

## **Plan for Today (>30% of TCEs are in)**

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### **Discussion about Final**

- 2 sides of 8.5x11 sheet of paper
- See posted notes for tomorrow's recitation for what could be on final.
- Friday will have review and can do examples.

### **PA5 Suggestions**

### **Register allocation for expressions**

### **What are IO monads?**

# Expression Evaluation

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## Sethi-Ullman Register allocation

```
label(node)
  if node is leaf then node.label = 1
  else if node is binary
    if node.left.label == node.right.label
      then node.label = node.left.label + 1
    else
      node.label = max( node.left.label,
                        node.right.label )
  else if node is unary
    node.label = node.child.label
```

## Questions to understand how to answer

- How many registers are needed if not using memory (aka push/pop)?
- Order of evaluation to make this register count work?
- What if an operator has more than 2 operands?

# THE IO MONAD

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Lightly edited with permission, Michelle Strout 4/13/15

Reading: "[Tackling the Awkward Squad](#)," Sections 1-2  
"[Real World Haskell](#)," Chapter 7: I/O

Thanks to Simon Peyton Jones for many of these slides.

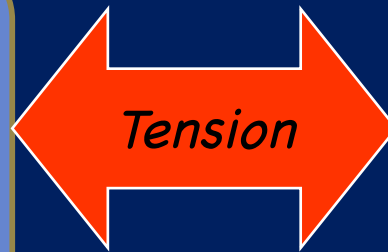
# Why Monads?

- Predictive parser
  - Passed around a list of tokens while processing.
- PA3 MeggyJava compiler
  - Passed around a number to create unique labels for code generation.
- PA4 MeggyJava compiler
  - Passing around a symbol table with parameter and method type and code generation information.
- Monads will help us abstract away some of that passing around.

# *Monadic Input and Output*

# The Problem

*A functional program defines a pure function, with **no side effects**.*



*The whole point of running a program is to have **some side effect**.*

# Monadic I/O: The Key Idea

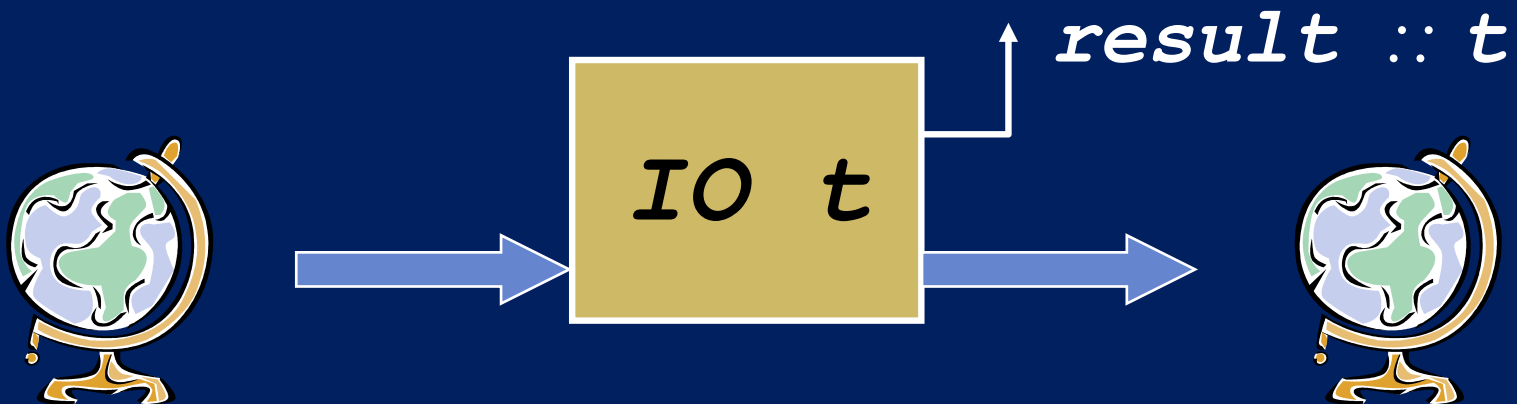
*A value of type  $(\text{IO } t)$  is an "action." When performed, it may do some input/output before delivering a result of type  $t$ .*



# A Helpful Picture

A value of type `(IO t)` is an "action." When performed, it may do some input/output before delivering a result of type `t`.

```
type IO t = World -> (t, World)
```





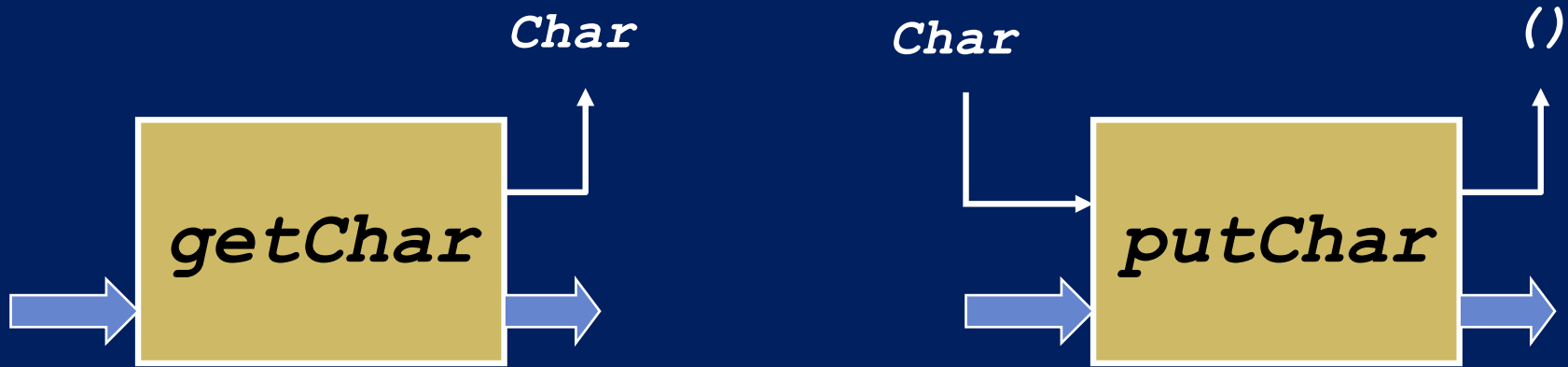
# Actions are First Class

*A value of type (`IO t`) is an "action." When performed, it may do some input/output before delivering a result of type `t`.*

```
type IO t = World -> (t, World)
```

- "Actions" are sometimes called "computations."
- An action is a **first-class value**.
- **Evaluating** an action has no effect; **performing** the action has the effect.

# Simple I/O



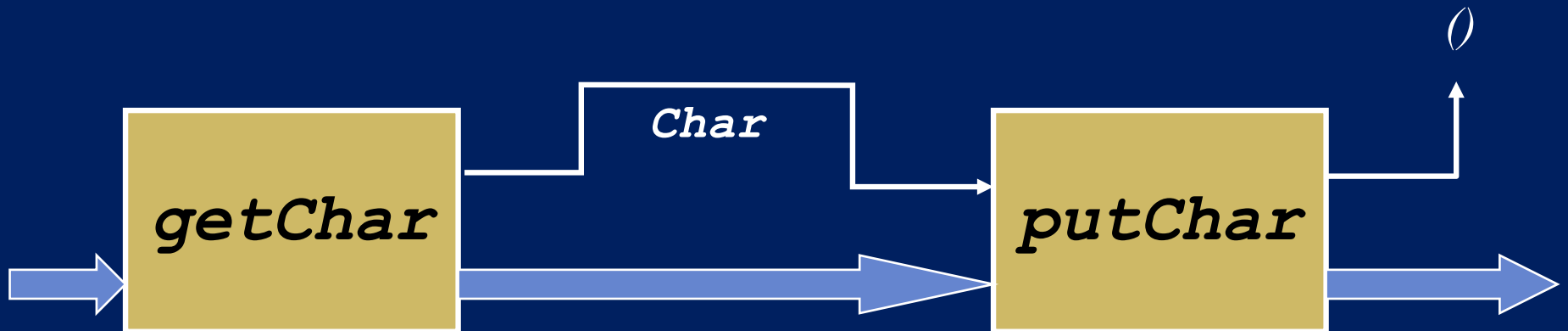
```
getChar :: IO Char  
putChar :: Char -> IO ()
```

```
main :: IO ()  
main = putChar 'x'
```

Main program is an  
action of type `IO ()`

# Connection Actions

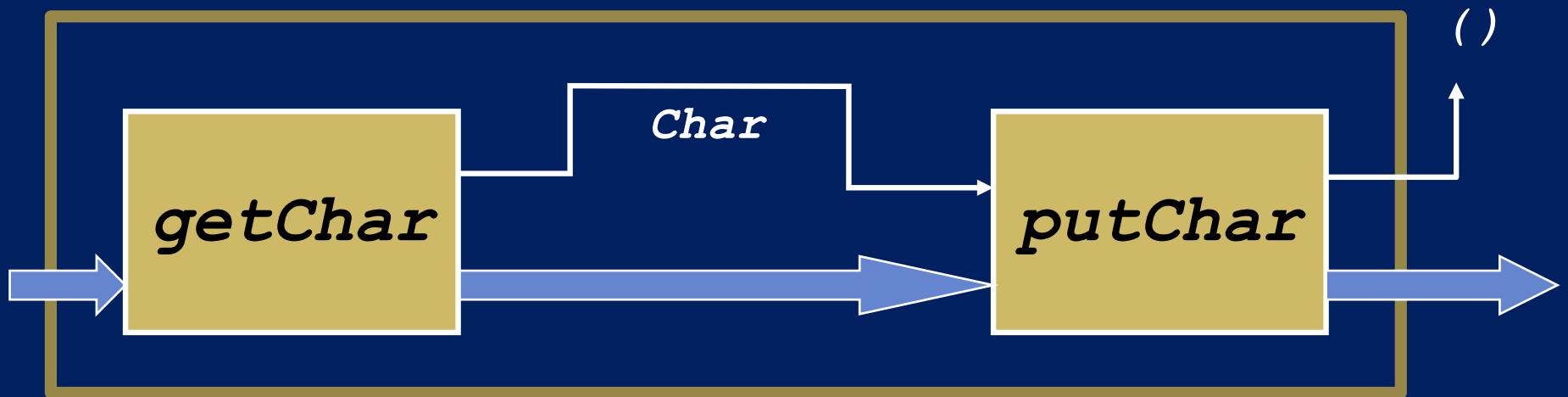
To read a character and then write it back out, we need to *connect* two actions.



The “bind” combinator lets us make these connections.

# The Bind Combinator ( $>>=$ )

```
 $(>>=) :: IO\ a \rightarrow (a \rightarrow IO\ b) \rightarrow IO\ b$ 
```

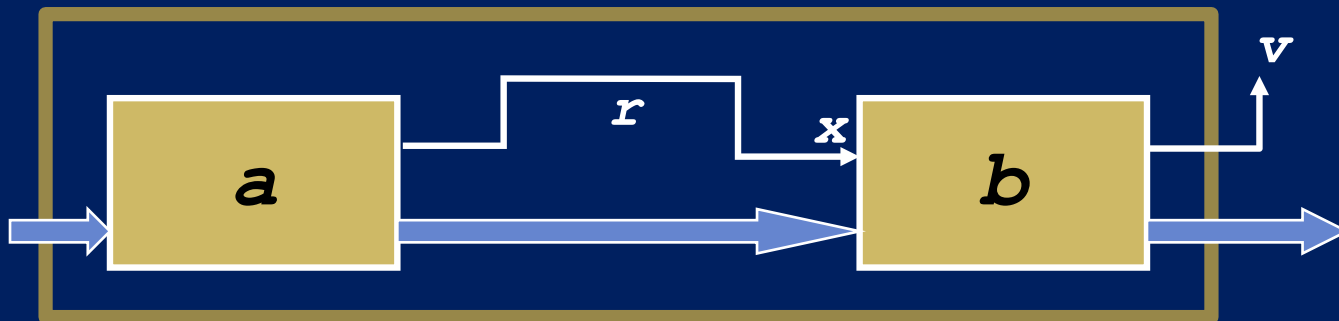


- We have connected two actions to make a new, bigger action.

```
echo :: IO ()  
echo = getChar >>= putChar
```

# The ( $>>=$ ) Combinator

- Operator is called **bind** because it *binds* the result of the left-hand action in the action on the right.
- Performing compound action  $a >>= \lambda x \rightarrow b$ :
  - performs action  $a$ , to yield value  $r$
  - applies function  $\lambda x \rightarrow b$  to  $r$
  - performs the resulting action  $b\{x \leftarrow r\}$
  - returns the resulting value  $v$



# Printing a Character Twice

```
echoDup :: IO ()  
echoDup = getChar    >>= (\c ->  
                putChar c >>= (\ () ->  
                putChar c ))
```

- The parentheses are optional because lambda abstractions extend “as far to the right as possible.”
- The `putChar` function returns unit, so there is no interesting value to pass on.

# The (>>) Combinator

- The “then” combinator (>>) does sequencing when there is no value to pass:

```
(>>) :: IO a -> IO b -> IO b  
m >> n = m >>= (\_ -> n)
```

```
echoDup :: IO ()  
echoDup = getChar >>= \c ->  
           putChar c >>  
           putChar c
```

```
echoTwice :: IO ()  
echoTwice = echo >> echo
```

# Getting Two Characters

```
getTwoChars :: IO (Char,Char)
getTwoChars = getChar    >>= \c1 ->
               getChar    >>= \c2 ->
               ?????
```

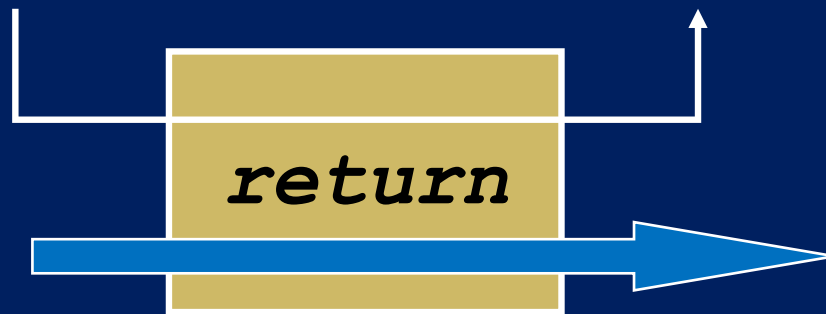
- We want to return `(c1,c2)`.
  - But, `(c1,c2) :: (Char, Char)`
  - And we need to return something of type `IO(Char, Char)`
- We need to have some way to convert values of “plain” type into the I/O Monad.



# The `return` Combinator

- The action `(return v)` does no IO and immediately returns `v`:

```
return :: a -> IO a
```



```
getTwoChars :: IO (Char, Char)  
getTwoChars = getChar >>= \c1 ->  
          getChar >>= \c2 ->  
          return (c1, c2)
```

# The "do" Notation

- The "do" notation adds syntactic sugar to make monadic code easier to read.

```
-- Plain Syntax
```

```
getTwoChars :: IO (Char,Char)
getTwoChars = getChar    >>= \c1 ->
               getChar    >>= \c2 ->
               return (c1,c2)
```

```
-- Do Notation
```

```
getTwoCharsDo :: IO (Char,Char)
getTwoCharsDo = do { c1 <- getChar ;
                    c2 <- getChar ;
                    return (c1,c2) }
```

- Do syntax designed to look imperative.

# Desugaring "do" Notation

- The "do" notation *only* adds syntactic sugar:

$$\mathit{do} \{ x \leftarrow e; es \} = e \gg= \backslash x \rightarrow \mathit{do} \{ es \}$$
$$\mathit{do} \{ e; es \} = e \gg \mathit{do} \{ es \}$$
$$\mathit{do} \{ e \} = e$$
$$\mathit{do} \{ \mathit{let} \ ds; es \} = \mathit{let} \ ds \ \mathit{in} \ \mathit{do} \{ es \}$$

*The scope of variables bound in a generator is the rest of the "do" expression.*

*The last item in a "do" expression must be an expression.*

# Syntactic Variations

- The following are equivalent:

```
do { x1 <- p1; ...; xn <- pn; q }
```

```
do  x1 <- p1; ...; xn <- pn; q
```

```
do x1 <- p1  
   ...  
   xn <- pn  
   q
```

*If the semicolons are omitted, then the generators must line up. The indentation replaces the punctuation.*

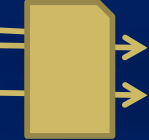
# Bigger Example


- The `getLine` function reads a line of input:

```
getLine :: IO [Char]
getLine = do { c <- getChar ;
              if c == '\n' then
                return []
              else
                do { cs <- getLine;
                   return (c:cs) }}
```

*Note the "regular" code mixed with the monadic operations and the nested "do" expression.*

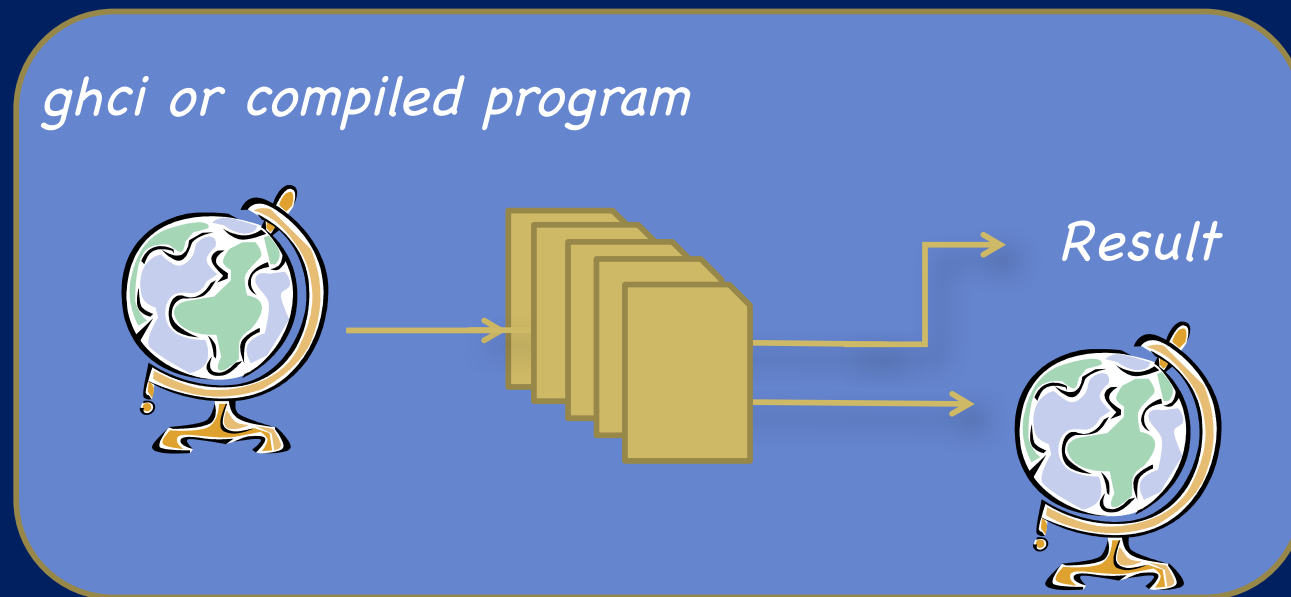
# An Analogy: Monad as Assembly Line

- Each action in the IO monad is a possible stage in an assembly line. 
- For an action with type `IO a`, the type
  - tags the action as suitable for the IO assembly line via the `IO` type constructor.
  - indicates that the kind of thing being passed to the next stage in the assembly line has type `a`.

- The `bind` operator “snaps” two stages `s1` and `s2` together to build a compound stage. 
- The `return` operator converts a pure value into a stage in the assembly line.
- The assembly line *does nothing* until it is turned on.
- The only safe way to “run” an IO assembly is to execute the program, either using `ghci` or running an executable.

# Powering the Assembly Line

- Running the program turns on the IO assembly line.
- The assembly line gets “the world” as its input and delivers a result and a modified world.
- The types guarantee that the world flows in a single thread through the assembly line.



# Implementation

- GHC uses world-passing semantics for the IO monad:

```
type IO t = World -> (t, World)
```

- It represents the “world” by an un-forgable token of type `World`, and implements `bind` and `return` as:

```
return :: a -> IO a
```

```
return a = \w -> (a, w)
```

```
(>>=) :: IO a -> (a -> IO b) -> IO b
```

```
(>>=) m k = \w -> case m w of (r, w') -> k r w'
```

- Using this form, the compiler can do its normal optimizations. The dependence on the world ensures the resulting code will still be single-threaded.
- The code generator then converts the code to modify the world “in-place.”



## But what does this mean?

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**Reference:** <http://stackoverflow.com/questions/3117583/is-haskell-truly-pure-is-any-language-that-deals-with-input-and-output-outside>

### Summary

- IO Monad value for Haskell main is like a program AST
- The IO Monad value is evaluated in the sense that IO actions are bound together sequentially including some IO actions that contain lambda function values based on input.
- The ghc compiler then converts these IO Monad values into C code and executes the C code at runtime.

**→ Show usage of monad for compiler example.**