
Principles of Programming Languages

Lecture 10

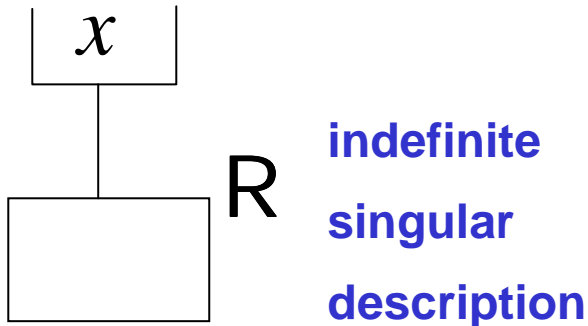
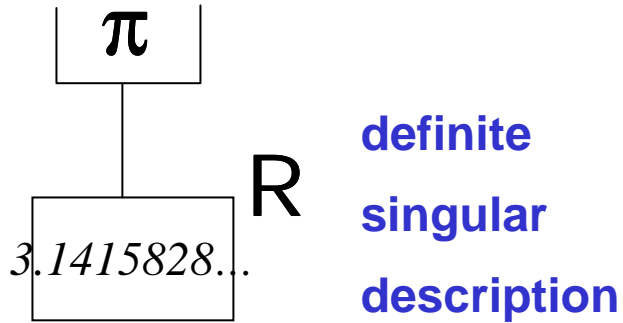
Variables and Assignment

Variable \neq Variable

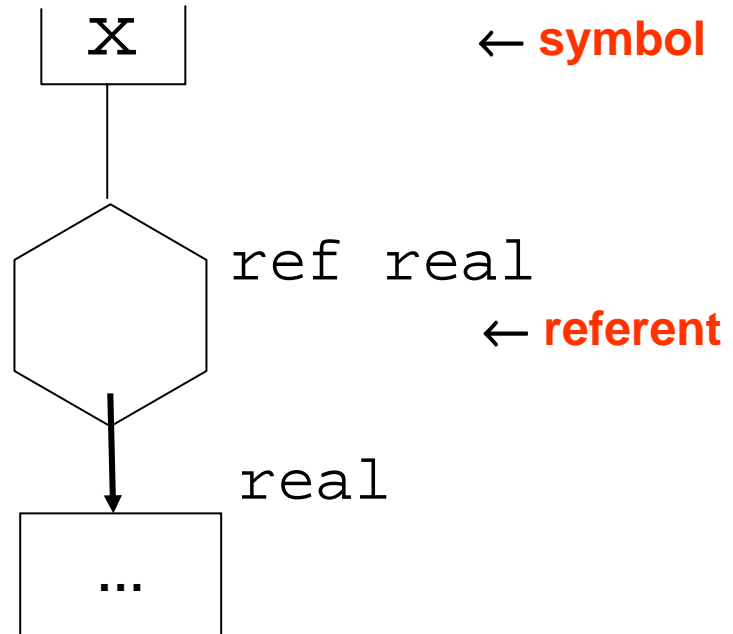
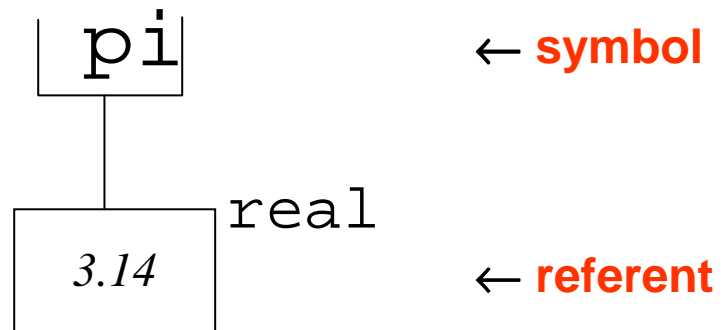
(Math)

(CS)

- **Math**



- **CS**

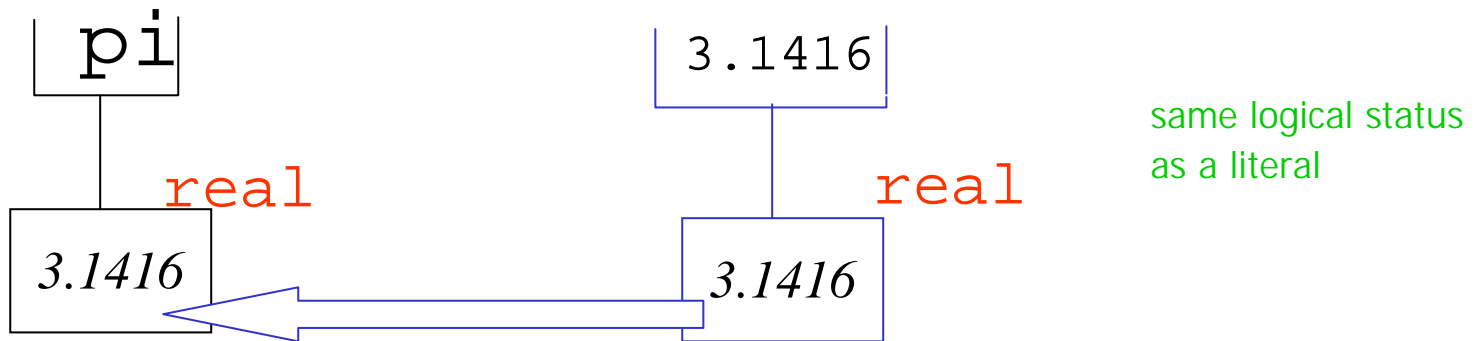


• e.g. Prolog & FP (IFP)

• e.g. Pascal & Algol 68

Constants

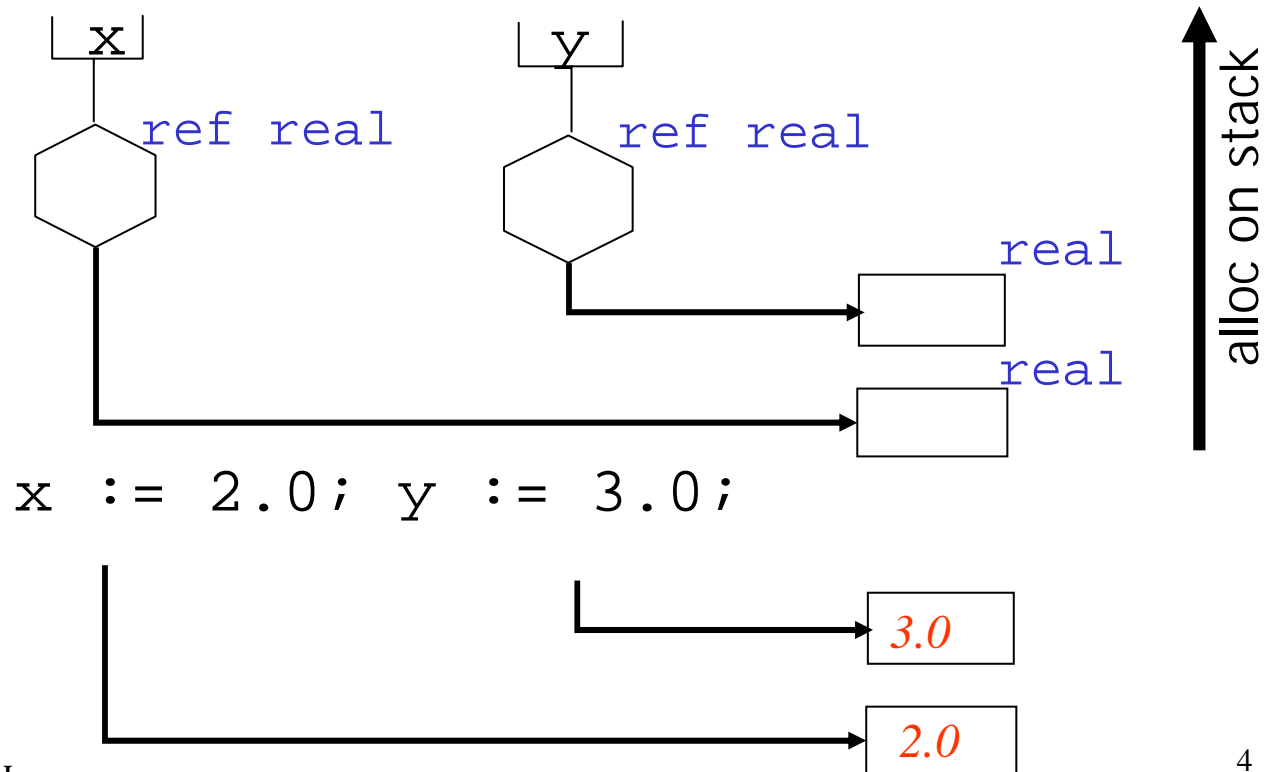
- Algol 68: `real pi = 3.1416;`
- C: (not quite) `#define PI 3.1416`
- Pascal: `const pi = 3.1416;`
- Ada: `pi: constant FLOAT := 3.1416;`



- No reference available
- No `ref real` object exists
- `pi` replaced by literal
- No stack (or any other) space allocated (stored in instruction)
- *Pure definition*

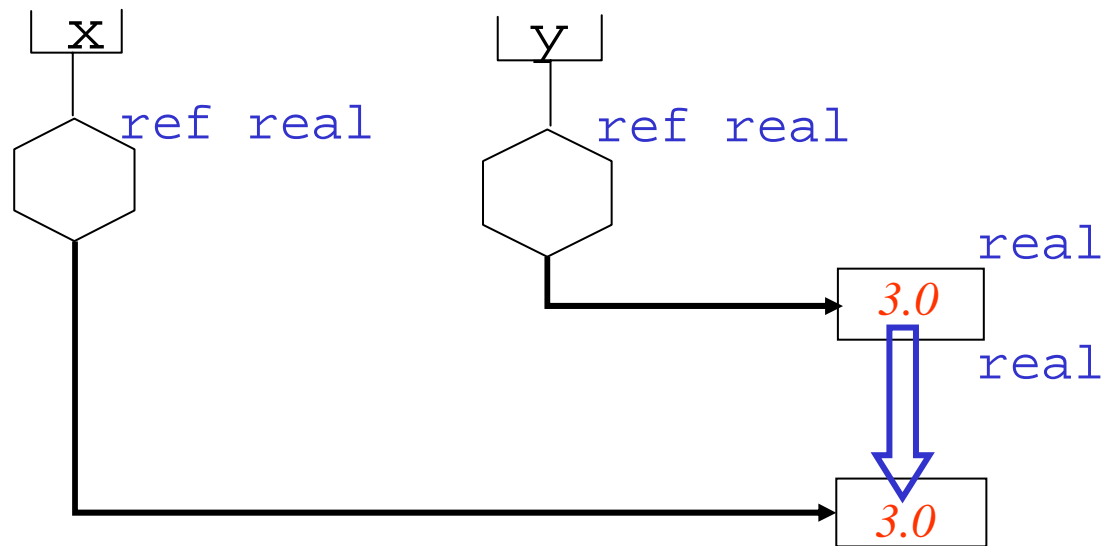
Variables

- Algol 68: `begin real x,y;`
- C: `{ float x,y;`
- Pascal: `var x,y: real;`
- Ada: `x,y: FLOAT;`



Assignment (after y gets 3)

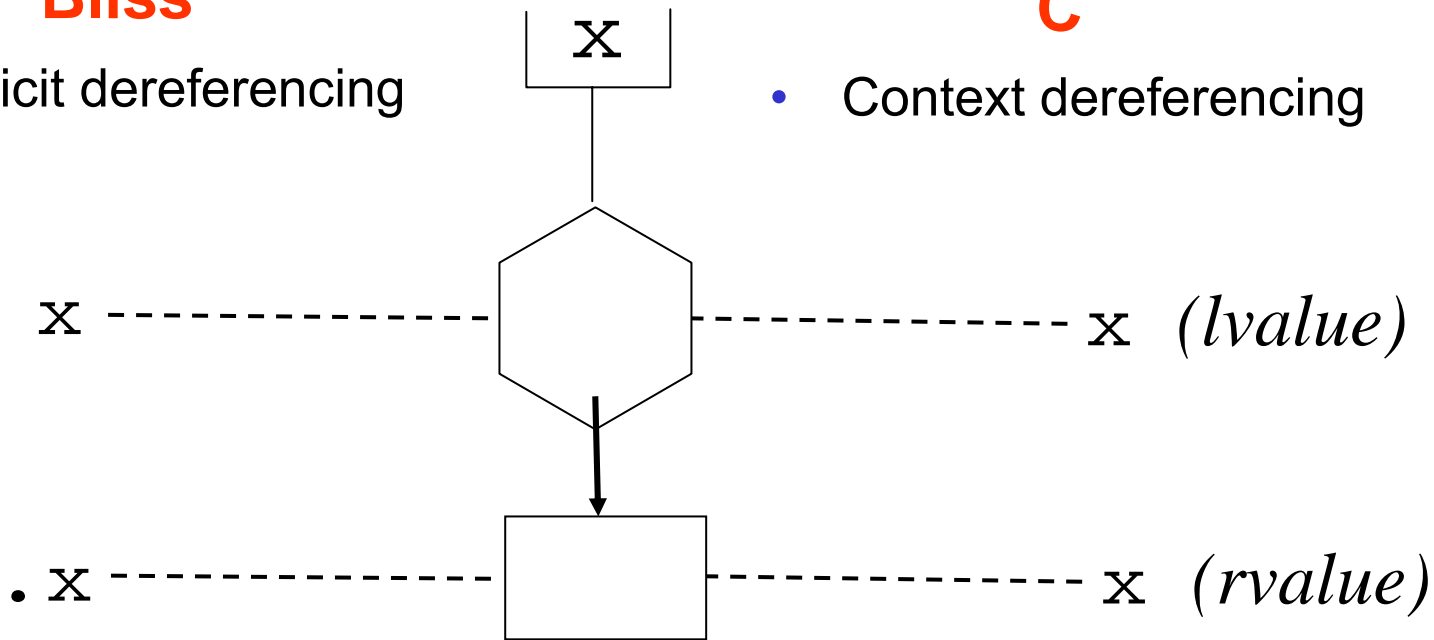
- Algol 68: $x := y;$
- C: $x = y;$
- Bliss: $x \leftarrow .y$



Variables: Bliss v. C

Bliss

- Explicit dereferencing



C

- Context dereferencing

`x = .y + .z`

`x = y + z`

`a + .j ← .(a + .i)`

`a[j] = a[i]`
lvalue *rvalue*

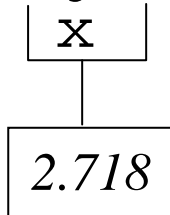
const

Synonym: `*(a+j) = *(a+i)`

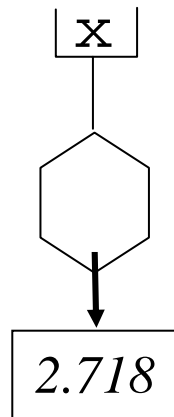
Only lvalue is actually computed

Algol 68 aside: identity declaration

- **real** x := c NOT same as **real** x = c
- Identity declaration **real** x = 2.718; means



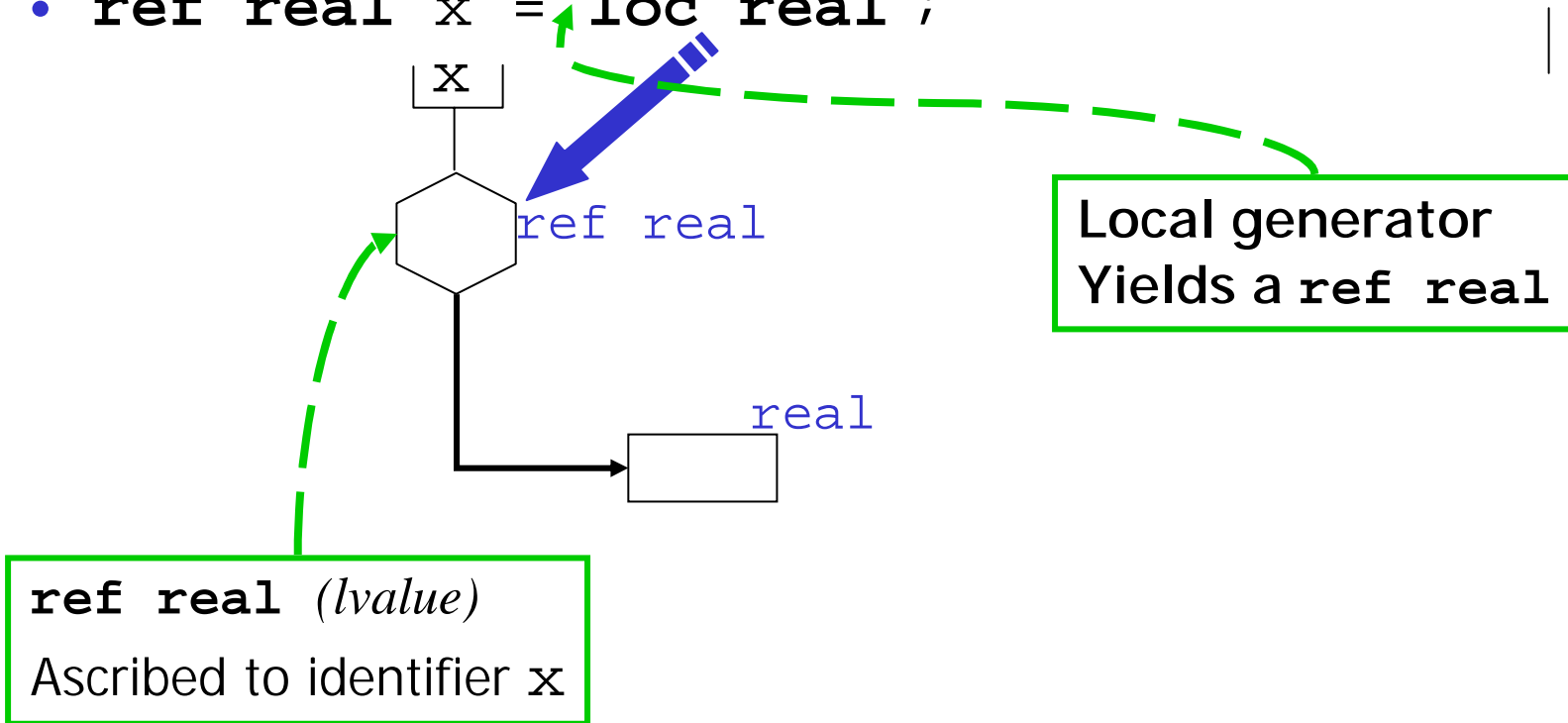
- Initialized variable **real** x := 2.718; means



- Why? Implicit *local* (stack) allocation

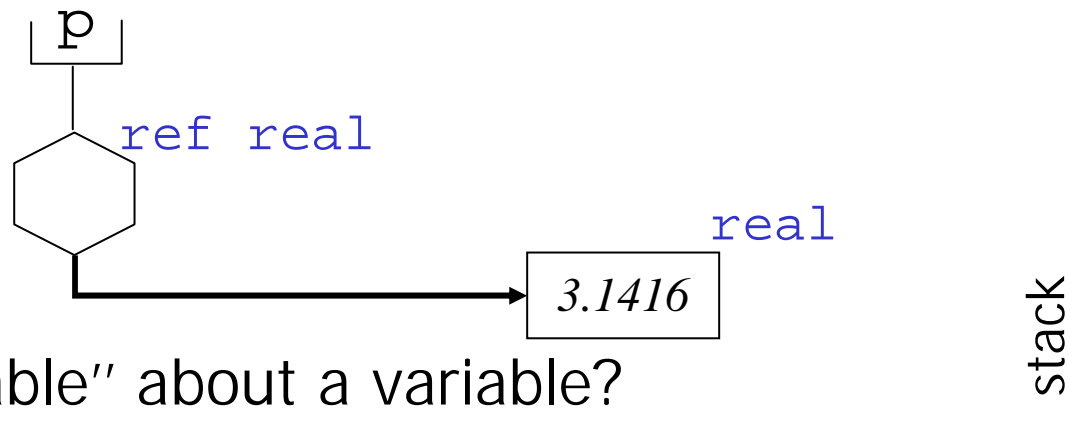
identity declaration (cont'd)

- `real x;` abbreviates `loc real x;` which abbreviates the following identity declaration:
- `ref real x = loc real ;`



Initialized Variables

- Algol 68: **begin** real p := 3.1416;
- C: { float p = 3.1416;
- Pascal: **var** p: real; **begin** p := 3.1416;
- Ada: p: FLOAT := 3.1416;

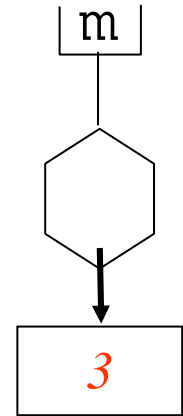
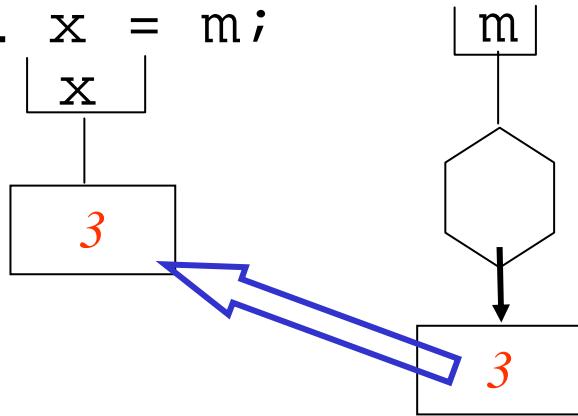


- What is “variable” about a variable?
 - Not reference (location, lvalue)
 - ◆ Ref cannot be changed; “ascribed to” or “possessed by” ident
 - Not identifier
 - Not type (in *most* languages)
 - The value (rvalue)

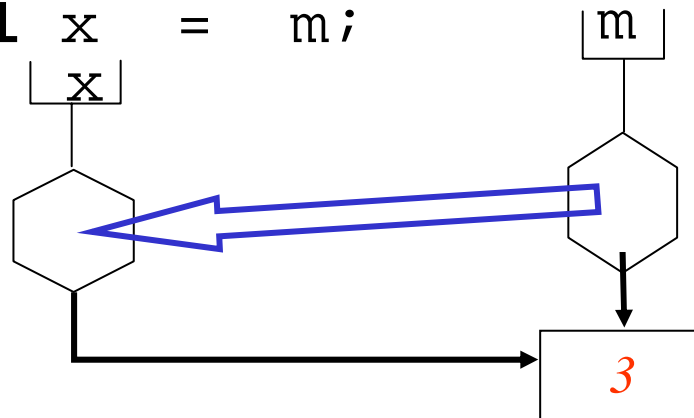
Algol 68 identity declaration: unifies variable & constant declaration

1. `loc real m := 3;`

2. `real x = m;`

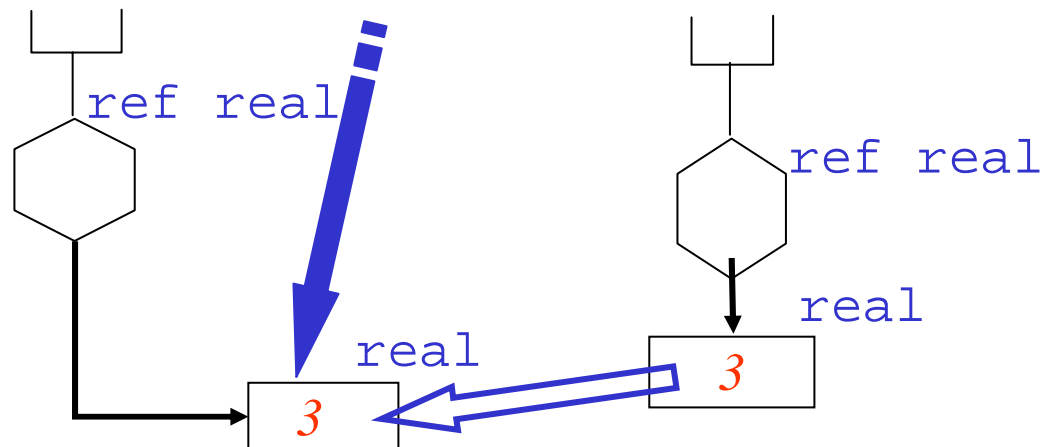


3. `ref real x = m;`

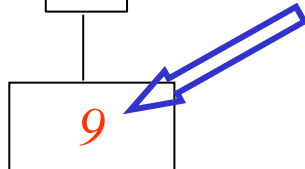


Algol 68 identity declaration (cont'd)

4. `real x := m; ⇒ loc real x:=m; ⇒
ref real x = loc real := m;`



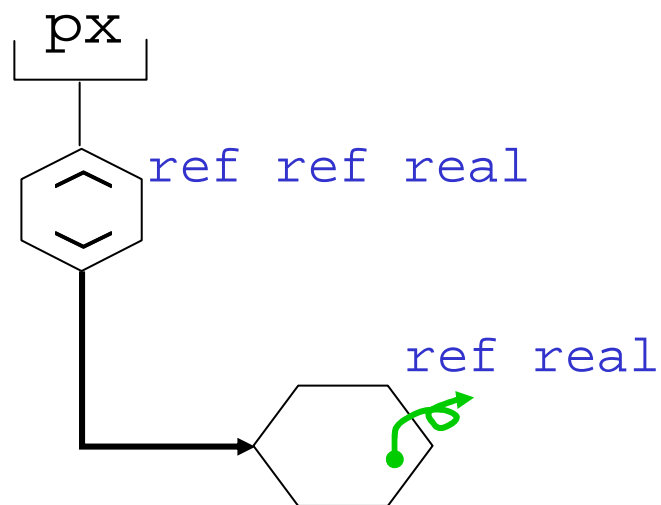
5. `real |t| = x*m;`



- **t const. In new environment**
- **x*m evaluated when declaration encountered**

Pointers

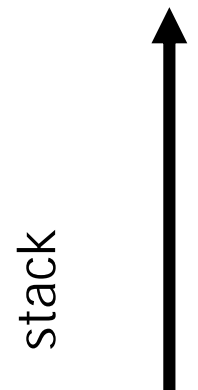
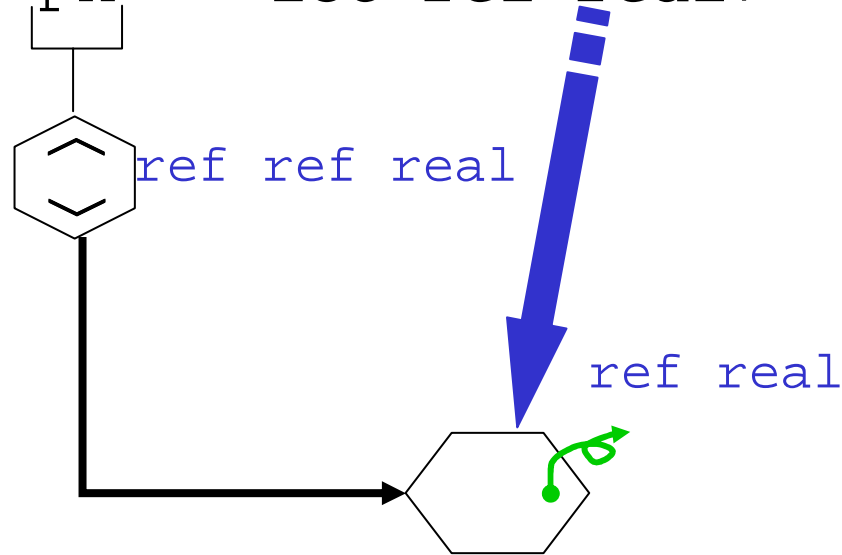
- Refs to variable refs
- Algol 68: `ref real px ;`
- C: `float *px ;`
- Pascal: `var px: ^real;`
- Ada: `type PTR is access FLOAT;`
`px : PTR; (PTR= ``access type'' to ``base type'')`



Algol 68 pointer declaration

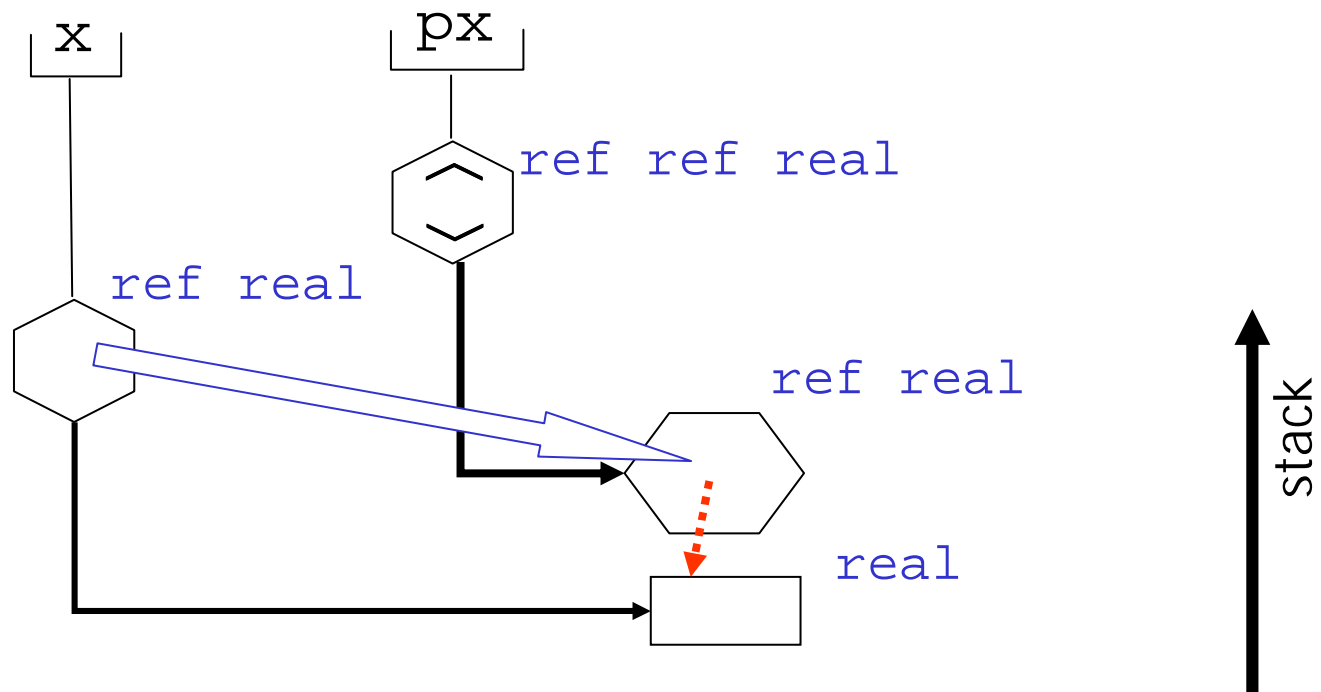
`ref real px` \Rightarrow

`ref ref real px = loc ref real;`



Pointers (cont'd)

- Algol 68: `px := x ;`
- C: `px = &x ;`
- Bliss: `px ← x`

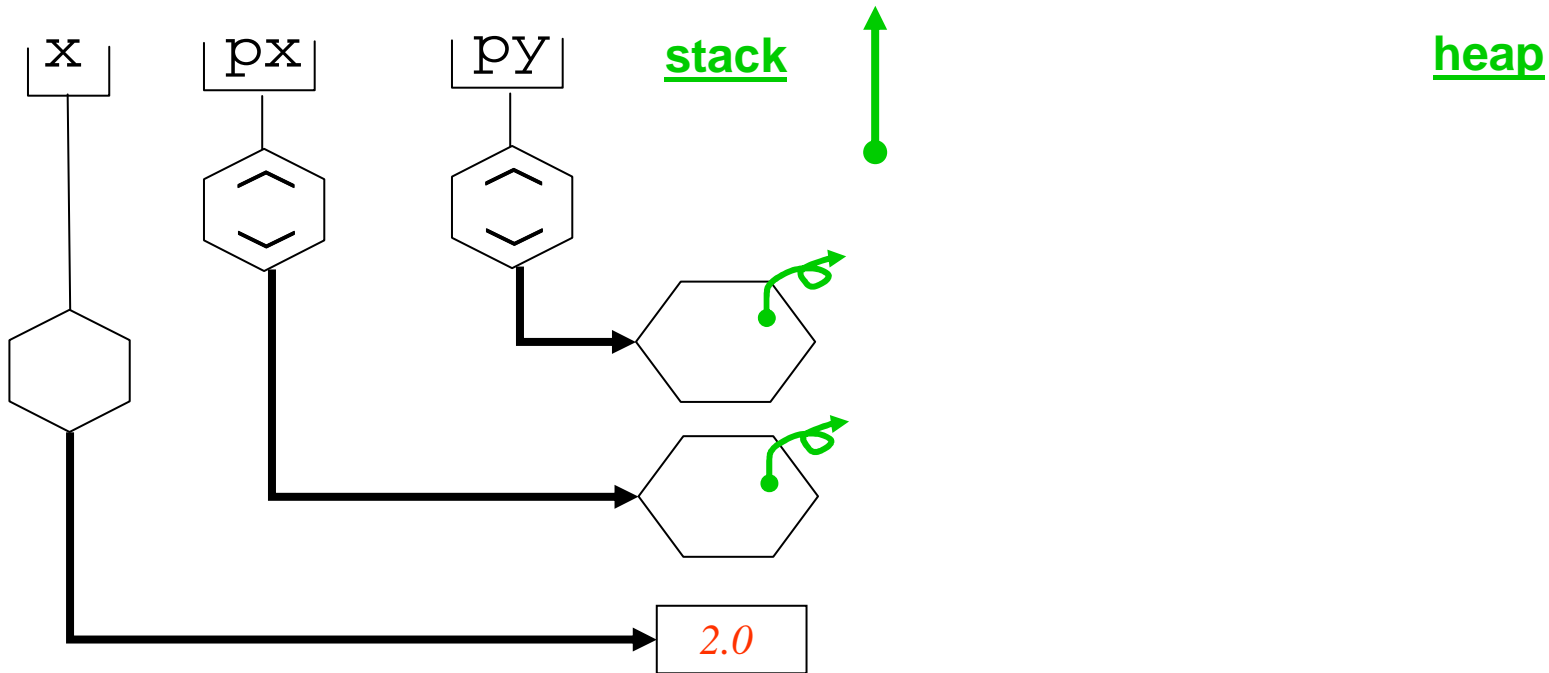


- NOT Pascal, Ada picture. Closest is:
 - Pascal: `new(px); px^ := x;` Ada: `px := new FLOAT(x);`

“anonymous variable”

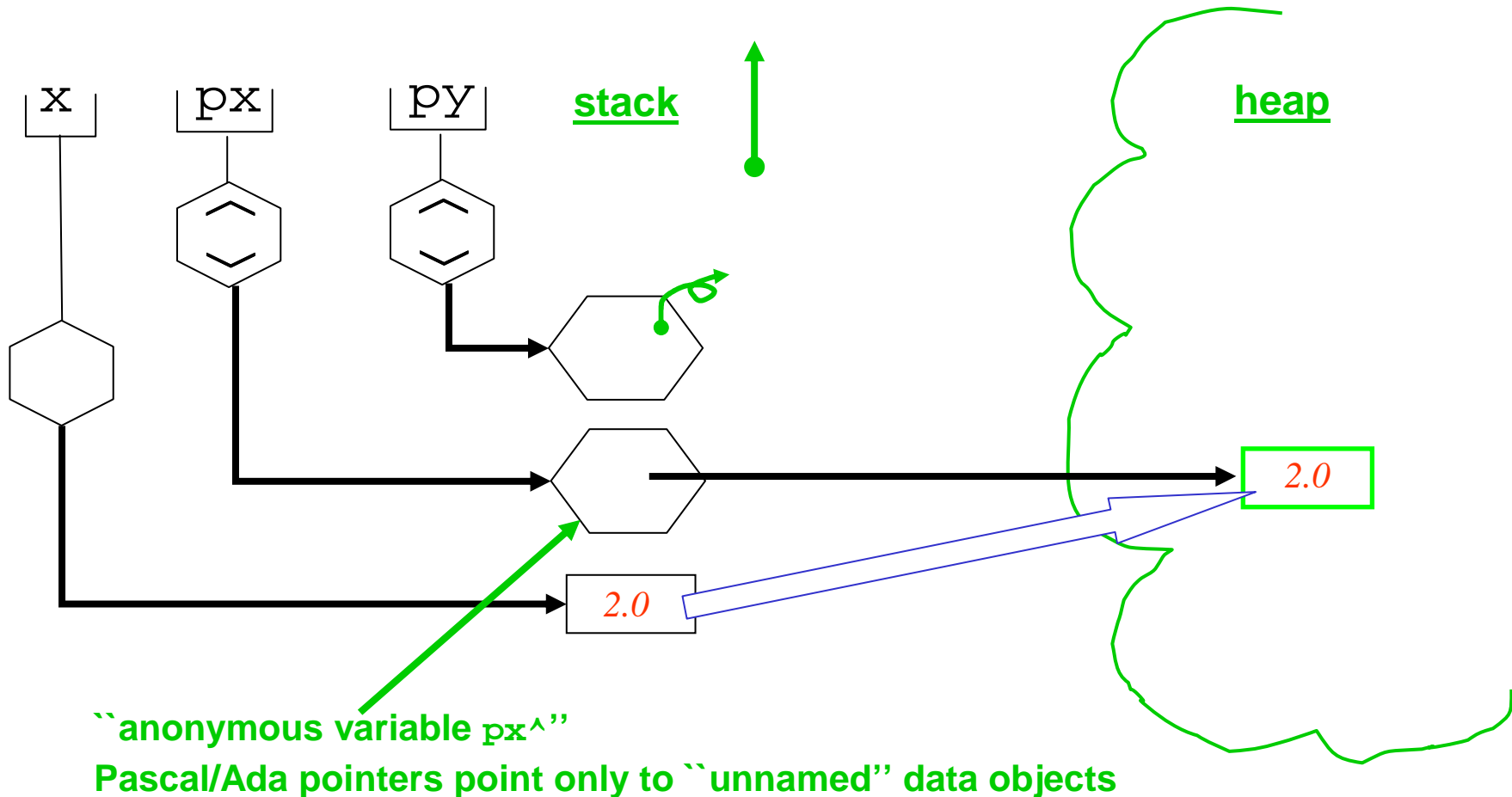
Pointers in Pascal/Ada: point into heap

- Defeats “dangling pointers” at block exit
- Pointers *cannot* point into the stack
- **var** x: real; **var** px, py: ^real;



Pascal/Ada pointers (cont'd)

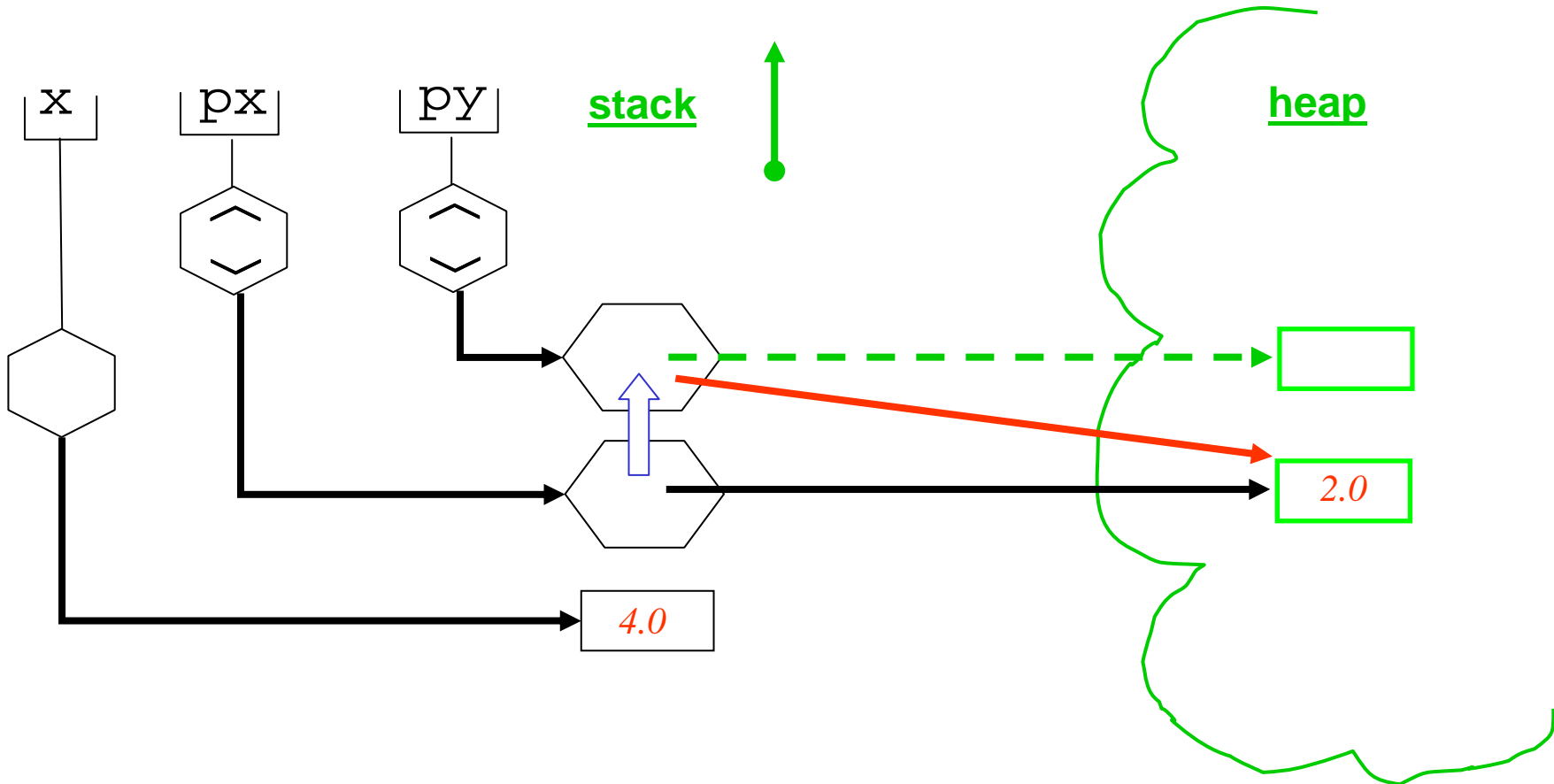
- `new(px); px^ := x;`



Pascal/Ada pointers (cont'd)

- `x := 4.0;`
- `new(py); py := px;`

1.

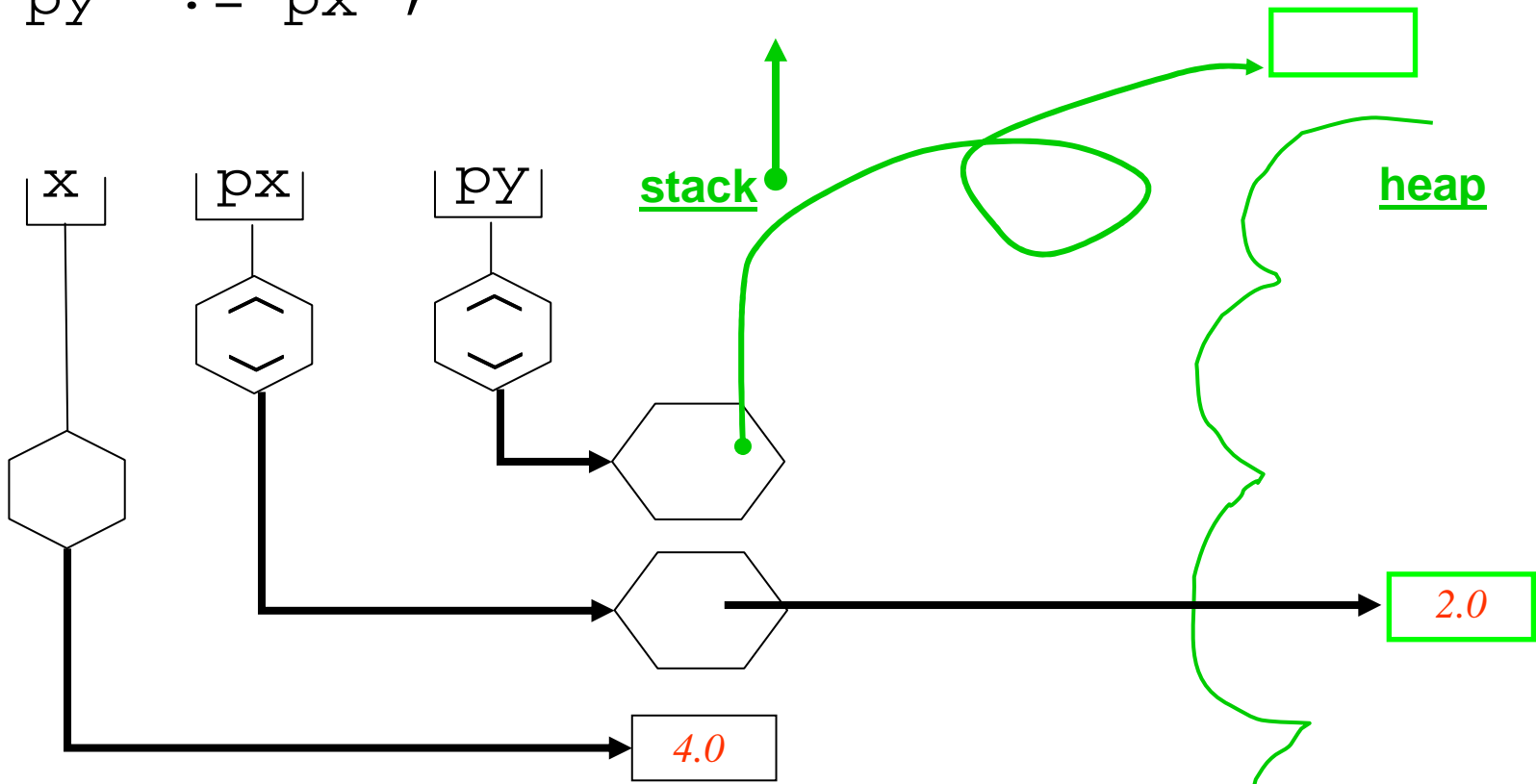


Pascal/Ada pointers (cont'd)

- `x := 4.0;`

2.

- `py^ := px^;`



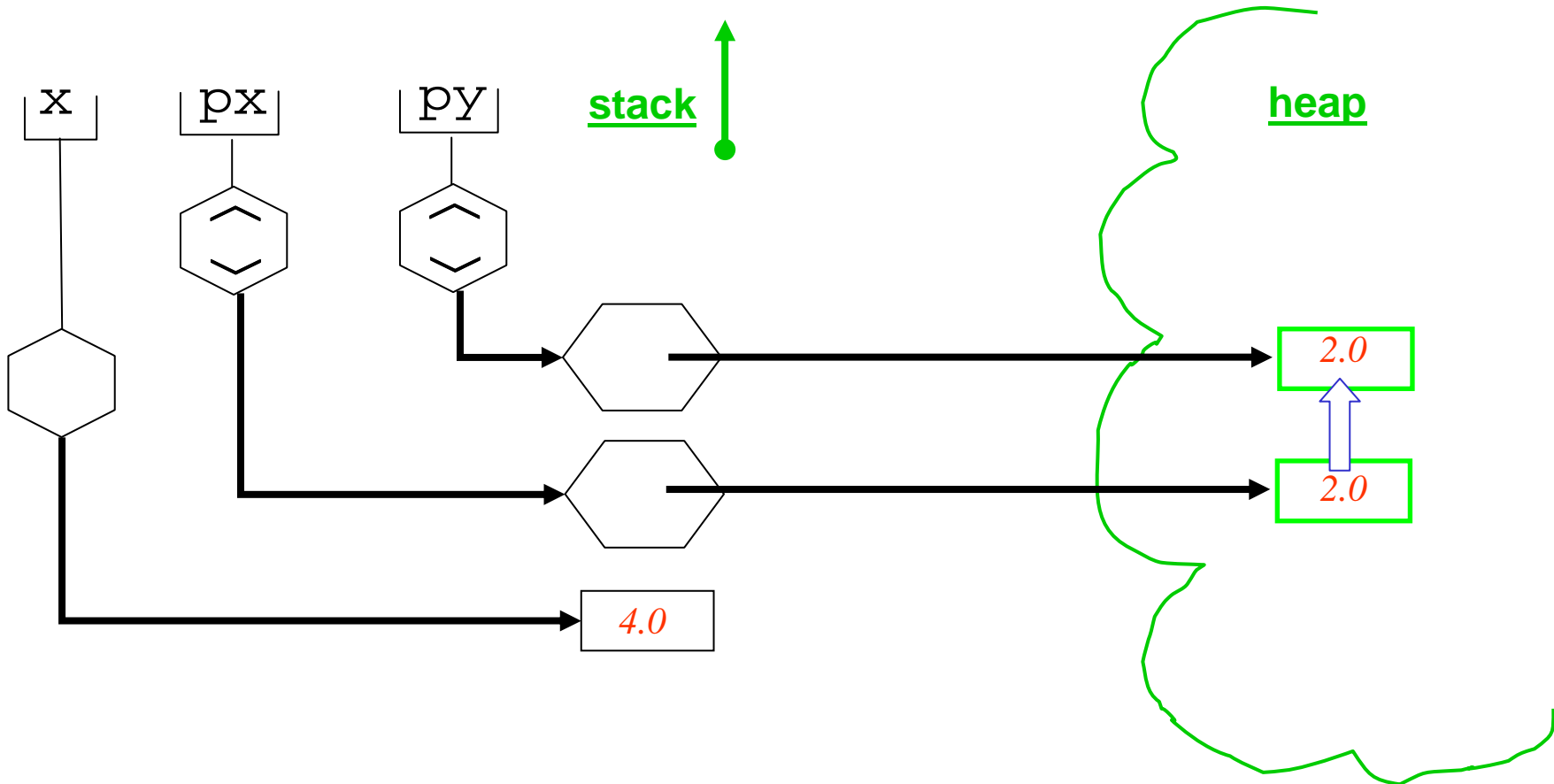
“pointer value out of legal range”

Pascal/Ada pointers (cont'd)

- `x := 4.0;`

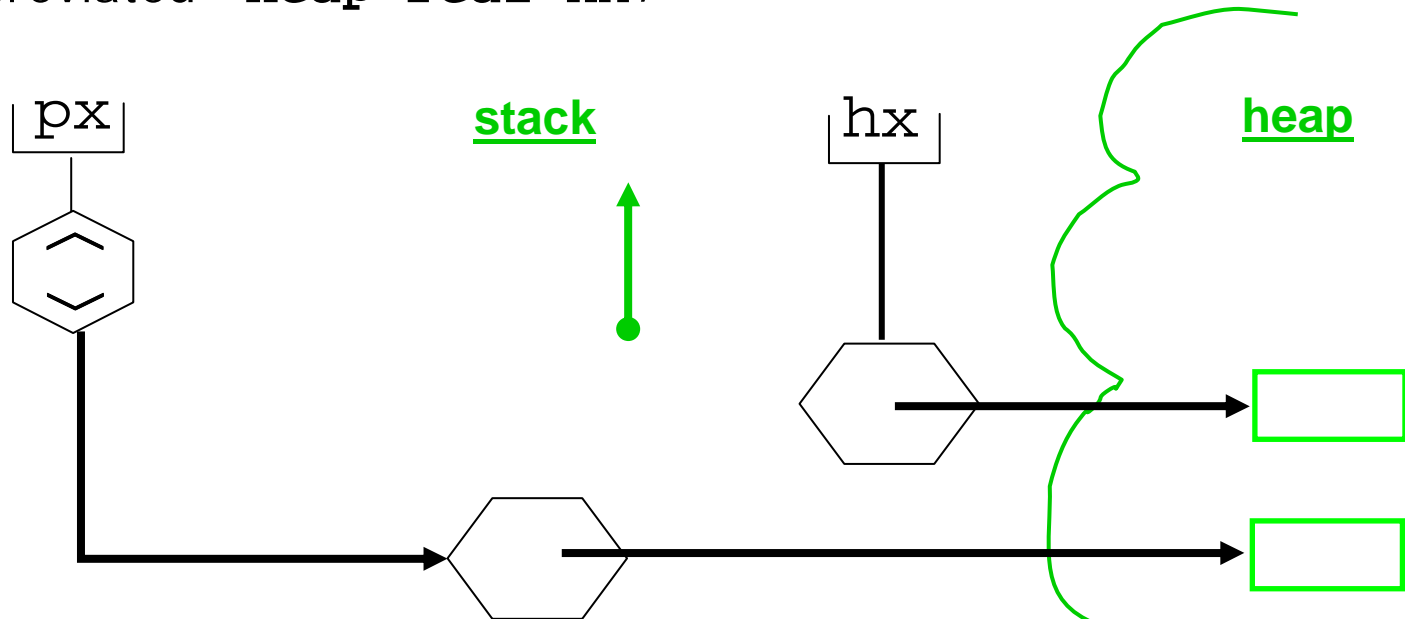
3.

- `new(py); py^ := px^;`



Algol68 has a heap allocator

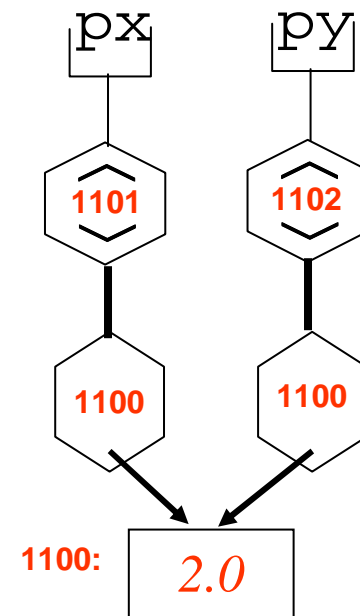
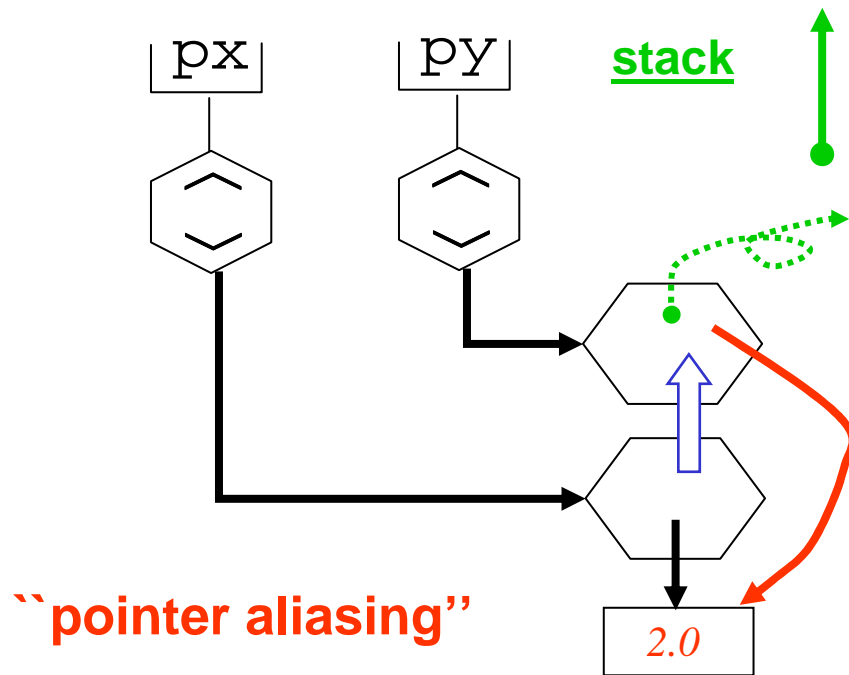
- `ref ref real px = heap real;`
- `ref real hx = heap real;`
 - Abbreviated `heap real hx;`



- Pascal: Ivalues of all var identifiers are in stack
 - Ivalues of all anonymous variables are in the heap
- Algol68: no restriction
 - allocation orthogonal to declaration

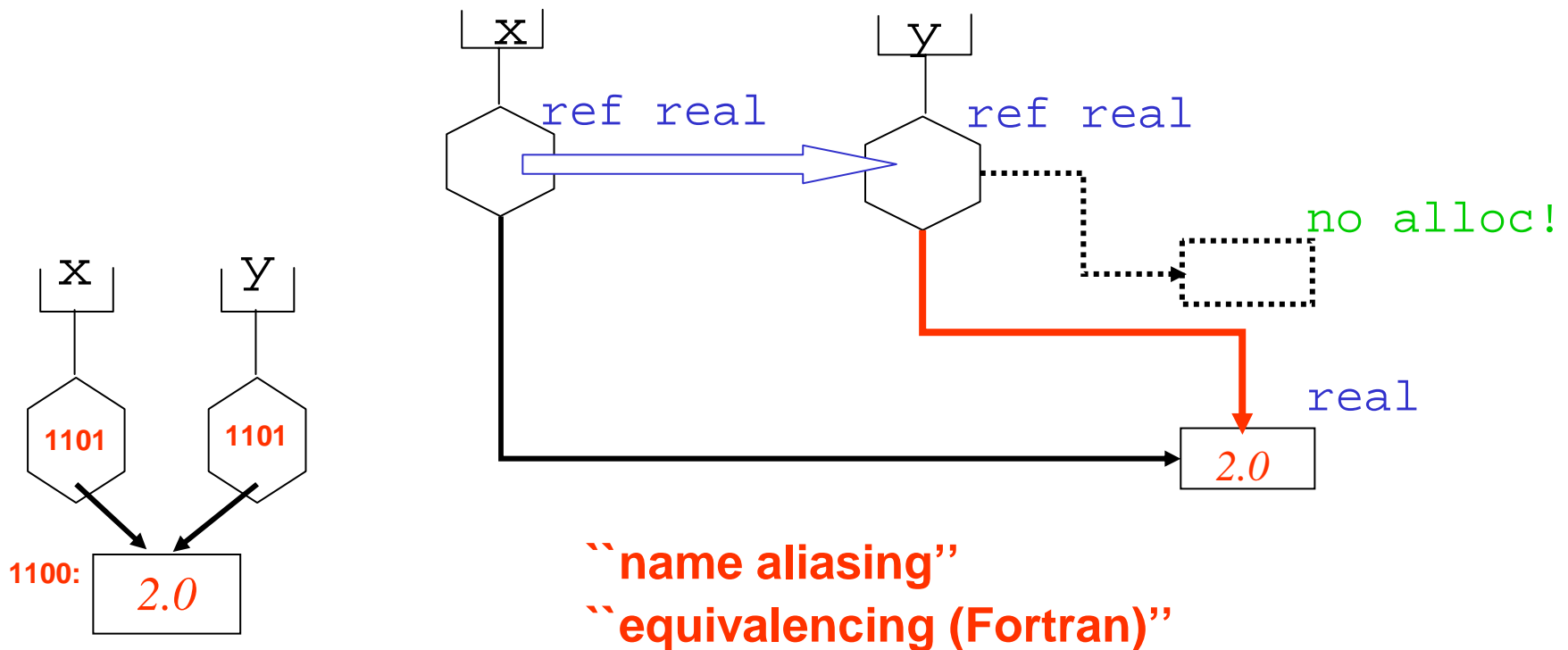
Refs and Aliasing

- Algol68: `ref real py; py := px;`

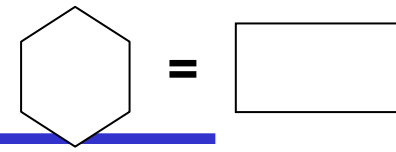


Refs and Aliasing (cont'd)

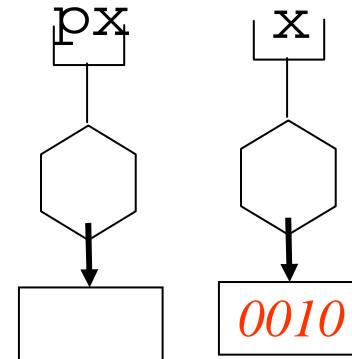
- Algol 68: `ref real y = x;` (identity declaration)



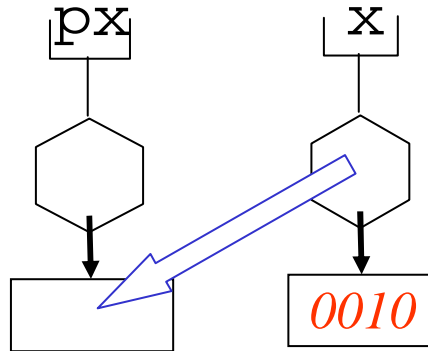
Bliss: untyped bit strings



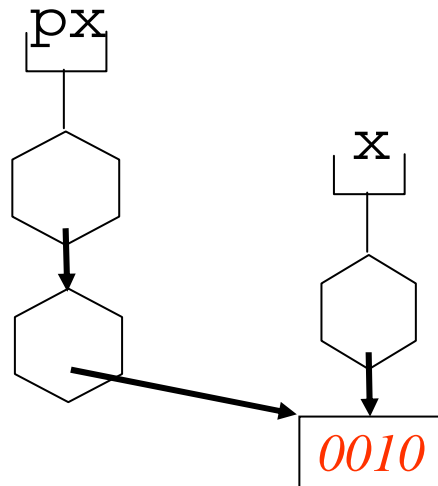
- `local x; x ← 2; local px;`



- `px ← x`



“ptr to x stored in px”



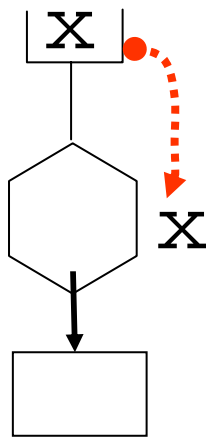
picture after assignment

Semantics of Variable Uses: C v. Bliss

- On LHS \boxed{X} = / ←

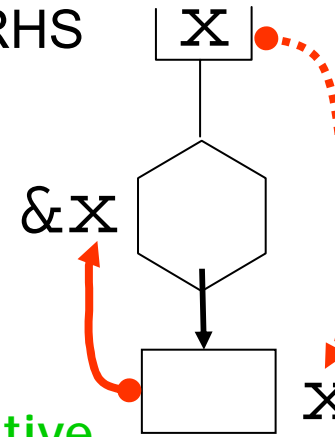
- On RHS \boxed{X}

C:



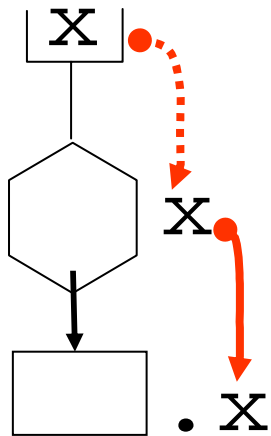
= / ←

=

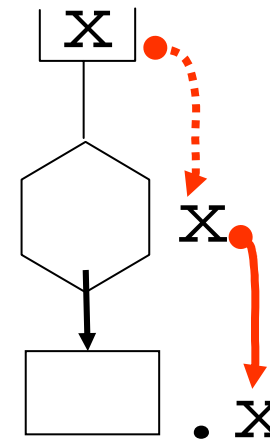


Meaning of x context-sensitive
Not ``referentially transparent''

Bliss:



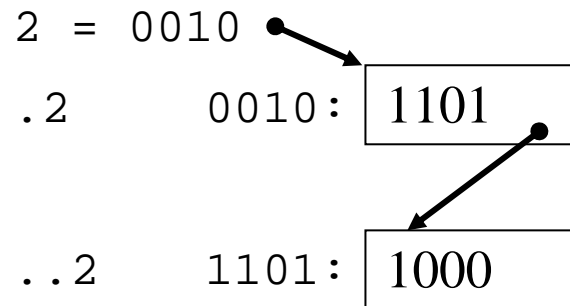
←



Meaning of x is ``referentially transparent''

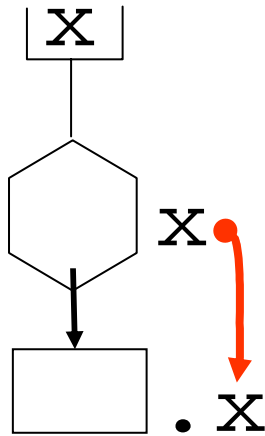
Variable Uses: C v. Bliss (cont'd)

- **C:** **&** **not** a unary operator
 - $\&(x+y)$ meaningless
 - $\&\&x$ meaningless (can't reverse a pointer)
 - $\&2$ meaningless
- **Bliss:** **•** **is** a unary operator
 - $\bullet(x+y)$ sensible
 - $\bullet \bullet x$ sensible
 - $\bullet 2$ sensible

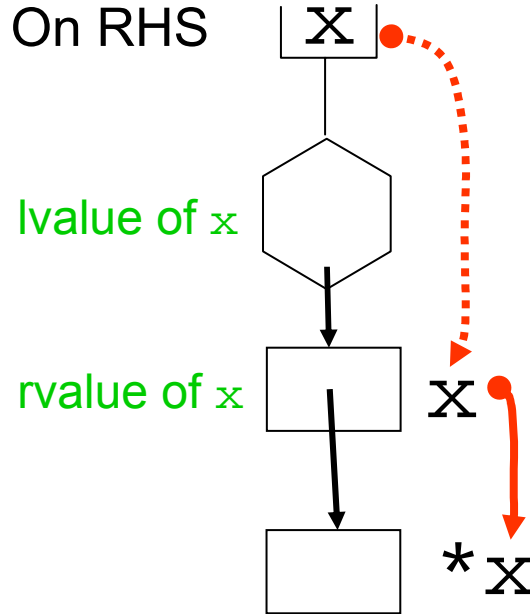


C * ≠ Bliss.

Bliss:



C: • On RHS



Bliss variable occurrences are “referentially transparent”
Semantics of x is independent of context