operations

- The following statement copies a file `f1` row-by-row to another file `f2`:

```
while write(f2, read(f1))
```

- Note that `read` is not a generator — hence the use of `while` rather than `every`.

10 Bang!

- `!S` Generates all the characters from the string `S`, or all the elements of the list/table/set `S`.
- `every write(!S)` writes all the characters from the string `S`, one character per line.
- If `S` is a variable then `!S` will generate variables that can be assigned to.

11 Backtracking

- `&fail` always fails.

```
][ &fail;
Failure
][ 3;
   r38 := 3 (integer)
][ 3 + &fail;
Failure
][ 3 + numeric("pi");
Failure
```

12 Bang! — Examples

- Different ways to write the elements of a list:

```
][ L := [1,2,3];
][ every i := !L do write(i);
1
2
3
```

```

][ every write(!L);
1
2
3

```

13 Bang! — Examples...

```

][ every write(L[1 to 3]);
1
2
3
][ write(!L) + &fail;
1
2
3

```

14 Bang!...

- If **L** in **!L** generates variables, then they can be assigned to:

```

][ every !L := 5;
][ L;
   r16 := L1:[5,5,5]
][ !L := 1;
][ L;
   r24 := L1:[1,5,5]

```

15 Bang!...

- Note that literal strings cannot be assigned to:

```

][ S := "bye";
][ write(!S);
b
][ every write(!S);
b
y
e
Failure
][ every !S := "m";
][ S;
   r30 := "mmm" (string)
][ every !"bye" := "m";
Run-time error 111

```

16 Other Built-In Generators

?S Generates random elements from the set, string, table, etc.**S**.

upto(C, S) Generate all the positions in the string **S**, where the characters in **C** occur. **C** is a special construction called a **cSet**, a set of characters. **Csets** are written in single quotes, strings in doubles.

```
12345678901234
upto ('xyz', "zebra-ox-young")
generates {1, 8, 10}
```

Alternation

17 Alternation

- `expr1 | expr2` generates the values from `expr1`, then from `expr2`.
- `1 | 2 | 3` is the same as `1 to 3`.
- `(1 to 3) | (4 to 6)` is the same as `1 to 6`.
- `&fail | 3` generates `3`.
- `(1=2) | 3` generates `3`.
- `(1=1) | 3` generates `1,3` (since `1=1` succeeds and produces `1`).

18 Variable generation

- The expression `x | y` generates the *variables* `x` and `y`.
- The expression `every (x | y) := 0` is equivalent to `x := 0; y := 0`

19 Terminating Execution

- The built-in procedure `stop(s)` writes `s` and terminates execution.
- A common idiom is `x := p() | stop("error")`. If `p()` fails, then stop and write `"error"`, otherwise assign the result of `p()` to `x`.

20 Variable generation

`every i := (0 | 1) do write (i)` First write `0` then `1`.

`every (x | y) := 0 x := 0; y := 0`

```
][ x := 1;
][ y:= 2;
][ every write(x|y);
1
2
][ every (x|y) := 42;
][ every write(x|y);
42
42
```

21 Examples

```
][ every write(1 | 2 | !"45" | 6);  
1  
2  
4  
5  
6  
][ write((1 | 2 | 3) > 2);  
2  
    r13 := 2  
][ write(2 < (1 | 2 | 3));  
3  
    r14 := 3
```

22 Examples

```
][ x := 5;  
    r16 := 5  
][ y := 6;  
    r19 := 6  
][ (x | y) = 6;  
    r20 := 6
```

Procedures as Generators

23 Procedures as Generators

Procedures are really generators; they can return 0, 1, or a sequence of results. There are three cases

fail The procedure fails and generates no value.

return e The procedure generates one value, **e**.

suspend e The procedure generates the value **e**, and makes itself ready to possibly generate more values.

24 Example

```
procedure To(i,j)
  while i <= j do {
    suspend i
    i+:= 1
  }
end

procedure main()
  every k := To(1,3) do
    write(k)
  end
```

25 simple.icn

```
procedure P()
  suspend 3
  suspend 4
  suspend 5
end

procedure main()
  every k := P() do
    write(k)
  end
```

26 simple.icn...

```
> setenv TRACE 100
> simple
simple.icn :      main()
simple.icn :      8 | P()
simple.icn :      2 | P suspended 3
3
simple.icn :      9 | P resumed
```

```
simple.icn : 3 | P suspended 4
4
simple.icn : 9 | P resumed
simple.icn : 4 | P suspended 5
5
simple.icn : 9 | P resumed
simple.icn : 5 | P failed
simple.icn : 10 main failed
```

27 simple.icn...

- Remember *goal-directed evaluation* — Icon will resume a generator as many times as necessary in order to try to make an expression succeed.
- The number of times a generator is invoked also depends on the context.

28 simple.icn...

```
][ .inc simple.icn;
][ P();
   r1 := 3 (integer)
][ every write(P());
3
4
5
][ P()=4;
   r3 := 4
][ P() + 10;
   r4 := 13
```


Bounded Expressions

29 Bounded Expressions

- Unlike Prolog, backtracking in Icon is *bounded*. This means that a generator that appears in certain parts of certain control constructs will never generate more than one value.
- `if e1 then e2 else e3` — `e1` is bounded, `e2` and `e3` are not.
- `while e1 do e2` — `e1` and `e2` are both bounded.
- `every e1 do e2` — `e1` is not bounded but `e2` is.
- `{e1, e2, ..., en}` — `e1, e2, ...` are bounded but `en` is not.

30 Example

```
][ if write(P()) then &fail else &fail;
3
Failure
][ (if P() then write(P()) else 1) & &fail;
3
4
5
Failure
```

31 Example...

```
][ every i := P() do write(i);
3
4
5
Failure
][ while i := P() do write(i);
3
3
3...
][ {write(P()); 42} & &fail;
3
```

32 Example...

```
][ every i := {write(1 to 5); 42} do write(i);
1
42
][ every i := {write(1 to 5); 10 to 12} do write(i);
1
10
```

```
11
12
][ every i := {write(1 to 5); write(100 to 105); 10 to 12} do write(i);
1
100
10
11
12
```

Summary

33 Readings

- Read [Christopher](#), pp. 35--42, 44, 56--57.
- Alternatively, read [Griswold&Griswold](#), pp. 87--95.

34 Acknowledgments

- Some material on these slides has been modified from William Mitchell's Icon notes: <http://www.cs.arizona.edu/classes/cs372/fall03/handouts.html>.
- Some material on these slides has been modified from Thomas W Christopher's Icon Programming Language Handbook, <http://www.tools-of-computing.com/tc/CS/iconprog.pdf>.