University of Arizona

CSc 620

Security Through Obscurity

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The Java VM

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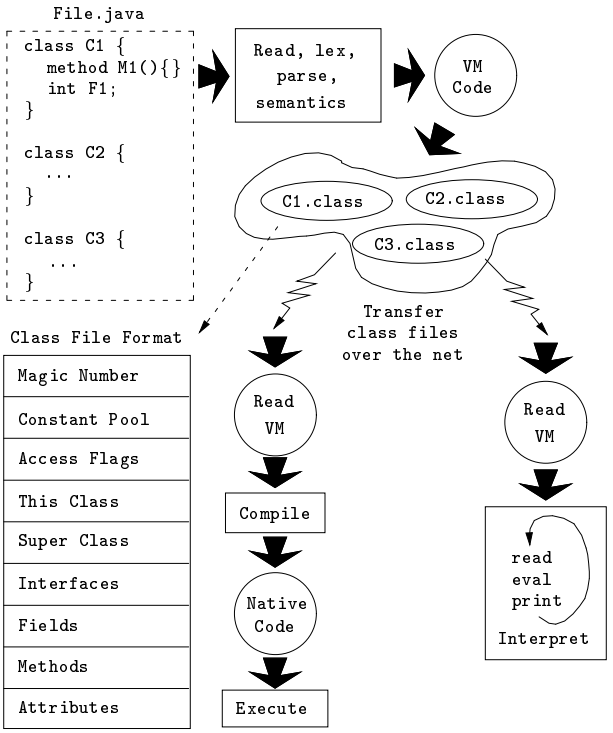
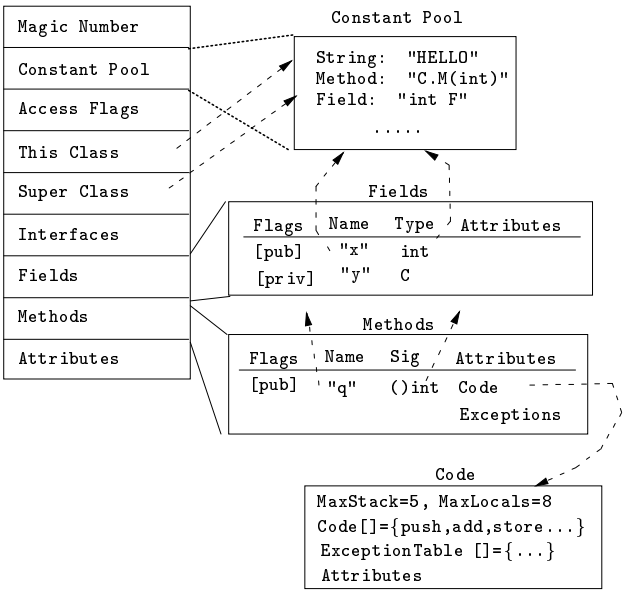
The Java Virtual Machine

- The Java VM has gone the "many complex instructions/large VM engine" way.
- Each Java source file may contain several Java classes. The Java compiler compiles each of these classes to a single Java *class file*.
- The Java class file stores all necessary data regarding the class. There is a symbol table (called the *Constant Pool*) which stores strings, large literal integers and floats, names and of all fields and methods.
- Each method is compiled to Java bytecode, a stack VM format.
- The class file is (almost) isomorphic to the source.

Compiling Java



The Java Class File Format



Java Byte Codes I

- The Java bytecodes can manipulate data in these formats: integers (32-bits), longs (64-bits), floats (32-bits), doubles (64-bits), shorts (16-bits), bytes (8-bits), object references (32/64-bit pointers), and arrays.
- The bytecodes are 1 byte wide.
- Each method can have up to 256 local variables and formal parameters. The bytecode reference these by number.
- Actually, we can have up to 65536 local vars. There is a special `wide` instruction that modifies load and store instructions to reference the high-numbered locals. Hack.
- The Java stack is 32-bits wide. Longs and doubles hence take two stack entries.

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Java Byte Codes II

- The bytecodes reference data from the class' constant pool. These references are 8 or 16 bits long. To push a reference to a literal string with constant pool # 4567, use `'ldc2 4567'`. If the # is 123, use `'ldc2 123'`.

$\texttt{int}_8$	An 8-bit integer value.
$\texttt{int}_{16}$	A 16-bit integer value.
$\texttt{int}_{32}$	A 32-bit integer value.
$\texttt{CP}_8$	An 8-bit constant pool index.
$\texttt{CP}_{16}$	A 16-bit constant pool index.
`FIdx`	An 8-bit local variable index.
$\texttt{FIdx}_{16}$	A 16-bit local variable index.
`CP[`i`]`	The i:th constant pool entry.
`Var[`i`]`	The i:th variable/formal parameter in the current method.

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Opcode	Mnemonic	Args	Stack	Description
0	`nop`		$[] \Rightarrow []$	
1	`aconst_null`		$[] \Rightarrow$ [null]	Push null object
2	`iconst_m1`		$[] \Rightarrow$ [-1]	Push -1
3...8	`iconst_`n		$[] \Rightarrow [n]$	Push integer constant $n, 0 \leq n \leq 5$
9...10	`lconst_`n		$[] \Rightarrow [n]$	Push long constant $n, 0 \leq n \leq 1$
11...13	`fconst_`n		$[] \Rightarrow [n]$	Push float constant $n, 0 \leq n \leq 2$
14...15	`dconst_`n		$[] \Rightarrow [n]$	Push double constant $n, 0 \leq n \leq 1$

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Opcode	Mnemonic	Args	Stack	Description
16	`bipush`	n:$\texttt{int}_8$	$[] \Rightarrow [n]$	Push 1-byte signed integer
17	`sipush`	n:$\texttt{int}_{16}$	$[] \Rightarrow [n]$	Push 2-byte signed integer
18	`ldc1`	n:$\texttt{CP}_8$	$[] \Rightarrow$ [CP[n]]	Push item from constant pool
19	`ldc2`	n:$\texttt{CP}_{16}$	$[] \Rightarrow$ [CP[n]]	Push item from constant pool
20	`ldc2w`	n:$\texttt{CP}_{16}$	$[] \Rightarrow$ [CP[n]]	Push long/double from constant pool

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Opcode	Mnemonic	Args	Stack
21 . . . 25	X`load`	n:`FIdx`	[] ⇒ [Var[n]]
	$X \in${i,l,f,d,a}, Load int, long, float, double, object from local var.		
26 . . . 29	`iload_`n		[] ⇒ [Var[n]]
	Load local integer var $n, 0 \leq n \leq 3$		
30 . . . 33	`lload_`n		[] ⇒ [Var[n]]
	Load local long var $n, 0 \leq n \leq 4$		
34 . . . 37	`fload_`n		[] ⇒ [Var[n]]
	Load local float var $n, 0 \leq n \leq 4$		
38 . . . 41	`dload_`n		[] ⇒ [Var[n]]
	Load local double var $n, 0 \leq n \leq 4$		

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Opcode	Mnemonic	Args	Stack
42 . . . 45	`aload_`n		[] ⇒ [Var[n]]
	Load local object var $n, 0 \leq n \leq 4$		
46 . . . 53	X`load`		$[A, I] \Rightarrow [V]$
	$X \in${ia,la,fa,da,aa,ba, ca,sa}. Push the value V (an int, long, etc.) stored at index I of array A.		
54 . . . 58	X`store`	n:`FIdx`	[Var[n]] ⇒ []
	$X \in${i,l,f,d,a}, Store int, long, float, double, object to local var.		
59 . . . 62	`istore_`n		[Var[n]] ⇒ []
	Store to local integer var $n, 0 \leq n \leq 3$		
63 . . . 66	`lstore_`n		[Var[n]] ⇒ []
	Store to local long var $n, 0 \leq n \leq 4$		

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Opcode	Mnemonic	Args	Stack
67 . . . 70	`fstore_`n		[Var[n]] ⇒ []
	Store to local float var $n, 0 \leq n \leq 4$		
71 . . . 74	`dstore_`n		[Var[n]] ⇒ []
	Store to local double var $n, 0 \leq n \leq 4$		
75 . . . 78	`astore_`n		[Var[n]] ⇒ []
	Store to local object var $n, 0 \leq n \leq 4$		
79 . . . 86	X`store`		$[A, I, V] \Rightarrow []$
	$X \in${ia,la,fa,da,aa,ba,ca,sa}. Store the value V (an int, long, etc.) at index I of array A.		
87	`pop`		$[A] \Rightarrow []$
	Pop top of stack.		

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Opcode	Mnemonic	Stack	Description
88	`pop2`	$[A, B] \Rightarrow []$	Pop 2 elements.
89	`dup`	$[V] \Rightarrow [V, V]$	
	Duplicate top of stack.		
90	`dup_x1`	$[B, V] \Rightarrow [V, B, V]$	Duplicate.
91	`dup_x2`	$[B, C, V] \Rightarrow [V, B, C, V]$	Duplicate.
92	`dup2`	$[V, W] \Rightarrow [V, W, V, W]$	Duplicate.
93	`dup2_x1`	$[A, V, W] \Rightarrow [V, W, A, V, W]$	Duplicate.
94	`dup2_x2`	$[A, B, V, W] \Rightarrow$ $[V, W, A, B, V, W]$	Duplicate.
95	`swap`	$[A, B] \Rightarrow [B, A]$	
	Swap top stack elements.		

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Opcode	Mnemonic	Stack	Description
96 . . . 99	X`add`	$[A, B] \Rightarrow [R]$	$X \in${i,l,d,f}. $R = A + B$
100 . . . 103	X`sub`	$[A, B] \Rightarrow [R]$	$X \in${i,l,d,f}. $R = A - B$
104 . . . 107	X`mul`	$[A, B] \Rightarrow [R]$	$X \in${i,l,d,f}. $R = A * B$
108 . . . 111	X`div`	$[A, B] \Rightarrow [R]$	$X \in${i,l,d,f}. $R = A/B$
112 . . . 115	X`mod`	$[A, B] \Rightarrow [R]$	$X \in${i,l,d,f}. $R = A\%B$
116 . . . 119	X`neg`	$[A] \Rightarrow [R]$	$X \in${i,l,d,f}. $R = -A$
120 . . . 121	X`shl`	$[A, B] \Rightarrow [R]$	$X \in${i,l}. $R = A << B$
122 . . . 123	X`shl`	$[A, B] \Rightarrow [R]$	$X \in${i,l}. $R = A >> B$
124 . . . 125	X`ushr`	$[A, B] \Rightarrow [R]$	$X \in${i,l}. $R = A >>> B$
126 . . . 127	X`and`	$[A, B] \Rightarrow [R]$	$X \in${i,l}. $R = A \&\& B$
128 . . . 129	X`or`	$[A, B] \Rightarrow [R]$	$X \in${i,l}. $R = A \| \| B$
130 . . . 131	X`xor`	$[A, B] \Rightarrow [R]$	$X \in${i,l}. $R = A$`xor`B

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Opcode	Mnemonic	Args	Stack
133 . . . 144	$X2Y$`cnv`		$[F] \Rightarrow [T]$
	Convert F from type X to T of type Y. $X \in${i,l,f,d}, $Y \in${i,l,f,d}.		
145 . . . 147	`i2`X		$[F] \Rightarrow [T]$
	$X \in${b,c,s}. Convert integer F to byte, char, or short.		
148,149,151	X`cmp`		$[A, B] \Rightarrow [V]$
	$X \in${l,f,d}. $A > B \Rightarrow V = 1$, $A < B \Rightarrow V = -1$, $A = B \Rightarrow V = 0$. $A =$ `NaN` $\vee B =$ `NaN` $\Rightarrow V = -1$		
150,152	X`cmp`		$[A, B] \Rightarrow [V]$
	$X \in${f,d}. $A > B \Rightarrow V = 1$, $A < B \Rightarrow V = -1$, $A = B \Rightarrow V = 0$. $A =$ `NaN` $\vee B =$ `NaN` $\Rightarrow V = 1$		
153 . . . 154	`if`$\diamond$	L:`int`$_{16}$	$[A] \Rightarrow []$
	$\diamond$={eq,ne,lt,ge,gt,le}. If $A \diamond 0$ goto $L +$ `pc`.		

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Opcode	Mnemonic	Args	Stack
159 . . . 164	`if_icomp`$\diamond$	L:`int`$_{16}$	$[A, B] \Rightarrow []$
	$\diamond$={eq,ne,lt,ge,gt,le}. If $A \diamond B$ goto $L +$ `pc`.		
165 . . . 166	`if_acmp`$\diamond$	L:`int`$_{16}$	$[A, B] \Rightarrow []$
	$\diamond$={eq,ne}. A, B are object refs. If $A \diamond B$ goto $L +$ `pc`.		
167	`goto`	I:`int`$_{16}$	$[] \Rightarrow []$
	Goto instruction I.		
168	`jsr`	I:`int`$_{16}$	$[] \Rightarrow []$
	Jump subroutine to instruction $I +$ `pc`.		
172 . . . 177	X`return`		$[V] \Rightarrow []$
	$X \in${i,f,l,d,a,v}. Return V.		
169	`ret`	L:`FIdx`	$[] \Rightarrow []$
	Return from subroutine. Address in local var L.		

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Opcode	Mnemonic	Args	Stack
170	`tableswitch`	D:`int`$_{32}$,l, h:`int`$_{32}$,o^{h-l+1}	$[K] \Rightarrow []$
	Jump through the K:th offset. Else goto D.		
171	`lookupswitch`	D:`int`$_{32}$,n:`int`$_{32}$,$(m, o)^n$	$[K] \Rightarrow []$
	If, for one of the (m, o) pairs, $K = m$, then goto o. Else goto D.		
178	`getstatic`	F:`CP`$_{16}$	$[] \Rightarrow [V]$
	Push value V of static field F.		
180	`getfield`	F:`CP`$_{16}$	$[R] \Rightarrow [V]$
	Push value V of field F in object R.		
179	`putstatic`	F:`CP`$_{16}$	$[] \Rightarrow [V]$
	Store value V into static field F.		
181	`putfield`	F:`CP`$_{16}$	$[R, V] \Rightarrow []$
	Store value V into field F of object R.		

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Opcode	Mnemonic	Args	Stack
182	`invokevirtual`	P:$\texttt{CP}_{16}$	$[R, A_1, A_2, \ldots] \Rightarrow []$
	Call virtual menthod P, with arguments $A_1 \cdots A_n$, through object reference R.		
183	`invokespecial`	P:$\texttt{CP}_{16}$	$[R, A_1, A_2, \ldots] \Rightarrow []$
	Call private/init/superclass menthod P, with arguments $A_1 \cdots A_n$, through object reference R.		
184	`invokestatic`	P:$\texttt{CP}_{16}$	$[A_1, A_2, \ldots] \Rightarrow []$
	Call static menthod P with arguments $A_1 \cdots A_n$.		
185	`invokeinterface`	P:$\texttt{CP}_{16}$,n:$\texttt{int}_{16}$	$[R, A_1, A_2, \ldots] \Rightarrow []$
	Call interface menthod P, with n arguments $A_1 \cdots A_n$, through object reference R.		
187	`new`	T:$\texttt{CP}_{16}$	$[] \Rightarrow [R]$
	Create a new object R of type T.		

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Opcode	Mnemonic	Args	Stack
188	`newarray`	T:$\texttt{int}_8$	$[C] \Rightarrow [R]$
	Allocate new array R, element type T, C elements long.		
191	`athrow`		$[R] \Rightarrow [?]$
	Throw exception.		
193	`instanceof`	C:$\texttt{CP}_{16}$	$[R] \Rightarrow [V]$
	Push 1 if object R is an instance of class C. Else push 0.		
194	`monitorenter`		$[R] \Rightarrow []$
	Get lock for object R.		
195	`monitorexit`		$[R] \Rightarrow []$
	Release lock for object R.		
196	`wide`	C:$\texttt{int}_8$,I:$\texttt{FIdx}_{16}$	$[] \Rightarrow []$
	Perform opcode C on variable `Var[`I`]`. C is one of the load/store instructions.		

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Opcode	Mnemonic	Args	Stack
197	`multianewarray`	T:$\texttt{CP}_{16}$,D:$\texttt{CP}_8$	$[d_1, d_2, \ldots] \Rightarrow [R]$
	Create new D-dimensional multidimensional array R. $d_1, d_2, \ldots$ are the dimension sizes.		
198	`ifnull`	L:$\texttt{int}_{16}$	$[V] \Rightarrow []$
	If $V =$ `null` goto L.		
199	`ifnonnull`	L:$\texttt{int}_{16}$	$[V] \Rightarrow []$
	If $V \neq$ `null` goto L.		
200	`goto_w`	I:$\texttt{int}_{32}$	$[] \Rightarrow []$
	Goto instruction I.		
201	`jsr_w`	I:$\texttt{int}_{32}$	$[] \Rightarrow []$
	Jump subroutine to instruction I.		

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JVM Example I

```
void spin() {
    int i;
    for (i = 0; i < 100; i++) {
        ;    // Loop body is empty
    }}
```

$\Downarrow$

```
0  iconst_0      // Push int constant 0
1  istore_1      // Store into local 1 (i=0)
2  goto 8        // First time through don't
                 // increment
5  iinc 1 1      // Increment local 1 by 1
                 // (i++)
8  iload_1       // Push local 1 (i)
9  bipush 100    // Push int constant (100)
11 if_icmplt 5   // Compare, loop
                 // if (i < 100)
14 return        // Return void when done
```

JVM Example II

```
    void dspin() {
        double i;
        for (i = 0.0; i < 100.0; i++) {
         ;       // Loop body is empty
    }}
```

```
0   dconst_0  // Push double constant 0.0
1   dstore_1  // Store into locals 1 and 2
              // (i = 0.0)
2   goto 9    // First time no incr
5   dload_1   // Push double
6   dconst_1  // Push double 1.o onto stack
7   dadd      // Add;
8   dstore_1  // Store result in locals
              // 1 and 2
9   dload_1   // Push local
10  ldc2_w #4 // Double 100.000000
13  dcmpg
14  iflt 5    // Compare, loop if <
              // (i < 100.000000)
17  return    // Return void when done
```

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JVM Example III

```
double doubleLocals(double d1, double d2) {
    return d1 + d2;
}
```

```
0 dload_1 // First argument in locals 1 and 2
1 dload_3 // Second argument in locals 3 and 4
2 dadd    // Each also uses two words on stack
3 dreturn
```

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JVM Example IV

```
    int align2grain(int i, int grain) {

        return ((i + grain-1) & ~(grain-1));
    }
```

```
0   iload_1
1   iload_2
2   iadd
3   iconst_1
4   isub
5   iload_2
6   iconst_1
7   isub
8   iconst_m1
9   ixor
10  iand
11  ireturn
```

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JVM Example V

```
    void useManyNumeric() {
        int i = 100; int j = 1000000;
        long l1 = 1; long l2 = 0xffffffff;
        double d = 2.2;
    }
```

```
0  bipush 100 // Push a small int
2  istore_1
3  ldc #1     // Integer 1000000; a larger
              // int value uses ldc
5  istore_2
6  lconst_1   // A tiny long value
7  lstore_3
8  ldc2_w #6  // A long 0xffffffff. A long
              // constant value.
11 lstore 5
13 ldc2_w #8  // Double 2.200000
16 dstore 7
```

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JVM Example VI

```
void whileInt() {
    int i = 0;
    while (i < 100) {
        i++;
    }
}
```

```
0   iconst_0
1   istore_1
2   goto 8
5   iinc 1 1
8   iload_1
9   bipush 100
11  if_icmplt 5
14  return
```

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JVM Example VII

```
int lessThan100(double d) {
    if (d < 100.0) {
        return 1;
    } else {
        return -1;
    }
}
```

```
0   dload_1
1   ldc2_w #4 // Double 100.000000
4   dcmpg     // Push 1 if d is NaN or
              //     d < 100.000000;
              // push 0 if d == 100.000000
5   ifge 10   // Branch on 0 or 1
8   iconst_1
9   ireturn
10  iconst_m1
11  ireturn
```

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JVM Example VIII

```
int add12and13() {return addTwo(12, 13);}
```

```
0  aload_0       // Push this local 0 (this) onto stack
1  bipush 12     // Push int constant 12 onto stack
3  bipush 13     // Push int constant 13 onto stack
5  invokevirtual #4 // Method Example.addtwo(II)I
8  ireturn       // Return int on top of stack; it is
                 // the int result of addTwo()
```

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JVM Example IX

```
Object create() {return new Object();}
```



```
0  new #1        // Class java.lang.Object
3  dup
4  invokespecial #4 // Method java.lang.Object.<init>()V
7  areturn
```

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JVM Example X

```
void createBuffer() {
   int buf[]; int bsz = 100; int val=12;
   buf = new int[bsz]; buf[10]=val;
   value = buf[11]; }
```

```
0   bipush 100    // Push bsz
2   istore_2      // Store bsz in local 2
3   bipush 12     // Push val
5   istore_3      // Store val in local 3
6   iload_2       // Push bsz...
7   newarray int  // and create new int array
9   astore_1      // Store new array in buf
10  aload_1       // Push buf
11  bipush 10     // Push constant 10
13  iload_3       // Push val
14  iastore       // Store val at buf[10]
15  aload_1       // Push buf
16  bipush 11     // Push constant 11
18  iaload        // Push value at buf[11]
19  istore_3      // ...and store it in value
20  return
```


JVM Example XI

```
int chooseNear(int i) {
    switch (i) {
        case 0:  return 0;
        case 2:  return 2;
        default: return -1;
}}
```

```
0   iload_1      // Load local 1 (argument i)
1   tableswitch 0 to 2:
    0: 28        // If i is 0, continue at 28
    1: 32        // If i is 1, continue at 34
    2: 30        // If i is 2, continue at 32
    default:34   // Otherwise, continue at 34
28  iconst_0     // i was 0; push int 0...
29  ireturn      // ...and return it
30  iconst_2     // i was 2; push int 2...
31  ireturn      // ...and return it
32  iconst_m1    // otherwise push int -1...
33  ireturn      // ...and return it
```