Superoptimizers A Look at the Smallest Program By Henry Massalin

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Problem Statement

- Find an optimal algorithm for compiling a function.
 - What does optimal mean?
 - Speed
 - Correctness
 - In this case, assembly output should be small as possible

Superoptimizer

- Given as input a set of instructions and a compiled function. The superoptimizer finds the smallest program that is equivalent to the input function.
- Most critical part of the tool is to determine if two pieces of code are equivalent a)Probabilistic Test

b)Pruning

Cool Example 1

- A function that returns the sign bit of X.
 if(x>0) return 1;
 else if(x < 0) return -1;
 else return 0;
- Typical assembly code

 le d0,0
 branch L1
 return 1
 L1: ge d0,0
 branch L2
 return -1
 L2: return 0

Superoptimizer
add d0,d0
subx.1 d1,d1
negx.1 d0
addx.1 d1,d1

Cool Example 2

- A function that returns the minimum of (x,y)
 if(x<y) return x;
 else
 return y;
- Typical assembly code

 ge d0,d1
 branch L1
 return d0
 L1: return d1

- Superoptimizer
 Sub.1 d1,d0 (x-y) subx.1 d2,d2 and.1 d2,d0 add.1 d1,d0
- This can be faster on a pipelined architecture where there might be a delay for a jump

Cool Example 3

 A function that returns the absolute value of x if(x>=0) return x;

else

return -x;

Typical assembly code

 lt d0,0
 branch L1
 return d0
 L1: neg d0
 return d0

Superoptimizer Move.1 d0,d1 add.1 d1,d1 subx.1 d1,d1 eor.1 d1,d0 sub.1 d1,d0

High Level Algorithm

- Choose a subset of the instruction set and store in a table for reference
- Superoptimizer consulates this table and generates all combinations of instructions
- Start with a instruction sets of length 1, length 2, etc.
- Test whether each program is equivalent to the original program
- Halt when an equivalent program is found.

Equivalence Test - Probabilistic Test

- Run the generated program on a select few smart inputs.
- If the output matches the source program, run the generated program through the entire test suite.
- Example inputs for signum a) Negative number, positive number, zero
- These test cases rules out programs that
 - a) Return same value regardless of input
 - b) Answers of the same sign
 - c) Return their argument

Pruning Search Space

- No need to evaluate sequences that cannot occur in a optimal solution
- Any longer sequences that have the same effect as shorter sequences cannot be part of the optimal

 a) move x,y; move x,y

b) and x,y; move z,y

- N-dimensional bit tables
 - For i<N, test a sequence of length i by accessing the ith dimension of the bit tables.
 - For each instruction in the sequence, if the lookup value is 1, no need to continue along this search path.

Limitations

- Even with optimizations, still an exponential search
 - 12 instruction functions take hours to find optimal solution
- Pointer verification is hard. Since pointers can point anywhere, all possibilities need to be checked (means all possible memory addresses)
 - On a 256-byte machine 2^{(2(256*8))} possible pointer locations
 - Probabilistic test is inconclusive with pointers
- Machine dependence, the input instruction set must be able to run on the machine Superoptimizer is running on.

Applications

- Superoptimizer is useful for
- Helping the compiler find optimizations for little tasks like

 a) multiplication by a constant. (E.g. array indexing)

b) Checking for equality between two values.

- 2) Equivalent sets of programs generated by Superoptimizer can aide the compiler in peephole optimization
- 3) Assembly programmers can use Superoptimizer to reduce the size of critical stdlib functionsa) Author gives example of rewriting printf function in 500
 - bytes.

Conclusion

- Superoptimizer is a tool that can be used to find optimal(shortest) equivalent programs.
- Even though it uses an exhaustive search
 - a) Probabilistic test and Pruning help reduce the time it takes to find the optimal solution.
 - b) Once equivalent optimal programs are found for common tasks they can be stored.
- The most interesting result is the relationship between arithmetic and logical instructions. And the ability they have to simulate branching instructions.
- Also, search structure lends itself to parallelism, slave processes can check for equivalence for a particular program and the input program. And report back to the master process.