
CSc 372

Comparative Programming Languages

33 : Icon — Generators

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Expressions as Generators

- Icon expressions are **generators**, they can return a sequence of values.
- Every expression has three possibilities: It can generate
 1. no values (\equiv failure),
 2. one value, or
 3. several values.

Expressions as Generators...

- Icon has many built-in generators, e.g. `i to j by k`.
The following two statements are equivalent:

```
every i := j to k do p(i)
every p(j to k)
```

- `every e` asks `e` to generate as many values as it possibly can, by backtracking into it until it fails.
- `every e1 do e2` evaluates `e2` for every value generated by `e1`.

Expressions as Generators...

- The number of values a generator will produce depends on the environment in which it is invoked:

```
][ write(1 to 5);  
1  
    r1 := 1 (integer)  
][ every write(1 to 5);  
1  
2  
3  
4  
5  
Failure
```

find

- `find(e1, e2)` returns the positions within the string `e2` where the string `e1` occurs.
- `find("wh", "who, what, when")` has three possible solutions and hence generates three values.

```
#                               123456789012345
|[ every i:=find("wh", "who, what, when")
    do write(i);
1
6
12
Failure
```

Goal-Directed Evaluation

- Expression evaluation in Icon is **goal-directed**; you always try to make every expression succeed and return a value, if at all possible.
- In the example below, **find** first returns 1. This makes **((i := ...) > 10)** fail. Next **find** generates 14 which makes **((i := ...) > 10)** succeed, and **write** is executed.

```
S := "where and at what time?"  
if (i := find("wh", S) > 10) then  
    write(i)
```

Goal-Directed Evaluation...

```
] [ 10 < (1 to 12);  
    r34 := 11  
] [ every write( 10 < (1 to 12));  
11  
12  
Failure
```

Counting Vowels

```
procedure main()  
  v := 0  
  while line := read() do {  
    every c := !line do  
      if c == !"aeiouAEIOU" then  
        v += 1  
    }  
  write("vowels=", v)  
end
```

```
> vowels  
hi there  
vowels=3
```


find...

- The expression

```
S := "where and at what time?"  
][ every i := 10 < find("wh", S) do write(i);  
14
```

can also be written

```
][ every write(10 < find("wh", S));  
14
```

File Operations

- The following statement copies a file `f1` row-by-row to another file `f2`:

```
while write(f2, read(f1))
```

- Note that `read` is not a generator — hence the use of `while` rather than `every`.

Bang!

- `!S` Generates all the characters from the string `S`, or all the elements of the list/table/set `S`.
- `every write(!S)` writes all the characters from the string `S`, one character per line.
- If `S` is a variable then `!S` will generate variables that can be assigned to.

Backtracking

● `&fail` always fails.

```
[[ &fail;  
Failure  
[[ 3;  
    r38 := 3 (integer)  
[[ 3 + &fail;  
Failure  
[[ 3 + numeric("pi");  
Failure
```

Bang! — Examples

- Different ways to write the elements of a list:

```
] [ L := [1, 2, 3];  
] [ every i := !L do write(i);  
1  
2  
3  
] [ every write(!L);  
1  
2  
3
```

Bang! — Examples...

```
] [ every write(L[1 to 3]);
```

```
1
```

```
2
```

```
3
```

```
] [ write(!L) + &fail;
```

```
1
```

```
2
```

```
3
```

Bang!...

- If L in $!L$ generates variables, then they can be assigned to:

```
[ [ every !L := 5 ;  
  [ L ;  
    r16 := L1 : [ 5 , 5 , 5 ]  
  [ !L := 1 ;  
    [ L ;  
      r24 := L1 : [ 1 , 5 , 5 ]
```

Bang!...

- Note that literal strings cannot be assigned to:

```
][ S := "bye" ;  
][ write(!S) ;  
b  
][ every write(!S) ;  
b  
y  
e  
Failure  
][ every !S := "m" ;  
][ S ;  
    r30 := "mmm"    (string)  
][ every !"bye" := "m" ;  
Run-time error 111
```


Other Built-In Generators

`?S` Generates random elements from the set, string, table, etc.`S`.

`upto(C, S)` Generate all the positions in the string `S`, where the characters in `C` occur. `C` is a special construction called a `CSet`, a set of characters. `CSets` are written in single quotes, strings in doubles.

```
          12345678901234
upto ( 'xyz' , "zebra-ox-young" )
      generates {1, 8, 10}
```

Alternation

Alternation

- `expr1 | expr2` generates the values from `expr1`, then from `expr2`.
- `1 | 2 | 3` is the same as `1 to 3`.
- `(1 to 3) | (4 to 6)` is the same as `1 to 6`.
- `&fail | 3` generates `3`.
- `(1=2) | 3` generates `3`.
- `(1=1) | 3` generates `1, 3` (since `1=1` succeeds and produces `1`).

Variable generation

- The expression $x \mid y$ generates the *variables* x and y .
- The expression $\text{every } (x \mid y) := 0$ is equivalent to $x := 0; y := 0$

Terminating Execution

- The built-in procedure `stop(s)` writes `s` and terminates execution.
- A common idiom is `x := p() | stop("error")`. If `p()` fails, then stop and write `"error"`, otherwise assign the result of `p()` to `x`.

Variable generation

`every i := (0 | 1) do write (i)` First write 0 then 1.

`every (x | y) := 0 x := 0; y := 0`

`] [x := 1;`

`] [y := 2;`

`] [every write(x|y);`

`1`

`2`

`] [every (x|y) := 42;`

`] [every write(x|y);`

`42`

`42`

Examples

```
] [ every write(1 | 2 | !"45" | 6) ;
```

```
1
```

```
2
```

```
4
```

```
5
```

```
6
```

```
] [ write((1 | 2 | 3) > 2) ;
```

```
2
```

```
    r13 := 2
```

```
] [ write(2 < (1 | 2 | 3)) ;
```

```
3
```

```
    r14 := 3
```

Examples

```
[ [ x := 5;  
    r16 := 5  
  ]  
[ [ y := 6;  
    r19 := 6  
  ]  
[ (x | y) = 6;  
  r20 := 6
```

Procedures as Generators

Procedures as Generators

Procedures are really generators; they can return 0, 1, or a sequence of results. There are three cases

`fail` The procedure fails and generates no value.

`return e` The procedure generates one value, `e`.

`suspend e` The procedure generates the value `e`, and makes itself ready to possibly generate more values.

Example

```
procedure To(i,j)
  while i <= j do {
    suspend i
    i+:= 1
  }
end

procedure main()
  every k := To(1,3) do
    write(k)
  end
```

simple.icn

```
procedure P()  
    suspend 3  
    suspend 4  
    suspend 5  
end
```

```
procedure main()  
    every k := P() do  
        write(k)  
end
```

simple.icn...

```
> setenv TRACE 100
```

```
> simple
```

```
      :      main( )
simple.icn :      8      | P( )
simple.icn :      2      | P suspended 3
3
simple.icn :      9      | P resumed
simple.icn :      3      | P suspended 4
4
simple.icn :      9      | P resumed
simple.icn :      4      | P suspended 5
5
simple.icn :      9      | P resumed
simple.icn :      5      | P failed
simple.icn :     10      | main failed
```

simple.icn...

- Remember **goal-directed evaluation** — Icon will resume a generator as many times as necessary in order to try to make an expression succeed.
- The number of times a generator is invoked also depends on the context.

simple.icn...

```
][ .inc simple.icn;  
][ P();  
    r1 := 3 (integer)  
][ every write(P());  
3  
4  
5  
][ P()=4;  
    r3 := 4  
][ P() + 10;  
    r4 := 13
```

Bounded Expressions

Bounded Expressions

- Unlike Prolog, backtracking in Icon is **bounded**. This means that a generator that appears in certain parts of certain control constructs will never generate more than one value.
- `if e1 then e2 else e3` — `e1` is bounded, `e2` and `e3` are not.
- `while e1 do e2` — `e1` and `e2` are both bounded.
- `every e1 do e2` — `e1` is not bounded but `e2` is.
- `{e1, e2, ..., en}` — `e1, e2, ...` are bounded but `en` is not.

Example

```
][ if write(P()) then &fail else &fail;
```

```
3
```

```
Failure
```

```
][ (if P() then write(P()) else 1) & &fail;
```

```
3
```

```
4
```

```
5
```

```
Failure
```

Example...

```
] [ every i := P() do write(i);
```

```
3
```

```
4
```

```
5
```

Failure

```
] [ while i := P() do write(i);
```

```
3
```

```
3
```

```
3...
```

```
] [ {write(P()); 42} & &fail;
```

```
3
```

Example...

```
] [ every i := {write(1 to 5); 42} do write(i);  
1  
42  
]  
] [ every i := {write(1 to 5); 10 to 12} do wri  
1  
10  
11  
12  
]  
] [ every i := {write(1 to 5); write(100 to 105)  
1  
100  
10  
11  
12
```

Summary

Readings

- Read Christopher, pp. 35--42, 44, 56--57.
- Alternatively, read Griswold&Griswold, pp. 87--95.

Acknowledgments

- Some material on these slides has been modified from William Mitchell's Icon notes:

<http://www.cs.arizona.edu/classes/cs372/fall03/handouts.html>.

- Some material on these slides has been modified from Thomas W Christopher's Icon Programming Language Handbook,

<http://www.tools-of-computing.com/tc/CS/iconprog.pdf>.